

# Binary products

In a category  $\mathcal{C}$ , a **product** for objects  $X, Y \in \mathcal{C}$  is a diagram  $X \xleftarrow{\pi_1} P \xrightarrow{\pi_2} Y$  with the universal property:

for all  $X \xleftarrow{f} Z \xrightarrow{g} Y$   
there is a unique  $h: Z \rightarrow P$

such that

commutes

# Binary products

In a category  $\mathcal{C}$ , a **product** for objects

$X, Y \in \mathcal{C}$  is a diagram  $X \xleftarrow{\pi_1} P \xrightarrow{\pi_2} Y$

satisfying

$(P, \pi_1, \pi_2)$  is terminal in the category with  
- objects  $(Z, f, g)$  where  $X \xleftarrow{f} Z \xrightarrow{g} Y$  in  $\mathcal{C}$

- morphisms  $h: (Z, f, g) \rightarrow (Z', f', g')$  are

$h \in \mathcal{C}(Z, Z')$  such that  $f' \circ h = f$  &  $g' \circ h = g$

- composition & identities as in  $\mathcal{C}$

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so if they exist, products are unique up to (unique) isomorphism.

# Binary products - notation

Usual notation for product of  $X$  &  $Y$  is

$$X \xleftarrow{\pi_1} X \times Y \xrightarrow{\pi_2} Y$$

and, given  $X \xleftarrow{f} Z \xrightarrow{g} Y$ , the  
unique  $h: Z \rightarrow X \times Y$  with  $\begin{cases} \pi_1 \circ h = f \\ \pi_2 \circ h = g \end{cases}$

will be written

$$\langle f, g \rangle : Z \rightarrow X \times Y$$

# Examples of products

In Set :  $X \times Y = \{(x, y) \mid x \in X \wedge y \in Y\}$

$$\pi_1(x, y) = x$$

$$\pi_2(x, y) = y$$

because ...

# Examples of products

$$\text{In Pre} : (P, \leq) \times (Q, \leq)$$

$$= (P \times Q, \leq) \quad \text{product in Set}$$

$$(p_1, q_1) \leq (p_2, q_2) \triangleq p_1 \leq p_2 \wedge q_1 \leq q_2$$

in P                      in Q

$$\pi_1(p, q) = p$$

$$\pi_2(p, q) = q$$

} are monotone functions

# Examples of products

In Mon :  $(M, \cdot, e) \times (N, \cdot, e)$

$$= (M \times N, \cdot, (e, e))$$

product  
in Set

$$(m_1, n_1) \cdot (m_2, n_2) \stackrel{\Delta}{=} (m_1 \cdot m_2, n_1 \cdot n_2)$$

in M

in N

unit for this is  $(e, e)$

$$\pi_1(m, n) = m$$

$$\pi_2(m, n) = n$$

} give monoid homomorphisms

# Examples of products

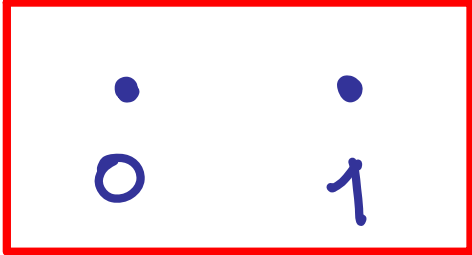
In a pre-ordered set  $(P, \leq)$ , regarded as a Category, the product of  $p, q \in P$  is a greatest lower bound (glb, or meet)  $p \wedge q$  :

$$p \wedge q \leq p \quad \& \quad p \wedge q \leq q$$

$$(\forall r \in P) \quad r \leq p \quad \& \quad r \leq q \quad \Rightarrow \quad r \leq p \wedge q$$



# Non-example

The poset , that is

$(\{0, 1\}, \leq)$  where  $0 \leq 0$  &  $1 \leq 1$

does not possess a product (=meet)  
for  $0$  &  $1$ .

# Duality

Binary product in  $\mathcal{C}^{\text{op}}$  is called  
binary **Coproduct** in  $\mathcal{C}$ .

E.g. coproduct of  $X, Y \in \text{Set}$  is

$$X \xrightarrow{i_1} X \amalg Y \xleftarrow{i_2} Y$$

$$i_1(x) \triangleq (x, 0)$$

$$i_2(y) \triangleq (y, 1)$$

$$\triangleq \{(x, 0) \mid x \in X\} \cup \{(y, 1) \mid y \in Y\}$$