

## Important mathematical jargon: Sets

Very roughly, sets are the mathematicians' data structures. Informally, we will consider a set as a (well-defined, unordered) collection of mathematical objects, called the elements (or members) of the set.

## Set membership

The symbol ‘ $\in$ ’ known as the *set membership* predicate is central to the theory of sets, and its purpose is to build statements of the form

$$x \in A$$

that are true whenever it is the case that the object  $x$  is an element of the set  $A$ , and false otherwise.

## Defining sets

The set	of even primes	is	{2}
	of booleans		{true, false}
	[−2..3]		{−2, −1, 0, 1, 2, 3}

## Set comprehension

The basic idea behind set comprehension is to define a set by means of a property that precisely characterises all the elements of the set.

Notations:

$$\{x \in A \mid P(x)\} \quad , \quad \{x \in A : P(x)\}$$

# Greatest common divisor

Given a natural number  $n$ , the set of its *divisors* is defined by set comprehension as follows

$$D(n) = \{ d \in \mathbb{N} : d \mid n \} .$$

## Example 52

1.  $D(0) = \mathbb{N}$

2.  $D(1224) = \left\{ \begin{array}{l} 1, 2, 3, 4, 6, 8, 9, 12, 17, 18, 24, 34, 36, 51, 68, \\ 72, 102, 136, 153, 204, 306, 408, 612, 1224 \end{array} \right\}$

**Remark** Sets of divisors are hard to compute. However, the computation of the greatest divisor is straightforward. :)

Going a step further, what about the *common divisors* of pairs of natural numbers? That is, the set

$$\text{CD}(m, n) = \{ d \in \mathbb{N} : d \mid m \wedge d \mid n \}$$

for  $m, n \in \mathbb{N}$ .

### **Example 53**

$$\text{CD}(1224, 660) = \{ 1, 2, 3, 4, 6, 12 \}$$

Since  $\text{CD}(n, n) = D(n)$ , the computation of common divisors is as hard as that of divisors. But, what about the computation of the *greatest common divisor*?

**Lemma 55 (Key Lemma)** *Let  $m$  and  $m'$  be natural numbers and let  $n$  be a positive integer such that  $m \equiv m' \pmod{n}$ . Then,*

$$CD(m, n) = CD(m', n) .$$

PROOF:

**Lemma 57** For all positive integers  $m$  and  $n$ ,

$$\text{CD}(m, n) = \begin{cases} D(n) & , \text{ if } n \mid m \\ \text{CD}(n, \text{rem}(m, n)) & , \text{ otherwise} \end{cases}$$



**Lemma 57** For all positive integers  $m$  and  $n$ ,

$$\text{CD}(m, n) = \begin{cases} D(n) & , \text{ if } n \mid m \\ \text{CD}(n, \text{rem}(m, n)) & , \text{ otherwise} \end{cases}$$

Since a positive integer  $n$  is the greatest divisor in  $D(n)$ , the lemma suggests a recursive procedure:

$$\text{gcd}(m, n) = \begin{cases} n & , \text{ if } n \mid m \\ \text{gcd}(n, \text{rem}(m, n)) & , \text{ otherwise} \end{cases}$$

for computing the *greatest common divisor*, of two positive integers  $m$  and  $n$ . This is

## Euclid's Algorithm

gcd

```
fun gcd( m , n )  
  = let  
    val ( q , r ) = divalg( m , n )  
  in  
    if r = 0 then n  
    else gcd( n , r )  
  end
```

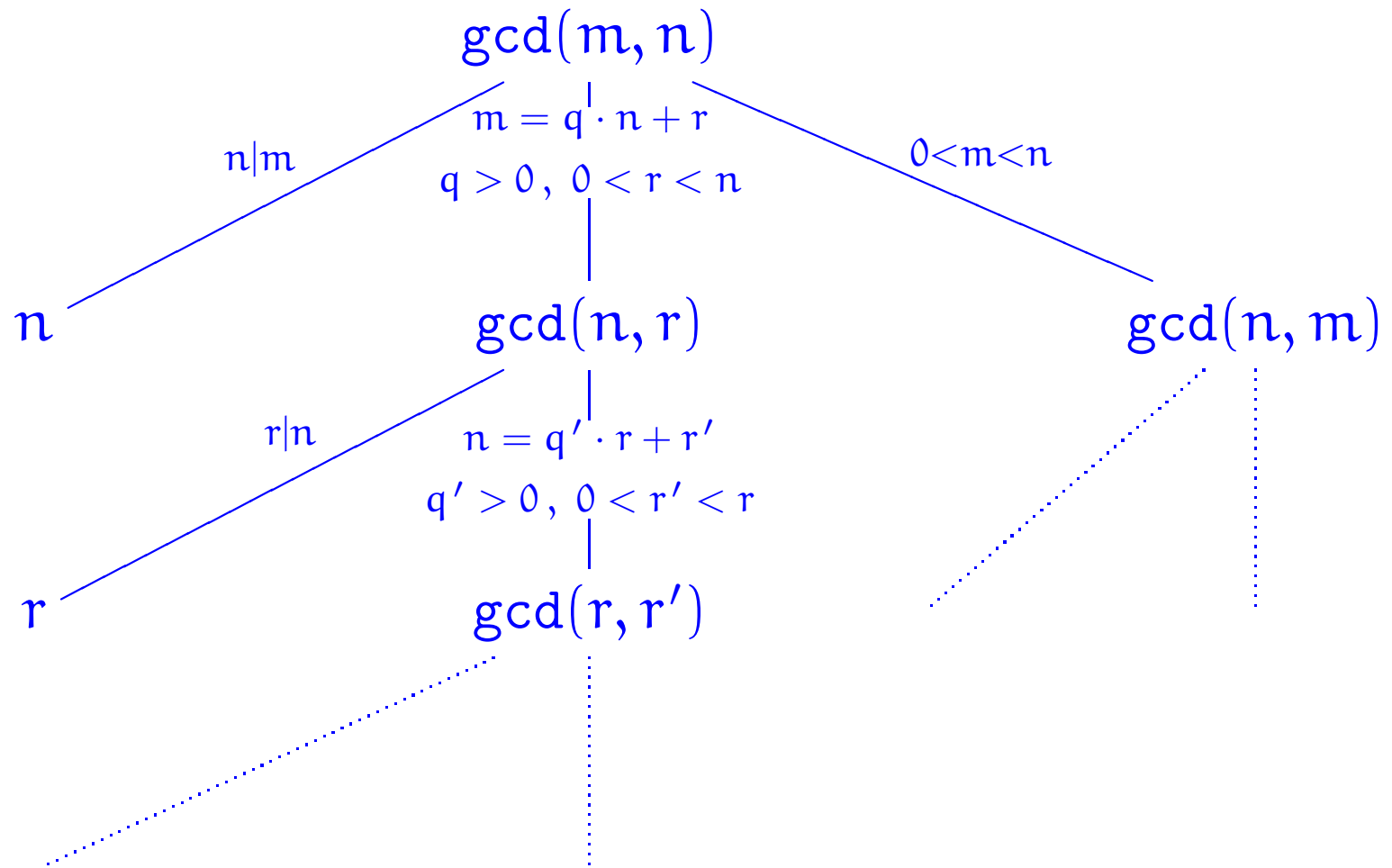
**Example 58** ( $\gcd(13, 34) = 1$ )

$$\begin{aligned}\gcd(13, 34) &= \gcd(34, 13) \\ &= \gcd(13, 8) \\ &= \gcd(8, 5) \\ &= \gcd(5, 3) \\ &= \gcd(3, 2) \\ &= \gcd(2, 1) \\ &= 1\end{aligned}$$

**Theorem 59** *Euclid's Algorithm  $\gcd$  terminates on all pairs of positive integers and, for such  $m$  and  $n$ ,  $\gcd(m, n)$  is the greatest common divisor of  $m$  and  $n$  in the sense that the following two properties hold:*

- (i) *both  $\gcd(m, n) \mid m$  and  $\gcd(m, n) \mid n$ , and*
- (ii) *for all positive integers  $d$  such that  $d \mid m$  and  $d \mid n$  it necessarily follows that  $d \mid \gcd(m, n)$ .*

PROOF:



## Fractions in lowest terms

```
fun lowterms( m , n )  
  = let  
    val gcdval = gcd( m , n )  
  in  
    ( m div gcdval , n div gcdval )  
  end
```

# Some fundamental properties of gcds

**Lemma 61** For all positive integers  $l$ ,  $m$ , and  $n$ ,

1. **(Commutativity)**  $\gcd(m, n) = \gcd(n, m)$ ,
2. **(Associativity)**  $\gcd(l, \gcd(m, n)) = \gcd(\gcd(l, m), n)$ ,
3. **(Linearity)<sup>a</sup>**  $\gcd(l \cdot m, l \cdot n) = l \cdot \gcd(m, n)$ .

PROOF:

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<sup>a</sup>Aka (Distributivity).