The Network Stack (1)

L41 Lecture 5
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Reminder: where we left off last term

- Long, long ago, but in a galaxy not so far away:
 - Lecture 3: The Process Model (1)
 - Lecture 4: The Process Model (2)
 - Lab 2: IPC buffer size and the probe effect
 - Lab 3: Micro-architectural effects of IPC
- Explored several implied (and rejected) hypotheses:
 - Larger IO/IPC buffer sizes amortise system-call overhead
 - A purely architectural view dominates performance
 - The probe effect doesn't matter in real-world workloads

This time: Introduction to Network Stacks

Rapid tour across hardware and software:

- Networking and the sockets API
- Network-stack design principles: 1980s vs. today
- Memory flow across hardware and software
- Network-stack construction and work flows
- Recent network-stack research

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Networking: A key OS function (1)

- Communication between computer systems
 - Local-Area Networks (LANs)
 - Wide-Area Networks (WANs)
- A network stack provides:
 - Sockets API and extensions
 - Interoperable, feature-rich, high-performance protocol implementations (e.g., IPv4, IPv6, ICMP, UDP, TCP, SCTP, 802.1, 802.11, ...)
 - Security functions (e.g., cryptographic tunneling, firewalls...)
 - Device drivers for Network Interface Cards (NICs)
 - Monitoring and management interfaces (BPF, ioctl)
 - Plethora of support libraries (e.g., DNS)

Networking: A key OS function (2)

• Dramatic changes over 30 years:

1980s: Early packet-switched networks, UDP+TCP/IP, Ethernet

1990s: Large-scale migration to IP; Ethernet VLANs

2000s: 1-Gigabit, then 10-Gigabit Ethernet; 802.11; GSM data

2010s: Large-scale deployment of IPv6; 40/100-Gbps Ethernet

... billions → trillions of devices?

- Vanishing technologies
 - UUCP, IPX/SPX, ATM, token ring, SLIP, ...

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The Berkeley Sockets API (1983)

close()
read()
write()
...
accept()
bind()
connect()
getsockopt
listen()
recv()
soloct()

getsockopt()
listen()
recv()
select()
send()
setsockopt()
socket()
...

- The Design and Implementation of the 4.3BSD Operating System
 - (but APIs/code first appeared in 4.2BSD)
- Now universal TCP/IP (POSIX, Windows)
- Kernel-resident network stack serves networking applications via system calls
- Reuses file-descriptor abstraction
 - Same API for local and distributed IPC
 - Simple, synchronous, copying semantics
 - Blocking/non-blocking I/O, select()
- Multi-protocol (e.g., IPv4, IPv6, ISO, ...)
 - TCP-focused but not TCP-specific
 - · Cross-protocal abstractions and libraries
 - Protocol-specific implementations
 - "Portable" applications

BSD network-stack principles (1980s-1990s)

Multi-protocol, packet-oriented network research framework:

- Object-oriented: multiple protocols, socket types, but one API
 - Protocol-independent: streams vs. datagrams, sockets, socket buffers, socket addresses, network interfaces, routing table, packets
 - Protocol-specific: connection lists, address/routing specialization, routing, transport protocol itself – encapsulation, decapsulation, etc.

• Packet-oriented:

- Packets and packet queueing as fundamental primitives
- Best effort: If there is a failure (overload, corruption), drop the packet
- · Work hard to maintain packet source ordering
- Differentiate 'receive' from 'deliver' and 'send' from 'transmit'
- Heavy focus on TCP functionality and performance
- Middle-node (forwarding), not just edge-node (I/O), functionality
- High-performance packet capture: Berkeley Packet Filter (BPF)

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FreeBSD network-stack principles (1990s-2010s)

All of the 1980s features and also ...

• Hardware:

- Multi-processor scalability
- NIC offload features (checksums, TSO/LRO, full TCP)
- Multi-queue network cards with load balancing/flow direction
- Performance to 10s or 100s of Gigabit/s
- · Wireless networking

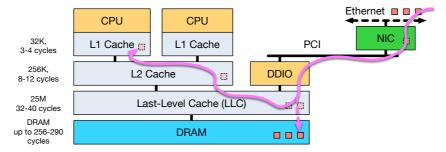
Protocols:

- Dual IPv4/IPv6
- · Pluggable congestion control
- Security/privacy: firewalls, IPSec, ...

Software model:

- Flexible memory model integrates with VM for zero-copy
- Network-stack virtualisation
- Userspace networking via netmap

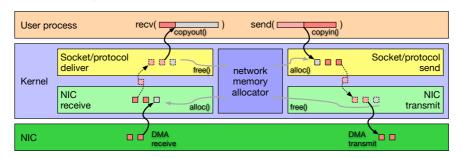
Memory flow in hardware



- Key idea: follow the memory
 - Historically, memory copying is avoided due to instruction count
 - Today, memory copying is avoided due to cache footprint
- Recent Intel CPUs push and pull DMA via the LLC ("DDIO")
 - If we differentiate 'send' and 'transmit', 'receive' vs. 'deliver', is this a good idea?
 - ... it depends on the latency between DMA and processing

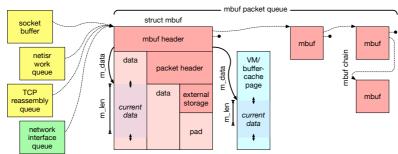
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Memory flow in software



- Socket API implies **one software-driven copy** to/from user memory
 - Historically, zero-copy VM tricks for socket API ineffective
- Network buffers cycle through the slab allocator
 - Receive: allocate in NIC driver, free in socket layer
 - Transmit: allocate in socket layer, free in NIC driver
- DMA performs second copy; can affect cache/memory bandwidth
 - NB: what if packet-buffer working set is larger than the cache?

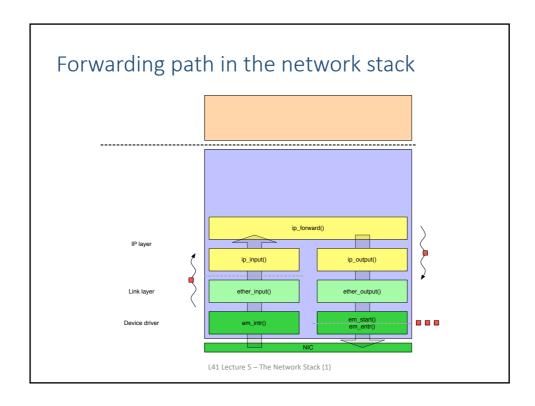
The mbuf abstraction

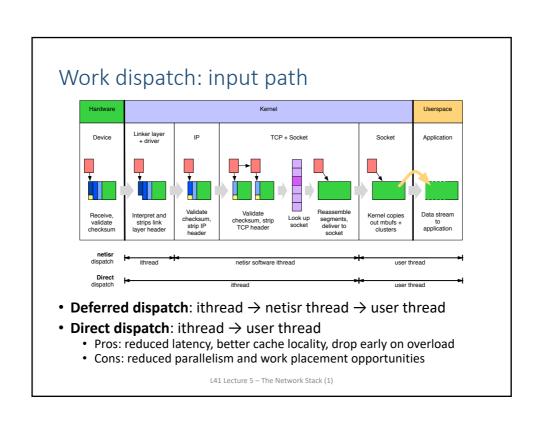


- Unit of work allocation and distribution throughout the stack
- mbuf chains represent in-flight packets, streams, etc.
 - Operations: alloc, free, prepend, append, truncate, enqueue, dequeue
 - Internal or external data buffer (e.g., VM page)
 - Reflects bi-modal packet-size distribution (e.g., TCP ACKs vs data)
- Similar structures in other OSes e.g., skbuff in Linux

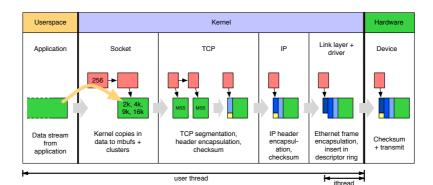
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Send/receive paths in the network stack Application recv() send() System call layer soreceive() scaend() Socket layer soreceive() scaend() TCP layer top_reas() top_reas() IP layer ip_input() ip_output() Device driver em_intr() em_ em_start() NIC L41 Lecture 5 – The Network Stack (1)





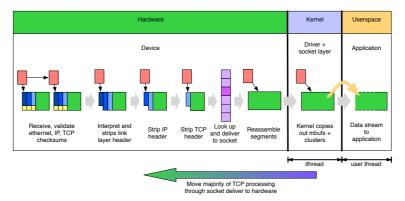
Work dispatch: output path



- Fewer deferred dispatch opportunities implemented
 - (Deferred dispatch on device-driver handoff in new iflib KPIs)
- · Gradual shift of work from software to hardware
 - Checksum calculation, segmentation, ...

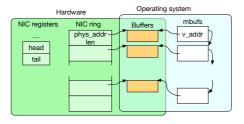
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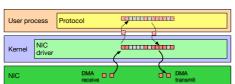
Work dispatch: TOE input path



- Kernel provides socket buffers and resource allocation
- Remainder, including state, retransmissions, etc., in NIC
- But: two network stacks? Less flexible/updateable structure?
 - Better with an explicit HW/SW architecture e.g., Microsoft Chimney

Netmap: a novel framework for fast packet I/O Luigi Rizzo, USENIX ATC 2012 (best paper).



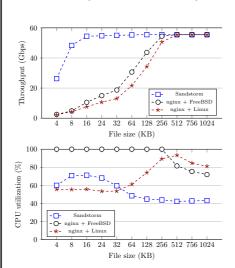


- Map NIC buffers directly into user process memory
- Not the sockets API: Zero copy to/from application
- System calls initiate DMA, block for NIC events
- Packets can be reinjected into normal stack
- Ships in FreeBSD; patch available for Linux
- Userspace network stack can be **specialised** to task (e.g., packet forwarding)

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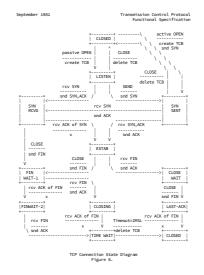
Network stack specialisation for performance

Ilias Marinos, Robert N. M. Watson, Mark Handley, SIGCOMM 2014, 2017.



- 30 years since the networkstack design developed
- Massive changes in architecture, microarchitecture, memory...
 - Optimising compilers
 - · Cache-centered CPUs
 - Multiprocessing, NUMA
 - DMA, multiqueue
 - 10 Gigabit/s Ethernet
- Performance lost to 'generality' throughout stack
- Revisit fundamentals through clean-slate stack
- Orders-of-magnitude performance gains

Next time: TCP protocol + implementation



- McKusick, et al: Chapter 14 (*Transport-Layer Protocols*)
- Transmission Control Protocol (TCP)
- TCP implementation
 - Buffers and input processing
 - Parallelism and performance
 - DoS resistance
- The final two labs