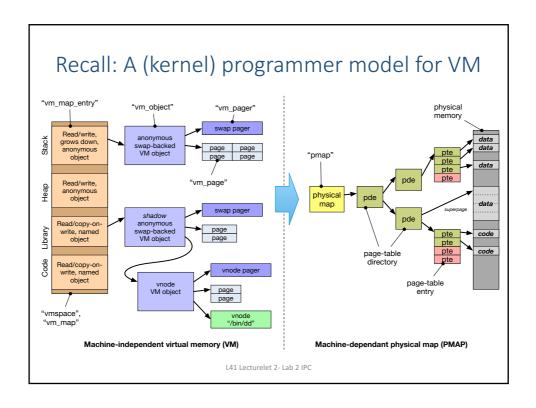
L41: Lab 2- IPC

Lecturelet 2
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L41: Lab 2 – Kernel implications of IPC

- A quick note on vm fault()
- Learn about (and trace) POSIX IPC
- Explore buffering and scheduler interactions
- Measure the probe effect
- This is the first of two labs contributing to Lab Report 2:
 - Lab 2 takes an OS-centric approach
 - Lab 3 takes a microarchitecture-centric approach
- Use data from both to write the lab report



The Mach VM fault handler (vm fault)

- Key goal of the Mach VM system: be as lazy as possible
 - Fill pages (with file data, zeroes, COW) on demand
 - Map pages into address spaces on demand
 - Flush TLB as infrequently as possible
- Any work avoided means reduced CPU cycles and less disk I/O
- Avoid as much work as possible when creating a mapping (e.g., mmap (), execve())
- Instead, do on-demand in the MMU trap handler,
 vm fault()
 - Machine-independent function drives almost all VM work
 - Input: faulting virtual address, output mapped page or signal
 - Look up object to find cached page; if none, invoke pager
 - May trigger behaviour such as zero filling or copy-on-write
- A good thing to probe with DTrace to understand VM traps

The benchmark

- Simple, bespoke IPC benchmark: pipes and sockets
- · Statically or dynamically linked
- Adjust user and kernel buffer sizes
- Various output modes

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The benchmark (2)

Three operational modes:

1thread IPC within a single thread of a single process 2thread IPC between two threads of a single process 2proc IPC between two threads in two processes

- Adjust IPC parameters:
 - -i pipe Use pipe() IPC
 - -i local Use socketpair() IPC
 - -b size Set user IPC buffer size
 - -t size Set total size across all IPCs
 - −s Also set in-kernel buffer size for sockets
 - -B Suppress quiescence (whole-program tracing)
- Output flags:
 - -q Suppress all output (whole-program tracing)
 - ¬∨ Verbose output (interactive testing)

The benchmark (3)

root@141-beaglebone ~/ipc:~ # ./ipc-static -v -i
pipe 1thread

Benchmark configuration:
 buffersize: 131072

totalsize: 16777216 blockcount: 128 mode: 1thread ipctype: pipe time: 0.033753791 485397.29 KBytes/sec

- Use verbose output
- Use pipe IPC
- Run bench mark in a single thread
- Use default buffersize of 128K, totalsize of 16M

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Experimental questions for the lab report

The full lab-report assignment will be distributed during the next lab.

The following questions are intended to help you gather data that you will need for that lab report:

- How does changing the buffer size affect IPC performance and why? For sockets, consider both with, and without, the –s flag.
- Is using multiple threads faster or slower than using multiple processes?

python-dtrace memory leak

```
# The benchmark has completed - stop
the DTrace instrumentation
```

```
dtrace_thread.stop()
dtrace_thread.join()
dtrace_thread.consumer.__del__()
```

- Memory leak in python-dtrace results in instability
- Work around by adding an explicit call to: dtrace thread.consumer. del ()
- Thanks to James Clarke for spotting the issue

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This lab session

- Use this session to continue to build experience:
 - Build and use the IPC benchmark
 - Use DTrace to analyse distributions of system calls, system-call execution times, and system-call arguments and return values
 - Use Jupyter/Python to analyse benchmark results
- Remember to consider the hypotheses the experimental questions are exploring.
- Use the tools in the most productive way:
 - Command line DTrace for quick exploration.
 - Jupyter for data capture, visualisation and analysis.
- Do ask us if you have any questions or need help