

Lecture 3: category of nominal sets

Category of nominal sets, **Nom**

- ▶ objects $X \in \mathbf{Nom}$ are nominal sets
- ▶ morphisms $f \in \mathbf{Nom}(X, Y)$ are functions $f \in Y^X$ that are **equivariant**:

$$(\forall \pi \in \mathbf{Perm} \mathbb{A}, x \in X) \pi \cdot (f x) = f(\pi \cdot x)$$

for all $\pi \in \mathbf{Perm} \mathbb{A}$, $x \in X$.

(Clearly these are closed under composition and include identity functions—so we do get a category.)

E.g. $\mathbf{Nom}(\mathbb{A}, \mathbb{A}) = ?$

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(Clearly these are closed under composition and include identity functions—so we do get a category.)

E.g. $\mathbf{Nom}(\mathbb{A}, \mathbb{A}) = \{id_{\mathbb{A}}\}$ because...

Lemma 1. If $f \in \mathbf{Nom}(X, Y)$ and $x \in X$, then

$$\mathit{supp}(f x) \subseteq \mathit{supp} x$$

For example, if $f \in \mathbf{Nom}(\mathbb{A}, \mathbb{A})$, then for any $a \in \mathbb{A}$

$$\{f a\} = \mathit{supp}(f a) \subseteq \mathit{supp} a = \{a\}$$

so that $f a = a$. Hence $f = id_{\mathbb{A}}$.

Lemma 1. If $f \in \text{Nom}(X, Y)$ and $x \in X$, then

$$\text{supp}(f x) \subseteq \text{supp } x$$

Proof. Suppose A supports x in X .

So for any $\pi \in \text{Perm } A$, if $(\forall a \in A) \pi a = a$, then $\pi \cdot x = x$ and hence $\pi \cdot (f x) = f(\pi \cdot x) = f x$.

Hence A also supports $f x$ in Y .

Taking $A = \text{supp } x$, we get that $\text{supp } x$ supports $f x$, so $\text{supp}(f x)$ is contained in $\text{supp } x$. \square

Nom is a cartesian closed category

Finite products: $X_1 \times \cdots \times X_n$ is given by cartesian product of sets with **Perm** \mathbb{A} -action

$$\pi \cdot (x_1, \dots, x_n) \triangleq (\pi \cdot x_1, \dots, \pi \cdot x_n)$$

which satisfies

$$\text{supp}(x, \dots, x_n) = (\text{supp } x_1) \cup \cdots \cup (\text{supp } x_n)$$

(Exercise)

Nom is a cartesian closed category

Exponentials: given $X, Y \in \mathbf{Nom}$, we get a **Perm** \mathbb{A} -action on the set Y^X of functions:

$$\pi \cdot f \triangleq \lambda(x \in X) \rightarrow \pi \cdot (f(\pi^{-1} \cdot x))$$

Why this definition? Whatever $\pi \cdot f$ is, want
application to be equivariant

$$\pi \cdot (f x) = (\pi \cdot f)(\pi \cdot x)$$

So we are forced to get

$$(\pi \cdot f) x = (\pi \cdot f)(\pi \cdot (\pi^{-1} \cdot x)) = \pi \cdot (f(\pi^{-1} \cdot x))$$

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Not every $f \in Y^X$ need have finite support wrt this action: let $X \rightarrow_{\text{fs}} Y$ be the subset of ones that do.

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E.g. given $a \in \mathbb{A}$, $K_a \triangleq \lambda(x \in \mathbb{A}) \rightarrow a$ is in $\mathbb{A} \rightarrow_{\text{fs}} \mathbb{A}$, because it is supported by $\{a\}$, since

$$(\pi \cdot K_a) x = \pi \cdot (K_a(\pi^{-1} \cdot x)) = \pi \cdot a = \pi a$$

and hence $\pi \cdot K_a = K_{\pi a}$.

Nom is a cartesian closed category

Exponential of $X, Y \in \mathbf{Nom}$ is $X \rightarrow_{\text{fs}} Y$ plus equivariant function $\mathbf{app} : (X \rightarrow_{\text{fs}} Y) \times X \rightarrow Y$
 $\mathbf{app}(f, x) = f x$

Given $f \in \mathbf{Nom}(Z \times X, Y)$, the unique $\hat{f} \in \mathbf{Nom}(Z, X \rightarrow_{\text{fs}} Y)$ making

$$\begin{array}{ccc} Z \times X & & \\ \hat{f} \times id_X \downarrow & \searrow f & \\ (X \rightarrow_{\text{fs}} Y) \times X & \xrightarrow{\mathbf{app}} & Y \end{array}$$

commute is given by currying: $\hat{f} z = \lambda(x \in X) \rightarrow f(z, x)$.

(Exercise)

Nom is a model of Church's higher order logic

[**Nom** is categorically equivalent to a well-known
Boolean topos, called the **Schanuel topos**.]

Nom is a model of Church's higher order logic

Subobject classifier: $\Omega = \{\text{true}, \text{false}\} \cong 1 + 1$
with trivial **Perm** \mathbb{A} -action: $\pi \cdot b \triangleq b$ (so $\text{supp } b = \emptyset$).

Power objects: $X \rightarrow_{\text{fs}} \Omega \cong \mathbf{P}_{\text{fs}} X$, the set of subsets $S \subseteq X$ that are finitely supported w.r.t. the **Perm** \mathbb{A} -action

$$\pi \cdot S \triangleq \{\pi \cdot x \mid x \in S\}$$

Nom is a model of Church's higher order logic

Coproducts are given by disjoint union.

Natural number object: $\mathbb{N} = \{0, 1, 2, \dots\}$ with trivial **Perm** \mathbb{A} -action: $\pi \cdot n \triangleq n$ (so $\text{supp } n = \emptyset$).

Nom \neq choice

Nom models classical higher-order logic, but not Hilbert's ε -operation, $\varepsilon x. \varphi(x)$ satisfying

$$(\forall x : X) \varphi(x) \Rightarrow \varphi(\varepsilon x. \varphi(x))$$

Theorem 1. There is no equivariant function $c : \{S \in \mathbf{P}_{\text{fs}} \mathbb{A} \mid S \neq \emptyset\} \rightarrow \mathbb{A}$ satisfying $c(S) \in S$ for all non-empty $S \in \mathbf{P}_{\text{fs}} \mathbb{A}$.

Proof. Suppose there were such a c . Putting $a \triangleq c \mathbb{A}$ and picking some $b \in \mathbb{A} - \{a\}$, we get a contradiction to $a \neq b$:

$$a = c \mathbb{A} = c((a \ b) \cdot \mathbb{A}) = (a \ b) \cdot c \mathbb{A} = (a \ b) \cdot a = b$$

Nom $\not\models$ choice

Nom models classical higher-order logic, but not Hilbert's ε -operation, $\varepsilon x. \varphi(x)$ satisfying

$$(\forall x : X) \varphi(x) \Rightarrow \varphi(\varepsilon x. \varphi(x))$$

In fact **Nom** does not model even very weak forms of choice, such as Dependent Choice.

The nominal set of names

Recall that \mathbb{A} is a nominal set once equipped with the action

$$\pi \cdot a = \pi(a)$$

which satisfies $\text{supp } a = \{a\}$.

Although $\mathbb{A} \in \mathbf{Set}$ is a countable, \mathbb{A} is not isomorphic to \mathbb{N} in \mathbf{Nom} . For any $f \in \mathbb{N} \rightarrow_{\text{fs}} \mathbb{A}$ has to satisfy

$$\{f\ n\} = \text{supp}(\text{app}(f, n)) \subseteq \\ \text{supp}(f, n) = \text{supp } f \cup \text{supp } n = \text{supp } f$$

for all $n \in \mathbb{N}$, and so f cannot be surjective.