# **Datacenters** (Optional fun)

#### What we will cover

(Datacenter Topic 7 is not examinable in 2013-14)

- Characteristics of a datacenter environment
  - goals, constraints, workloads, etc.
- How and why DC networks are different (vs. WAN)
  - e.g., latency, geo, autonomy, ...
- How traditional solutions fare in this environment
  - e.g., IP, Ethernet, TCP, ARP, DHCP
- Not details of how datacenter networks operate

#### Disclaimer

Material is emerging (not established) wisdom

- Material is incomplete
  - many details on how and why datacenter networks operate aren't public

# Why Datacenters?

Your <public-life, private-life, banks, government> live in my datacenter.

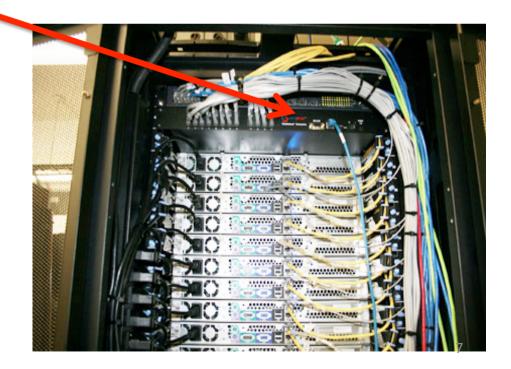
Security, Privacy, Control, Cost, Energy, (breaking) received wisdom; all this and more come together into sharp focus in datacenters.

Do I need to labor the point?

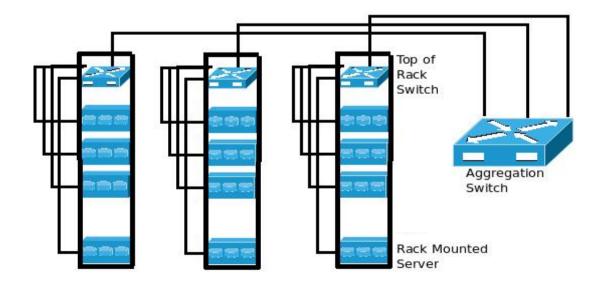
Servers organized in racks



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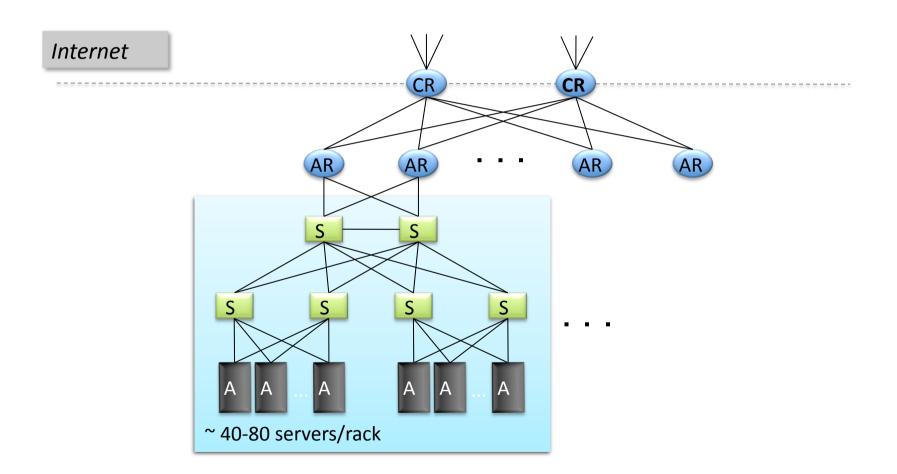


- Servers organized in racks
- Each rack has a `Top of Rack' (ToR) switch
- An `aggregation fabric' interconnects ToR switches
- Connected to the outside via `core' switches
  - note: blurry line between aggregation and core
- With network redundancy of ~2x for robustness

Example 1 Campus Network Brocade MLXe Brocade JCX 6650 Brocade Brocade ICX 6650 MLXe Brocade\* ICX 6650 Brocade ICX 6650 Data Center Brocade ICX 6650 Aggregation/Core Brocade ICX 6650 10 GbE Servers 10 GbE Data Center 40 GbE Access

Brocade reference design

# Example 2



Cisco reference design

#### Observations on DC architecture

- Regular, well-defined arrangement
- Hierarchical structure with rack/aggr/core layers
- Mostly homogenous within a layer
- Supports communication between servers and between servers and the external world

Contrast: ad-hoc structure, heterogeneity of WANs

#### Datacenters have been around for a while





# SCALE!



# How big exactly?

- 1M servers [Microsoft]
  - less than google, more than amazon

> \$1B to build one site [Facebook]

• >\$20M/month/site operational costs [Microsoft '09]

But only O(10-100) sites

#### What's new?

- Scale
- Service model
  - user-facing, revenue generating services
  - multi-tenancy
  - jargon: SaaS, PaaS, DaaS, laaS, ...

# **Implications**

- Scale
  - need scalable solutions (duh)
  - improving efficiency, lowering cost is critical
  - → `scale out' solutions w/ commodity technologies
- Service model
  - performance means \$\$
  - virtualization for isolation and portability

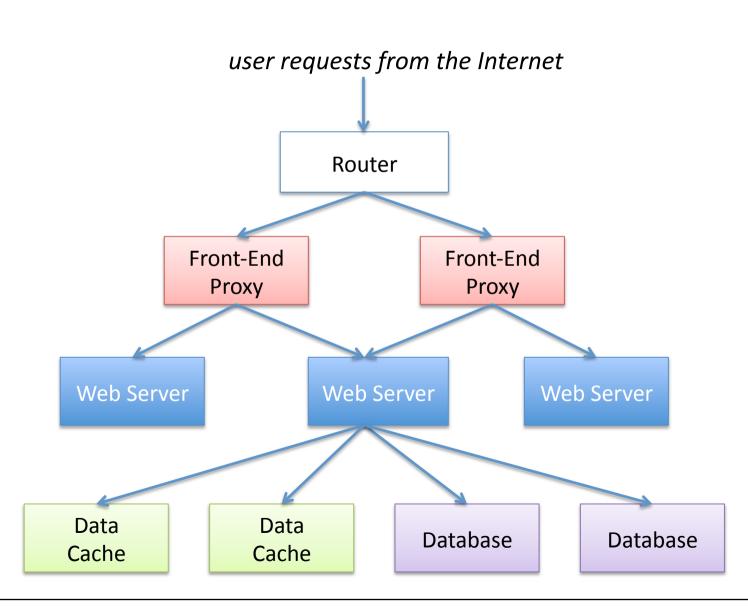
# Multi-Tier Applications

- Applications decomposed into tasks
  - Many separate components
  - Running in parallel on different machines

# Componentization leads to different types of network traffic

- "North-South traffic"
  - Traffic between external clients and the datacenter
  - Handled by front-end (web) servers, mid-tier application servers, and back-end databases
  - Traffic patterns fairly stable, though diurnal variations

# North-South Traffic



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# Componentization leads to different types of network traffic

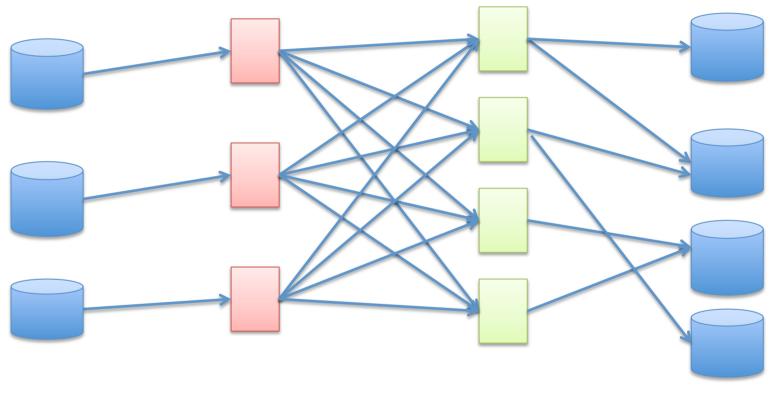
#### "North-South traffic"

- Traffic between external clients and the datacenter
- Handled by front-end (web) servers, mid-tier application servers, and back-end databases
- Traffic patterns fairly stable, though diurnal variations

#### "East-West traffic"

- Traffic between machines in the datacenter
- Comm within "big data" computations (e.g. Map Reduce)
- Traffic may shift on small timescales (e.g., minutes)

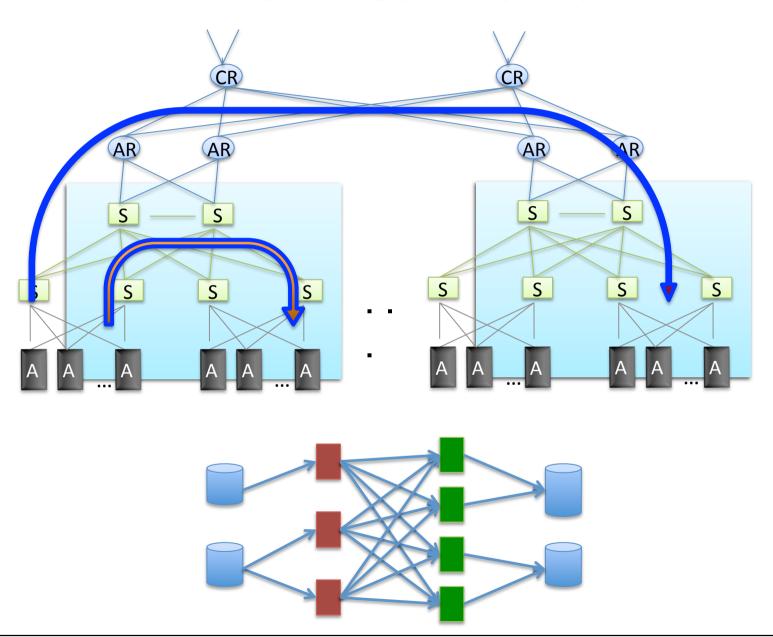
#### **East-West Traffic**



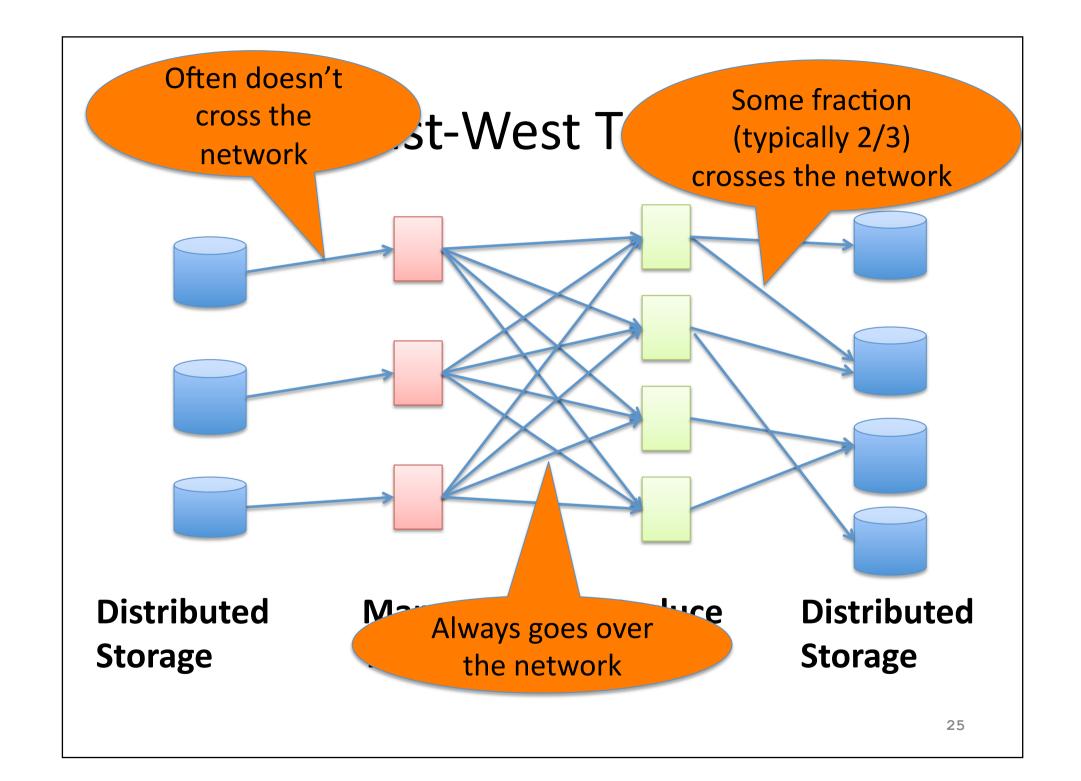
Distributed Storage

Map Tasks Reduce Tasks Distributed Storage

# **East-West Traffic**



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- Huge scale:
  - ~20,000 switches/routers
  - contrast: AT&T ~500 routers

- Huge scale:
- Limited geographic scope:
  - High bandwidth: 10/40/100G
  - Contrast: Cable/aDSL/WiFi
  - Very low RTT: 10s of microseconds
  - Contrast: 100s of milliseconds in the WAN

- Huge scale
- Limited geographic scope
- Single administrative domain
  - Can deviate from standards, invent your own, etc.
  - "Green field" deployment is still feasible

- Huge scale
- Limited geographic scope
- Single administrative domain
- Control over one/both endpoints
  - can change (say) addressing, congestion control, etc.
  - can add mechanisms for security/policy/etc. at the endpoints (typically in the hypervisor)

- Huge scale
- Limited geographic scope
- Single administrative domain
- Control over one/both endpoints
- Control over the placement of traffic source/sink
  - e.g., map-reduce scheduler chooses where tasks run
  - alters traffic pattern (what traffic crosses which links)

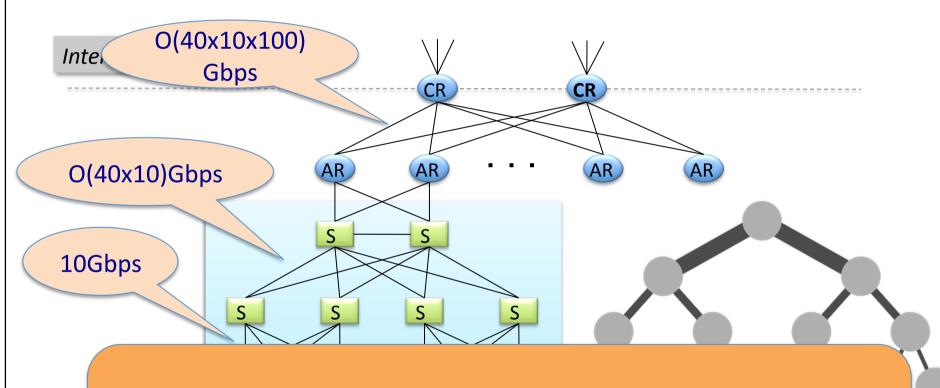
- Huge scale
- Limited geographic scope
- Single administrative domain
- Control over one/both endpoints
- Control over the placement of traffic source/sink
- Regular/planned topologies (e.g., trees/fat-trees)
  - Contrast: ad-hoc WAN topologies (dictated by real-world geography and facilities)

- Huge scale
- Limited geographic scope
- Single administrative domain
- Control over one/both endpoints
- Control over the placement of traffic source/sink
- Regular/planned topologies (e.g., trees/fat-trees)
- Limited heterogeneity
  - link speeds, technologies, latencies, ...

#### <u>Goals</u>

- Extreme bisection bandwidth requirements
  - recall: all that east-west traffic
  - target: any server can communicate at its full link speed
  - problem: server's access link is 10Gbps!

#### **Full Bisection Bandwidth**



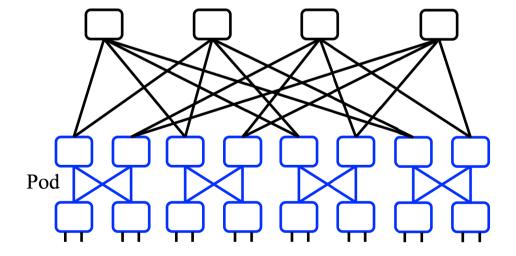
#### Traditional tree topologies "scale up"

- full bisection bandwidth is expensive
- typically, tree topologies "oversubscribed"

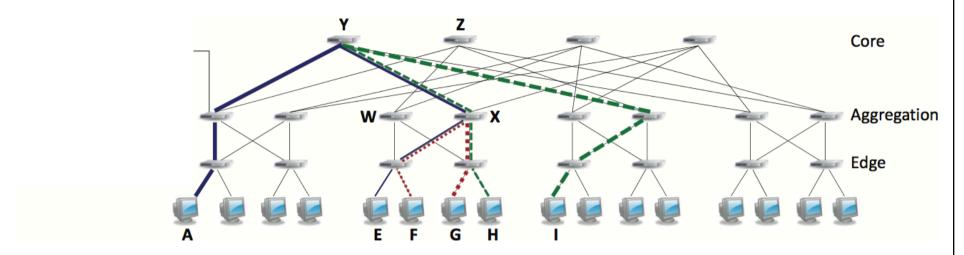
# A "Scale Out" Design

- Build multi-stage `Fat Trees' out of k-port switches
  - k/2 ports up, k/2 down
  - Supports  $k^3/4$  hosts:
    - 48 ports, 27,648 hosts

All links are the same speed (e.g. 10Gps)



#### Full Bisection Bandwidth Not Sufficient



- To realize full bisectional throughput, routing must spread traffic across paths
- Enter load-balanced routing
  - How? (1) Let the network split traffic/flows at random (e.g., ECMP protocol -- RFC 2991/2992)
  - How? (2) Centralized flow scheduling?
  - Many more research proposals

#### <u>Goals</u>

- Extreme bisection bandwidth requirements
- Extreme latency requirements
  - real money on the line
  - current target: 1μs RTTs
  - how? cut-through switches making a comeback
    - reduces switching time

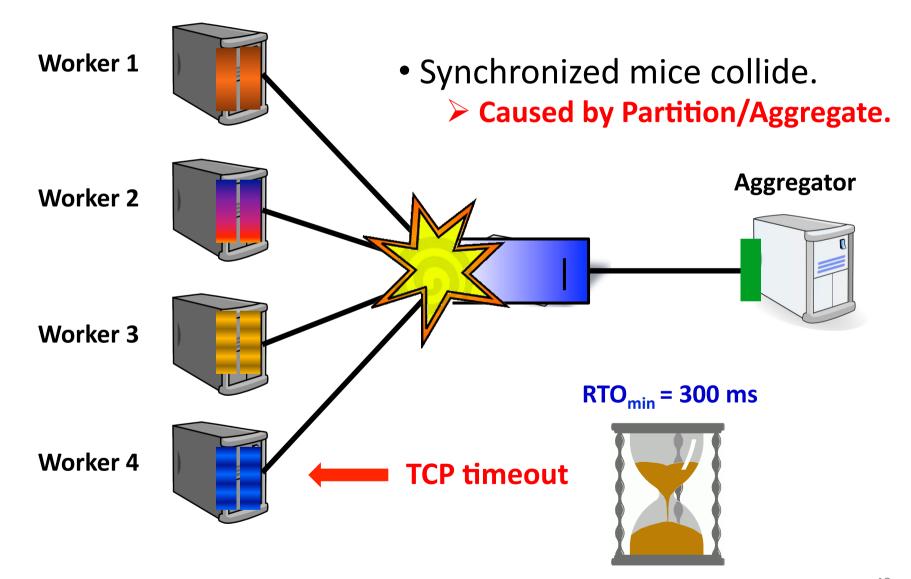
### <u>Goals</u>

- Extreme bisection bandwidth requirements
- Extreme latency requirements
  - real money on the line
  - current target: 1μs RTTs
  - how? cut-through switches making a comeback
  - how? avoid congestion
    - reduces queuing delay

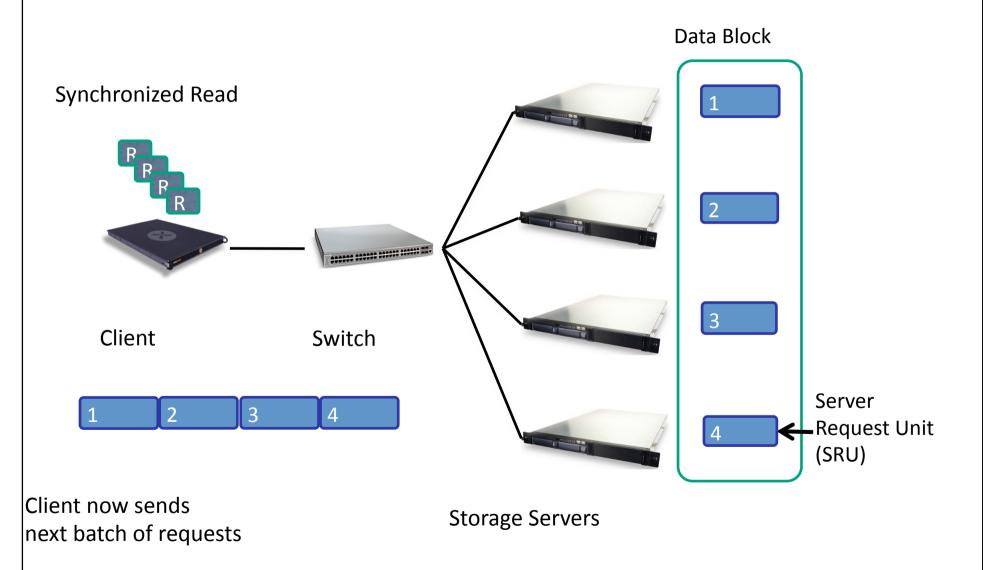
### Goals

- Extreme bisection bandwidth requirements
- Extreme latency requirements
  - real money on the line
  - current target: 1μs RTTs
  - how? cut-through switches making a comeback (lec. 2!)
  - how? avoid congestion
  - how? fix TCP timers (e.g., default timeout is 500ms!)
  - how? fix/replace TCP to more rapidly fill the pipe

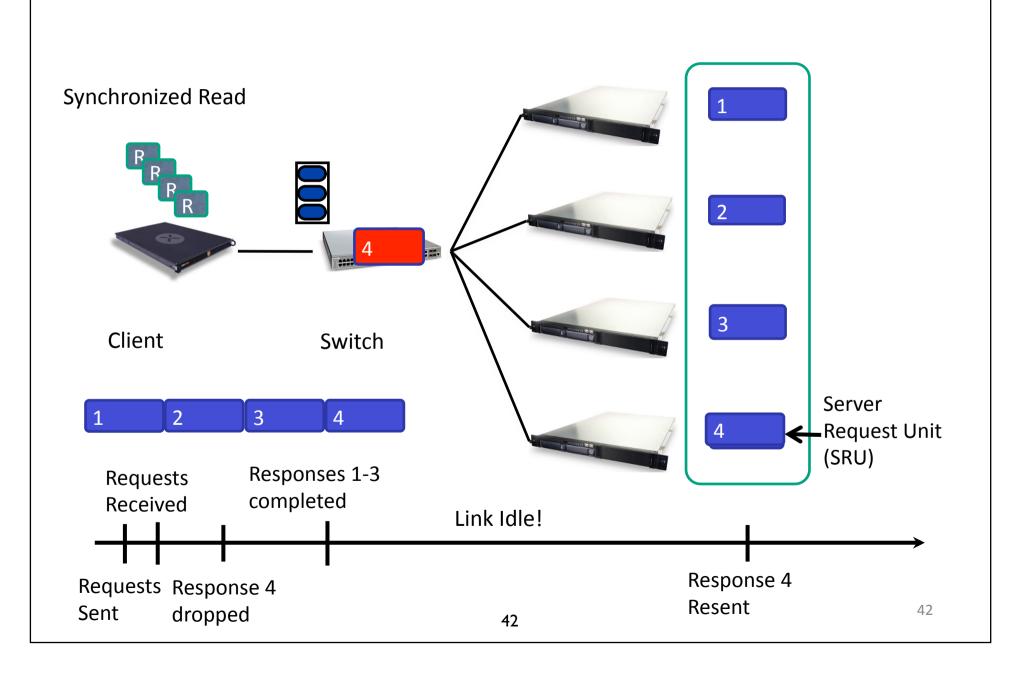
## An example problem at scale - INCAST



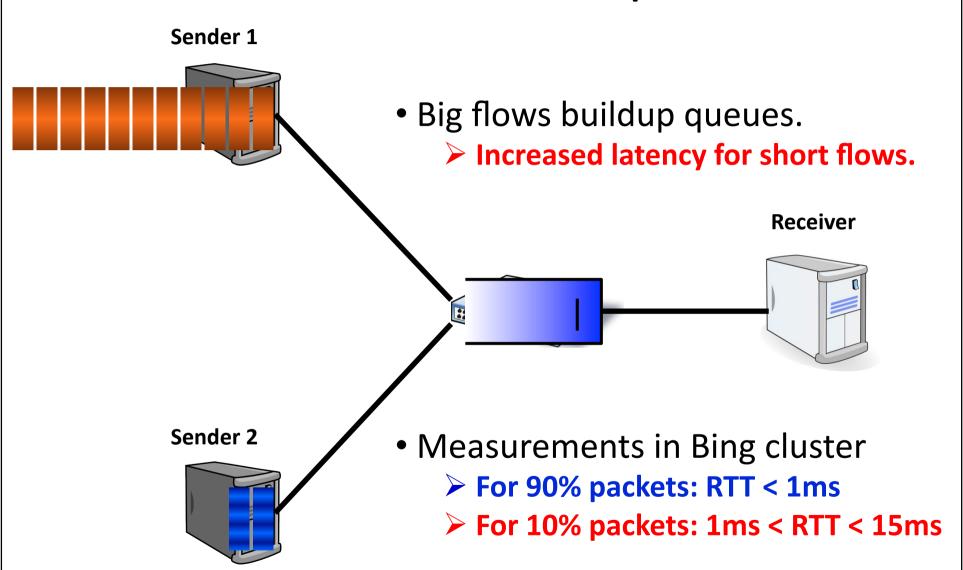
## The Incast Workload



## **Incast Workload Overfills Buffers**



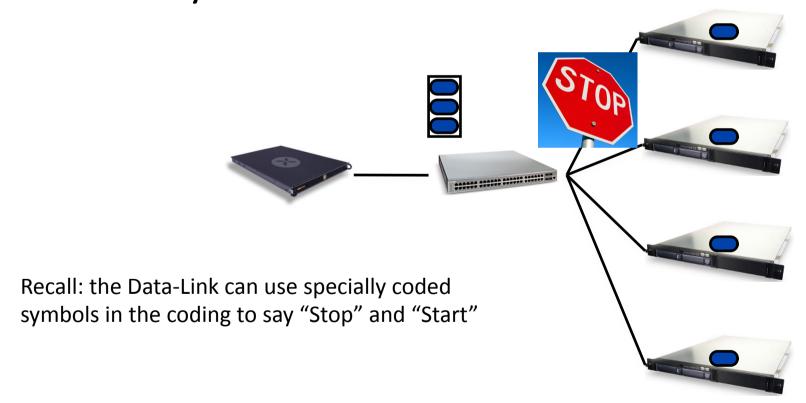
## Queue Buildup



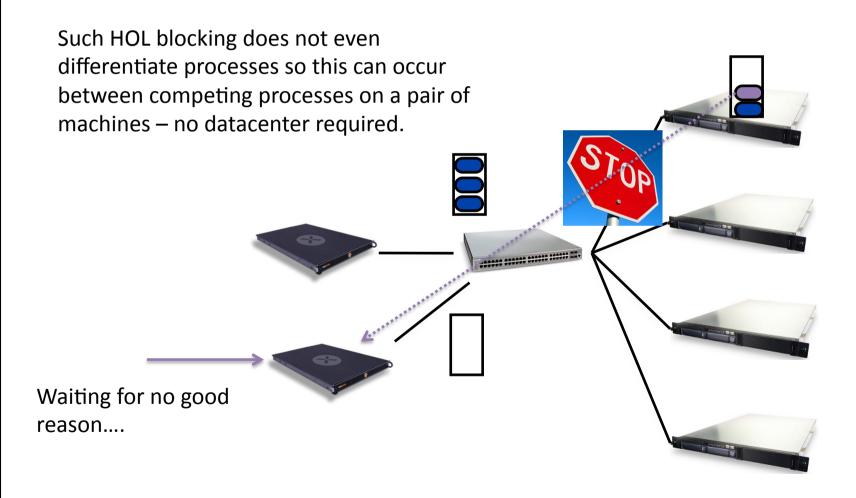
## Link-Layer Flow Control

Common between switches but this is flow-control to the end host too...

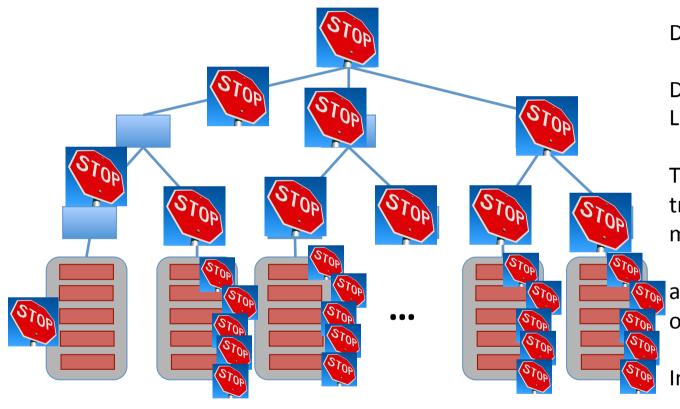
 Another idea to reduce incast is to employ Link-Layer Flow Control.....



## Link Layer Flow Control – The Dark side Head of Line Blocking....



# Link Layer Flow Control But its worse that you imagine....



Double down on trouble....

Did I mention this is Link-Layer!

That means no (IP) control traffic, no routing messages....

a whole system waiting for one machine

Incast is very unpleasant.

Reducing the impact of HOL in Link Layer Flow Control can be done through priority queues and *overtaking*....

### Goals

- Extreme bisection bandwidth requirements
- Extreme latency requirements
- Predictable, deterministic performance
  - "your packet will reach in Xms, or not at all"
  - "your VM will always see at least YGbps throughput"
  - Resurrecting `best effort' vs. `Quality of Service' debates
  - How is still an open question

#### Goals

- Extreme bisection bandwidth requirements
- Extreme latency requirements
- Predictable, deterministic performance
- Differentiating between tenants is key
  - e.g., "No traffic between VMs of tenant A and tenant B"
  - "Tenant X cannot consume more than XGbps"
  - "Tenant Y's traffic is low priority"

### <u>Goals</u>

- Extreme bisection bandwidth requirements
- Extreme latency requirements
- Predictable, deterministic performance
- Differentiating between tenants is key
- Scalability (of course)
  - Q: How's Ethernet spanning tree looking?

### <u>Goals</u>

- Extreme bisection bandwidth requirements
- Extreme latency requirements
- Predictable, deterministic performance
- Differentiating between tenants is key
- Scalability (of course)
- Cost/efficiency
  - focus on commodity solutions, ease of management
  - some debate over the importance in the network case

## Summary

- new characteristics and goals
- some liberating, some constraining
- scalability is the baseline requirement
- more emphasis on performance
- less emphasis on heterogeneity
- less emphasis on interoperability