

## Big unions

**Definition 86** Let  $U$  be a set. For a collection of sets  $\mathcal{F} \in \mathcal{P}(\mathcal{P}(U))$ , we let the big union (relative to  $U$ ) be defined as

$$\bigcup \mathcal{F} = \{x \in U \mid \exists A \in \mathcal{F}. x \in A\} \in \mathcal{P}(U) .$$

## Big intersections

**Definition 88** Let  $U$  be a set. For a collection of sets  $\mathcal{F} \subseteq \mathcal{P}(U)$ , we let the big intersection (relative to  $U$ ) be defined as

$$\bigcap \mathcal{F} = \{x \in U \mid \forall A \in \mathcal{F}. x \in A\} .$$

$$U = \mathbb{R}$$

closure property

**Theorem 89** Let

$$\mathcal{F} = \left\{ S \subseteq \mathbb{R} \mid (0 \in S) \wedge (\forall x \in \mathbb{R}. x \in S \implies (x+1) \in S) \right\} .$$

Then, (i)  $\mathbb{N} \in \mathcal{F}$  and (ii)  $\mathbb{N} \subseteq \bigcap \mathcal{F}$ . Hence,  $\bigcap \mathcal{F} = \mathbb{N}$ .

PROOF:

- $\bigcap \mathcal{F} \subseteq \mathbb{N}$   $x \in \bigcap \mathcal{F} \implies \forall S \in \mathcal{F}. x \in S$

$\mathbb{R} \in \mathcal{F}$   $\rightsquigarrow$  characterizes  $\mathbb{N}$ .  
 $\mathbb{Q} \in \mathcal{F}$   
 $\mathbb{N} \in \mathcal{F}$

- $\mathbb{N} \subseteq \bigcap \mathcal{F}$  But  $\mathbb{N} \in \mathcal{F}$ , hence  $x \in \mathbb{N}$

$\forall n \in \mathbb{N}. n \in \bigcap \mathcal{F} \iff \forall n \in \mathbb{N}. \boxed{\forall S \in \mathcal{F}. n \in S}$  Exercise

$\equiv P(n)$

By induction on  $n \in \mathbb{N}$ .

Examples

$$\cup \{\} = \emptyset$$

$$\cup \{X\} = X$$

$$\cup \{X, Y\} = X \cup Y$$

$$x \in X \cup Y \text{ iff } x \in X \\ \vee x \in Y$$

**Union axiom**

Every collection of sets has a union.

$\cup \mathcal{F}$

$$x \in \cup \mathcal{F} \iff \exists X \in \mathcal{F}. x \in X$$

For non-empty  $\mathcal{F}$  we also have

$$\bigcap \mathcal{F}$$

defined by

$$\forall x. x \in \bigcap \mathcal{F} \iff (\forall X \in \mathcal{F}. x \in X) .$$

Analogous to  
datatype  $\alpha$   $\beta$  disjoint union = one of  $\alpha$  | two of  $\beta$ .

## Disjoint unions

**Definition 90** The disjoint union  $A \uplus B$  of two sets  $A$  and  $B$  is the set

$$A \uplus B = (\{1\} \times A) \cup (\{2\} \times B) .$$

Thus,

$$A \uplus B = \{ (1, a) \mid a \in A \} \cup \{ (2, b) \mid b \in B \} .$$

$$\forall x. x \in (A \uplus B) \iff (\exists a \in A. x = (1, a)) \vee (\exists b \in B. x = (2, b)) .$$

Remark :  $\{*\} \times X = \{ (*, x) \mid x \in X \}$

**Proposition 92** For all finite sets  $A$  and  $B$ ,

$$A \cap B = \emptyset \implies \#(A \cup B) = \#A + \#B .$$

PROOF IDEA:  $A$  and  $B$  are disjoint.

$$A = \{a_1, \dots, a_m\}$$

$$B = \{b_1, \dots, b_n\}$$

$$\forall i, j. a_i \neq b_j$$

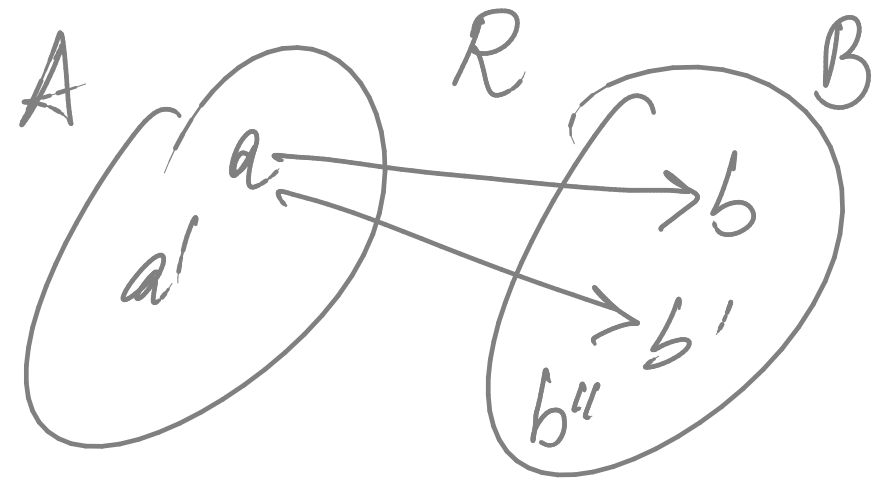
$$A \cup B = \{ \overbrace{a_1, \dots, a_m}^m, \overbrace{b_1, \dots, b_n}^n \} \longleftarrow m+n$$



**Corollary 93** For all finite sets  $A$  and  $B$ ,

$$\#(A \uplus B) = \#A + \#B .$$

$(a, b) \in R$  notation  $a R b$   
 $(a, b') \in R$   $a R b'$



## Relations

**Definition 95** A (binary) relation  $R$  from a set  $A$  to a set  $B$

$$R : A \rightarrow B \quad \text{or} \quad R \in \text{Rel}(A, B) ,$$

is

$$R \subseteq A \times B \quad \text{or} \quad R \in \mathcal{P}(A \times B) .$$

**Notation 96** One typically writes  $a R b$  for  $(a, b) \in R$ .



## Examples:

- ▶ Empty relation.

$$\emptyset : A \dashrightarrow B$$

$$(a \emptyset b \iff \text{false})$$

- ▶ Full relation.

$$(A \times B) : A \dashrightarrow B$$

$$(a (A \times B) b \iff \text{true})$$

- ▶ Identity (or equality) relation.

$$\text{id}_A = \{ (a, a) \mid a \in A \} : A \dashrightarrow A$$

$$(a \text{id}_A a' \iff a = a')$$

- ▶ Integer square root.

$$R_2 = \{ (m, n) \mid m = n^2 \} : \mathbb{N} \dashrightarrow \mathbb{Z}$$

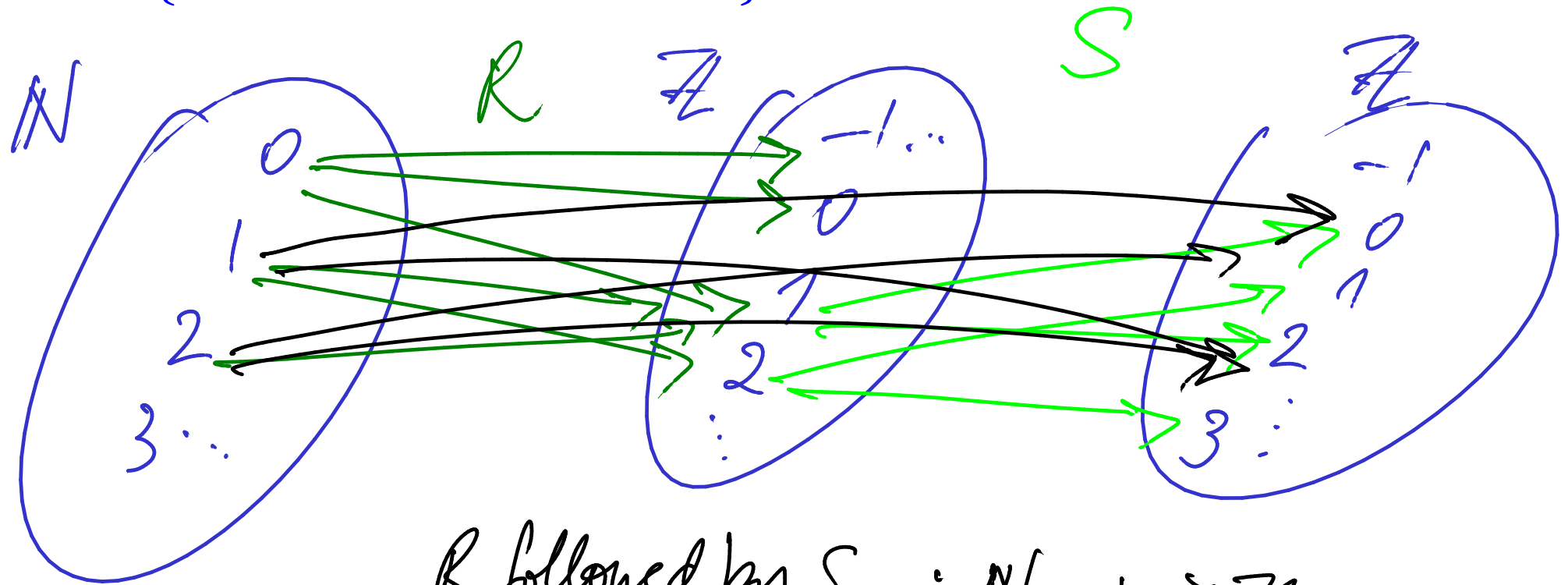
$$(m R_2 n \iff m = n^2)$$

# Internal diagrams

**Example:**

$$R = \{ (0, 0), (0, -1), (0, 1), (1, 2), (1, 1), (2, 1) \} : \mathbb{N} \dashrightarrow \mathbb{Z}$$

$$S = \{ (1, 0), (1, 2), (2, 1), (2, 3) \} : \mathbb{Z} \dashrightarrow \mathbb{Z}$$



$R$  followed by  $S : \mathbb{N} \dashrightarrow \mathbb{Z}$

## Relational composition

$$R: A \rightarrow B$$

$$S: B \rightarrow C$$

$$(S \circ R): A \rightarrow C$$

Def  $a \in A, c \in C$

$$a (S \circ R) c \stackrel{\text{def}}{\iff} \exists b \in B.$$

$$a R b \wedge b S c.$$

**Theorem 98** *Relational composition is associative and has the identity relation as neutral element.*

► *Associativity.*

For all  $R : A \twoheadrightarrow B$ ,  $S : B \twoheadrightarrow C$ , and  $T : C \twoheadrightarrow D$ ,

$$(T \circ S) \circ R = T \circ (S \circ R)$$

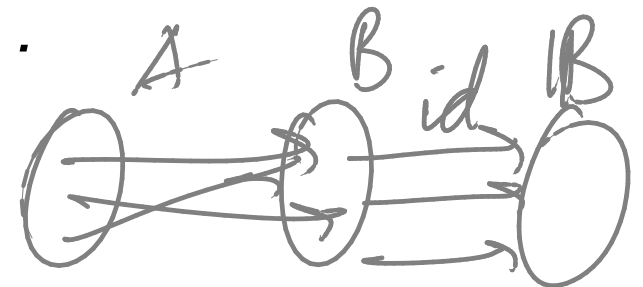
► *Neutral element.*

For all  $R : A \twoheadrightarrow B$ ,

$$R \circ \text{id}_A = R = \text{id}_B \circ R .$$

↳ justifies the notation

$T \circ S \circ R$



## Relational extensionality

$$R = S : A \rightarrow B$$

iff

$$\forall a \in A. \forall b \in B. a R b \iff a S b$$



$$\text{Def } [n] = \{0, \dots, n-1\} \quad \# [n] = n$$

Relations from  $[m]$  to  $[n]$  and  $(m \times n)$ -matrices over Booleans provide two alternative views of the same structure.

This carries over to identities and to composition/multiplication .

$$R : [m] \rightarrow [n] \rightsquigarrow \underline{\text{mat}}(R)$$

$$(\underline{\text{mat}}(R))_{i,j} = \text{true}$$

$$\iff \text{def } (i,j) \in R$$

$$\rightsquigarrow M$$

$$\underline{\text{rel}}(M)$$

$$(i,j) \in \text{rel}(M)$$

$$\iff \text{def } M_{i,j} = \underline{\text{true}}$$

