

Unique existence

The notation

$$\exists! x. P(x)$$

stands for

the *unique existence* of an x for which the property $P(x)$ holds .

That is,

$$\exists x. P(x) \wedge \underbrace{\left(\forall y. \forall z. (P(y) \wedge P(z)) \implies y = z \right)}_{\text{Uniqueness}}$$

Disjunction

Disjunctive statements are of the form

$$P \text{ or } Q$$

or, in other words,

either P , Q , or both hold

or, in symbols,

$$P \vee Q$$

The main proof strategy for disjunction:

To prove a goal of the form

$$P \vee Q$$

you may

1. try to prove P (if you succeed, then you are done); or
2. try to prove Q (if you succeed, then you are done);
otherwise
3. break your proof into cases; proving, in each case,
either P or Q .

Proposition 25 For all integers n , either $n^2 \equiv 0 \pmod{4}$ or $n^2 \equiv 1 \pmod{4}$.

PROOF: $\forall \text{int } n. [n^2 \equiv 0 \pmod{4}] \vee [n^2 \equiv 1 \pmod{4}]$

Assume arbitrary int. n .

Try to show: $n^2 \equiv 0 \pmod{4}$ ~~X~~

Try to show: $n^2 \equiv 1 \pmod{4}$ ~~X~~

We consider two cases:

(1) n even and (2) n is odd.

Case (1) $n = 2k$ for some k

$$n^2 = 4k^2 \equiv 0 \pmod{4} \quad \checkmark$$

Case (2) $n = 2l + 1$ for some l

$$n^2 = (2l + 1)^2 = 4l^2 + 4l + 1$$

$$= 4(l^2 + l) + 1 \equiv 1 \pmod{4}$$



Assumptions

⋮
 $P_1 \vee P_2$
⋮

Goals
 Q

The use of disjunction:

To use a disjunctive assumption

$P_1 \vee P_2$

to establish a goal Q , consider the following two cases in turn: (i) assume P_1 to establish Q , and (ii) assume P_2 to establish Q .

<u>Assume</u>	<u>Goal</u>		<u>Assume</u>	<u>Goal</u>
⋮ P_1	Q		⋮ P_2	Q .

Scratch work:

Before using the strategy

Assumptions

Goal

Q

\vdots

$P_1 \vee P_2$

After using the strategy

Assumptions

Goal

Q

\vdots

P_1

Assumptions

Goal

Q

\vdots

P_2

Proof pattern:

In order to prove Q from some assumptions amongst which there is

$$P_1 \vee P_2$$

write: We prove the following two cases in turn: (i) that assuming P_1 , we have Q ; and (ii) that assuming P_2 , we have Q . Case (i): Assume P_1 . **and provide a proof of Q from it and the other assumptions.** Case (ii): Assume P_2 . **and provide a proof of Q from it and the other assumptions.**

$$\binom{p}{m} = \frac{p!}{m!(p-m)!} \quad p \in \mathbb{N}$$

A little arithmetic

notations: C_m^p
 ${}^p C_m$

Lemma 27 For all positive integers p and natural numbers m , if $m = 0$ or $m = p$ then $\binom{p}{m} \equiv 1 \pmod{p}$.

PROOF: \forall pos. int p . \forall nat. m .

$$(m=0 \vee m=p) \Rightarrow \binom{p}{m} \equiv 1 \pmod{p}$$

Assume p arbitrary pos. int. and m arbitrary natural number.

Assume: $(m=0 \vee m=p)$

RTP: $\binom{p}{m} \equiv 1 \pmod{p}$

We proceed by cases.

(1) Assume $m=0$.

$$\text{Then } \binom{p}{m} = \binom{p}{0} = 1 \equiv 1 \pmod{p}$$

(2) Assume $m=p$

$$\text{Then } \binom{p}{m} = \binom{p}{p} = 1 \equiv 1 \pmod{p}$$

NB

\equiv is a predicate and $n \equiv n \pmod{k}$
for all n and k ☐

Lemma 28 For all integers p and m , if p is prime and $0 < m < p$ then $\binom{p}{m} \equiv 0 \pmod{p}$.

PROOF: Assume p integer and m integer.

Assume p is prime and $0 < m < p$.

RTP: $\binom{p}{m} \equiv 0 \pmod{p} \Leftrightarrow \binom{p}{m}$ is a multiple of p .

$$\binom{p}{m} = \frac{p!}{m!(p-m)!} = p \cdot \left[\frac{(p-1)!}{m!(p-m)!} \right]$$

$\binom{p}{m} = p \cdot k$ and we are done. $k \times$
We need show k is an integer!

$$\binom{p}{m} = p \cdot \frac{(p-1)!}{m! (p-m)!} \quad m \cdot (m-1) \cdot \dots \cdot 1$$

So

$$p \cdot (p-1)! = \binom{p}{m} \cdot m! (p-m)!$$

need to argue/show that $p \mid \binom{p}{m}$

p prime and $p \nmid m! (p-m)! \Rightarrow p \mid \binom{p}{m}$
 Euclid's Theorem.

Proposition 29 For all prime numbers p and integers $0 \leq m \leq p$, either $\binom{p}{m} \equiv 0 \pmod{p}$ or $\binom{p}{m} \equiv 1 \pmod{p}$.

PROOF: Let p be a prime and m be an integer between 0 and p .

RTP: $\binom{p}{m} \equiv 0 \pmod{p} \vee \binom{p}{m} \equiv 1 \pmod{p}$

(1) Case $m=0$: $\binom{p}{m} \equiv 1 \pmod{p}$ by Prop 27

(2) Case $0 < m < p$: $\binom{p}{m} \equiv 0 \pmod{p}$ by Lemma 28

(3) Case $m=p$: $\binom{p}{m} \equiv 1 \pmod{p}$ by Prop 27. ◻

A little more arithmetic

Corollary 33 (The Freshman's Dream) For all natural numbers m , n and primes p ,

$$(m + n)^p \equiv m^p + n^p \pmod{p} .$$

PROOF: Let m, n be nat. Let p prime.

$$\begin{aligned} (m+n)^p &= \sum_{i=0}^p \binom{p}{i} m^i n^{p-i} \\ &= m^p + n^p + \sum_{i=1}^{p-1} \binom{p}{i} m^i n^{p-i} \end{aligned}$$

$$\forall i=1, \dots, p-1. \quad \binom{p}{i} \equiv 0 \pmod{p}$$

$$\Rightarrow \binom{p}{i} m^i n^{p-i}$$

$$\equiv 0 \cdot m^i \cdot n^{p-i}$$

$$\equiv 0 \pmod{p}$$

$$\Rightarrow \sum_{i=1}^{p-1} \binom{p}{i} m^i n^{p-i} \equiv 0 \pmod{p}$$

Lemma

$$a \equiv b \pmod{m}$$

$$a' \equiv b' \pmod{m}$$

\Rightarrow

$$a \cdot a' \equiv b \cdot b' \pmod{m}$$

$$a + a' \equiv b + b' \pmod{m}$$



Corollary 34 (The Dropout Lemma) For all natural numbers m and primes p ,

$$(m + 1)^p \equiv m^p + 1 \pmod{p} .$$

Proposition 35 (The Many Dropout Lemma) For all natural numbers m and i , and primes p ,

$$(m + i)^p \equiv m^p + i \pmod{p} .$$

PROOF: Idea: $(m + i)^p = \left(m + \underbrace{(1 + 1 + \dots + 1)}_{i \text{ times}} \right)^p$
 $= \left((m + 1 + 1 + \dots + 1) + 1 \right)^p$

$$\begin{aligned} & \overbrace{(m+1+\dots+1)}^i{}^p \\ & \equiv \underbrace{(m+1+\dots+1)}_{i-1}{}^p + 1 \end{aligned}$$

Formally
one proceeds
by induction!

$$\equiv \underbrace{(m+1+\dots+1)}_{i-2}{}^p + \underbrace{1+1}_2$$

...

$$\equiv \underbrace{(m + \underbrace{1+\dots+1}_{i-(i-1)})}_{i}{}^p + \underbrace{1+\dots+1}_{i-1}$$

$$\equiv m^p + \underbrace{1+\dots+1}_i = m^p + i$$



The Many Dropout Lemma (Proposition 35) gives the first part of the following very important theorem as a corollary.

Theorem 36 (Fermat's Little Theorem) *For all natural numbers i and primes p ,*

1. $i^p \equiv i \pmod{p}$, and

2. $i^{p-1} \equiv 1 \pmod{p}$ whenever i is not a multiple of p .

The fact that the first part of Fermat's Little Theorem implies the second one will be proved later on .

Btw

1. Fermat's Little Theorem has applications to:
 - (a) primality testing^a,
 - (b) the verification of floating-point algorithms, and
 - (c) cryptographic security.

^aFor instance, to establish that a positive integer m is not prime one may proceed to find an integer i such that $i^m \not\equiv i \pmod{m}$.