$$gcd$$
fun gcd(m, n) $M = q \cdot n + r$

$$= let$$
val (q, r) = divalg(m, n)
in
if r = 0 then n $CD(m_{1}n) = CD(q \cdot n_{1}n) = D(n)$
else gcd(n, r)
end $CD(m_{1}n) = CD(n_{1}r)$
 $Key Lemma$

Theorem 60 Euclid's Algorithm gcd terminates on all pairs of positive integers and, for such m and n, gcd(m, n) is the greatest common divisor of m and n in the sense that the following two properties hold:

- (i) both gcd(m, n) | m and gcd(m, n) | n, and
- (ii) for all positive integers d such that d | m and d | n it necessarily follows that d | gcd(m, n).

PROOF: [(i) drd (ii)] dre aquivallent xercit $(\forall d. d|m, d|n \Leftrightarrow d|gcd(m,n))$ CD(m, n) = D(gcd(m, n))which is equivalent

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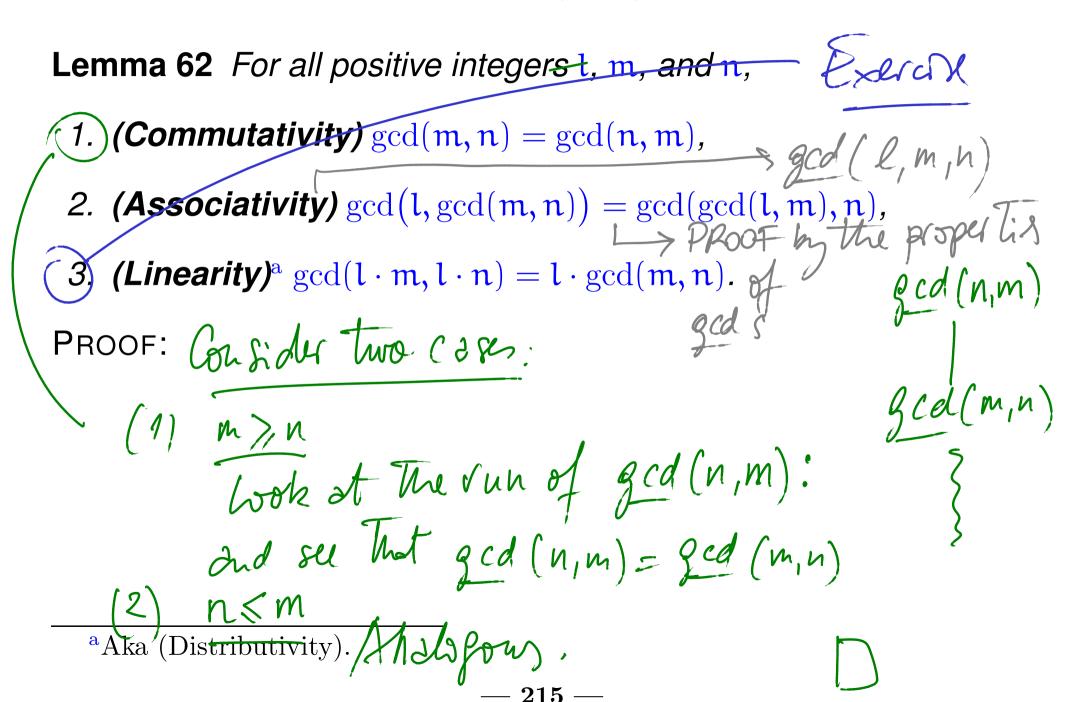
=>q'=1=)r'= n-r < n/2 Run "AR'L N/ gcd (Frith Frith) gcd(m, n) $m = q \cdot n + r$ r/n/m for Filoshacci 0 < m < n6 q > 0, 0 < r < nmun by g. gcd(n,m)gcd(n, r)n $n = q' \cdot r + r'$ r|n|q' > 0, 0 < r' < rqcd (m,n) gcd(r,r')is $O(log(n \rightarrow (m, n)))$ How moh shaller is r Than n? r n/2 r1 CAN gcd ('m,n) $r' \leq n/2$ in O(max(m,n))

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Fractions in lowest terms

```
fun lowterms( m , n )
= let
    val gcdval = gcd( m , n )
    in
    ( m div gcdval , n div gcdval )
    end
```

Some fundamental properties of gcds



Corollary: Pellinia for parime and OKMKp (Exercise) **Theorem 63** For positive integers k, m, and n, if k | (m · n) and gcd(k, m) = 1 then $k \mid n$. PROOF: Let k, m, n be pritive in tipers. Assume Okl(m.n) and Oged(k,m)=1 RTP: KIN By (2) ond linedity: n.gcd (k,m)=n $gcd(n\cdot k, n\cdot m)$ By (i) $m \cdot n = k \cdot q$ for some integer q. Here, $n = g \cdot cd(n,k,k \cdot q) \stackrel{\text{by lineasity}}{=} k \cdot g \cdot cd(n,q)$.

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Corollary 64 (Euclid's Theorem) For positive integers m and n, and prime p, if $p \mid (m \cdot n)$ then $p \mid m$ or $p \mid n$.

Now, the second part of Fermat's Little Theorem follows as a corollary of the first part and Euclid's Theorem.

PROOF: Let M, M be portier integers and p & prime. Assume pl(m.n) By coses : (1) plm and we are done. $(2) \neg (p|m)$ The ped(p,m) = 1and by the previous the pln.

$$NB: k \equiv [k]_{m} (mdm)$$

$$a \equiv b (modm) \Longrightarrow k \cdot a \equiv k \cdot b (modm)$$

$$\begin{cases} 0, 1, \dots, p-1 \\ 0, \dots, p-1 \\$$

Extended Euclid's Algorithm

Example 67 (egcd(34, 13) = ((5, -13), 1))

	gcd(34 , 13)	34	=	2.	13	+	8
=	gcd(13 , 8)	13	=	1.	8	+	5
—	gcd(8,5)	8	=	1.	5	+	3
=	gcd(5,3)	5	=	1.	3	+	2
=	gcd(3,2)	3	=	1.	2	+	1
=	gcd(2 , 1)	2	=	2.	1	+	0

= 1

Extended Euclid's Algorithm

Example 67 (egcd(34, 13) = ((5, -13), 1))

	gcd(34 , 13)	34	=	2.	13	+	8	8	=	34	-2.	13
=	gcd(13 , 8)	13	=	1.	8	+	5	5	=	13	-1.	8
=	gcd(8 , 5)	8	=	1.	5	+	3	3	=	8	-1.	5
=	gcd(5,3)	5	—	1.	3	+	2	2	=	5	-1.	3
=	gcd(3,2)	3	—	1.	2	+	1	1	=	3	-1.	2
=	gcd(2,1)	2	=	2.	1	+	0					

= 1

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— 232-а —

— 232-c —

$$\begin{array}{l} \underbrace{\operatorname{gcd}(m,n)}_{i} = l_{1} \cdot m + l_{2} \cdot n & \operatorname{fn} \operatorname{shne} l_{i} \operatorname{ad} l_{2} \\ \operatorname{gcd}(34,13) \\ = \operatorname{gcd}(34,13) \\ = \operatorname{gcd}(13,8) \\ = \operatorname{gcd}(13,8) \\ = \operatorname{gcd}(8,5) \\ = \operatorname{gcd}(8,5) \\ = \operatorname{gcd}(8,5) \\ = \operatorname{gcd}(8,5) \\ = \operatorname{gcd}(5,3) \\ = \operatorname{gcd}(5,3) \\ = \operatorname{gcd}(3,2) \\ \operatorname{fn} \operatorname{fn$$