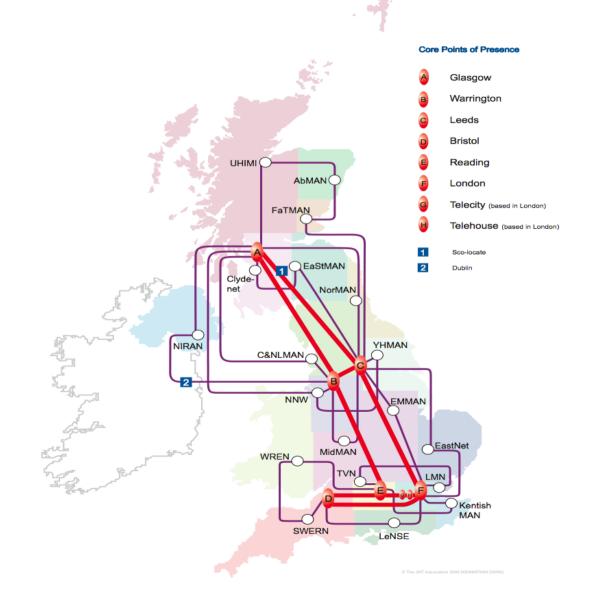
## L11: BGP Lecture 14 2013

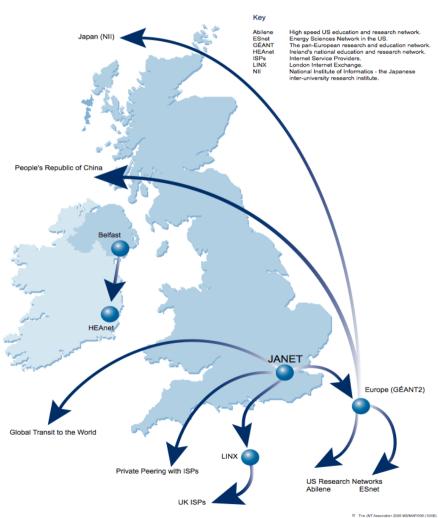
Timothy G. Griffin Computer Lab Cambridge UK

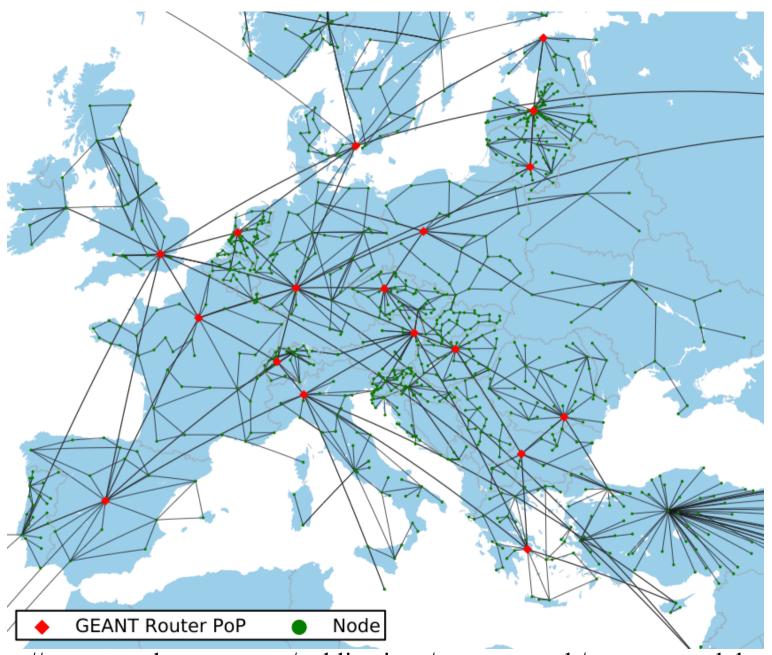
#### **JANET**



#### **JANET** and the Internet

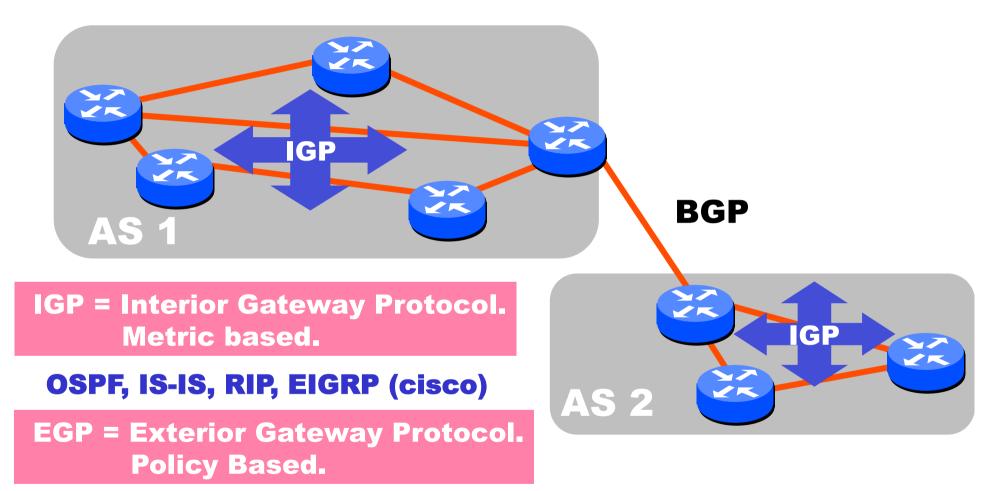
#### **JANET External Network Access Provision**





http://www.topology-zoo.org/publications/eu\_nren\_tech/eu\_nren\_tech.html

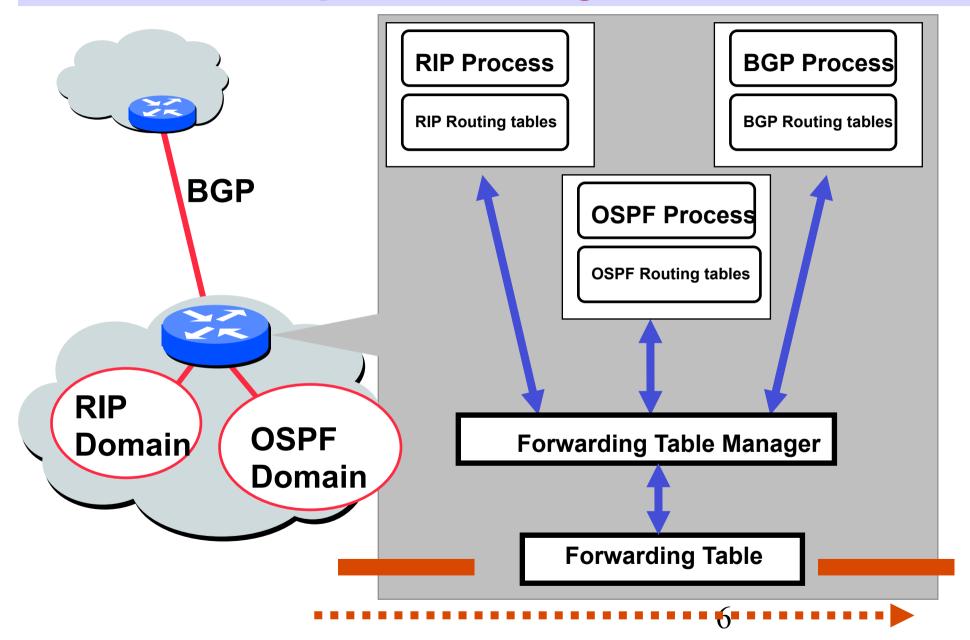
# Architecture of Dynamic Routing



**Only one: BGP** 

The Routing Domain of BGP is the entire Internet

## Happy Packets: The Internet Does Not Exist Only to Populated Routing Tables



## **Autonomous Routing Domains**

A collection of physical networks glued together using IP, that have a unified administrative routing policy.

- Campus networks
- Corporate networks
- ISP Internal networks

• ...

## **Autonomous Systems (ASes)**

An autonomous system is an autonomous routing domain that has been assigned an Autonomous System Number (ASN).

... the administration of an AS appears to other ASes to have a single coherent interior routing plan and presents a consistent picture of what networks are reachable through it.

RFC 1930: Guidelines for creation, selection, and registration of an Autonomous System

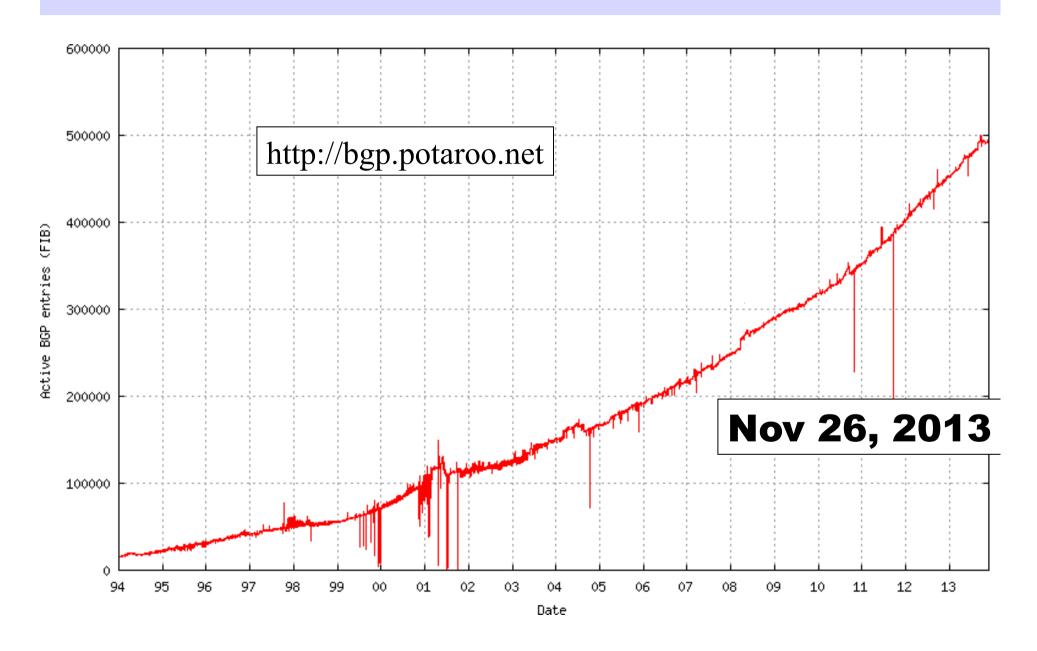
## **AS Numbers (ASNs)**

```
    JANET: 786
```

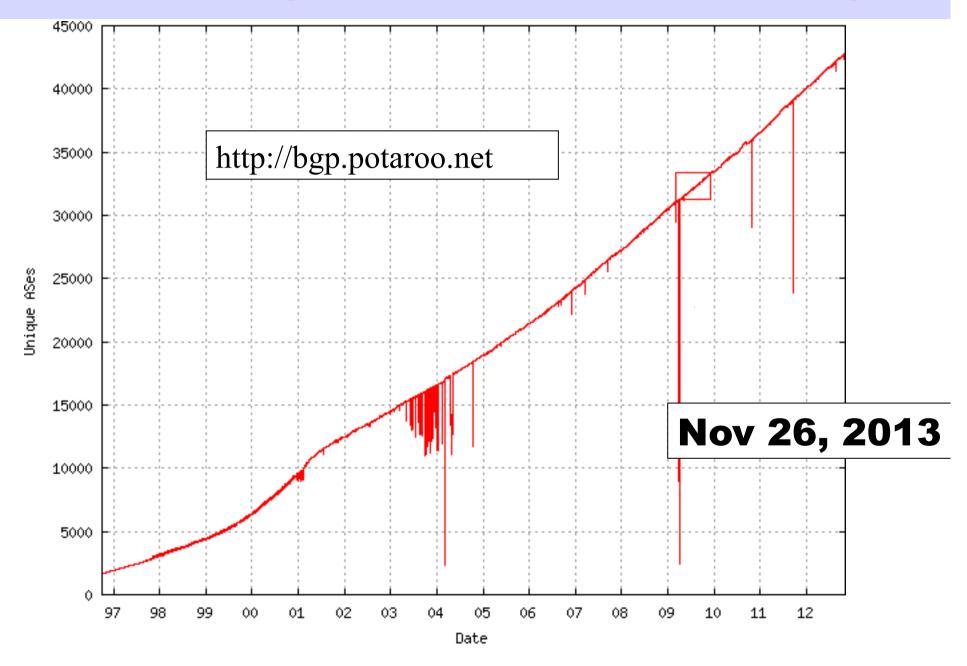
- MIT: 3
- Harvard: 11
- UC San Diego: 7377
- AT&T: 7018, 6341, 5074, ...
- UUNET: 701, 702, 284, 12199, ...
- Sprint: 1239, 1240, 6211, 6242, ...
- ...

**ASNs** represent units of routing policy

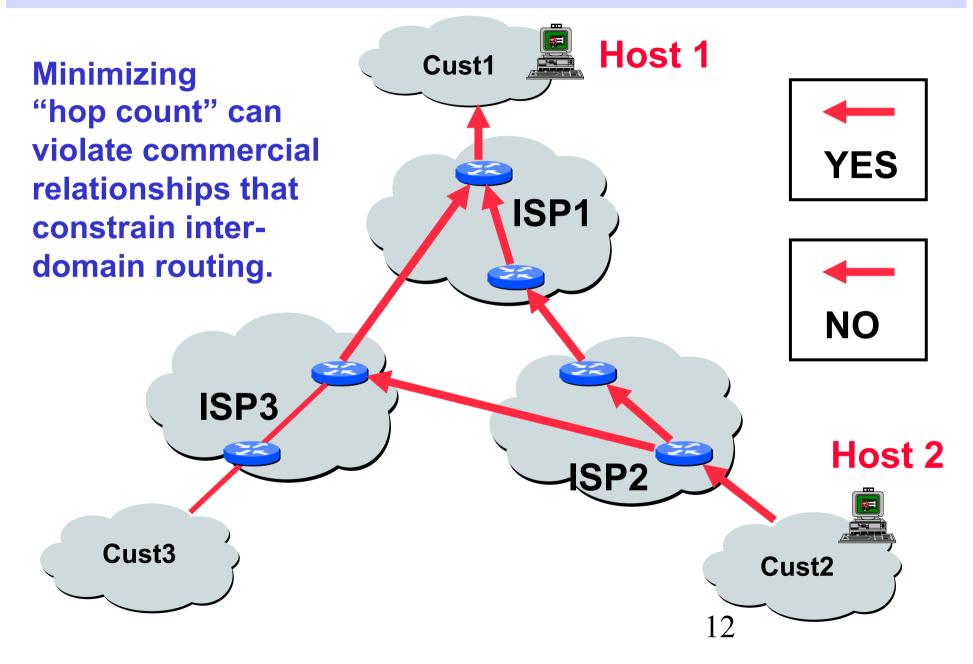
### How many prefixes are used today?



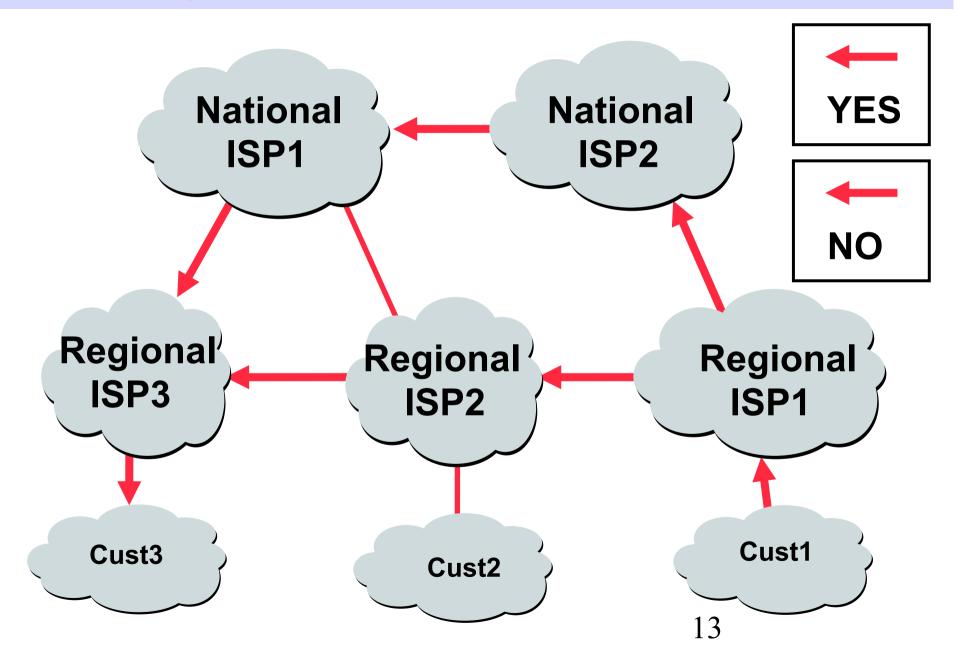
### How many ASNs are used today?



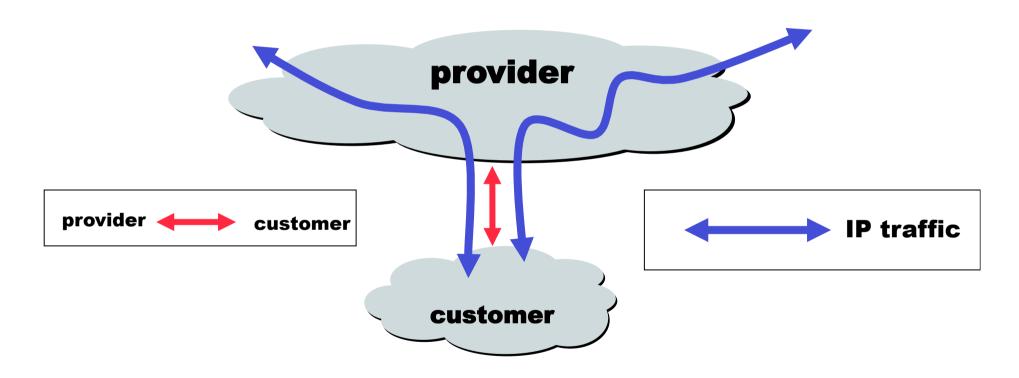
#### Policy-Based vs. Distance-Based Routing?



#### Why not minimize "AS hop count"?

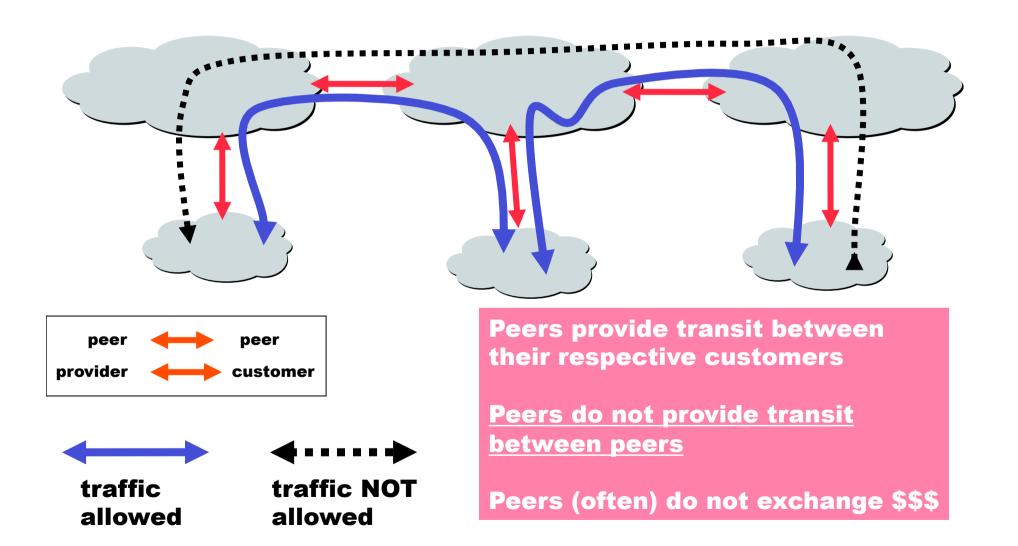


#### **Customers and Providers**

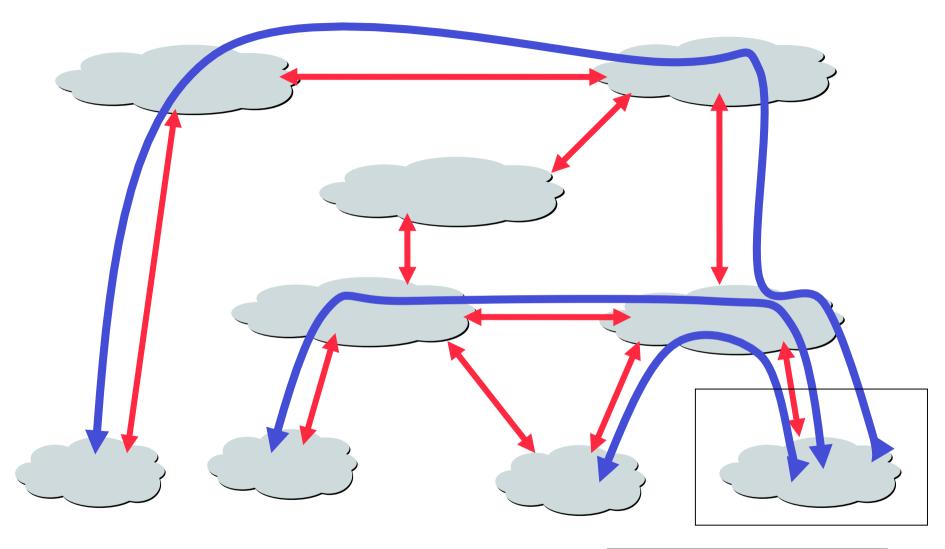


**Customer pays provider for access to the Internet** 

## The "Peering" Relationship



## **Peering Provides Shortcuts**



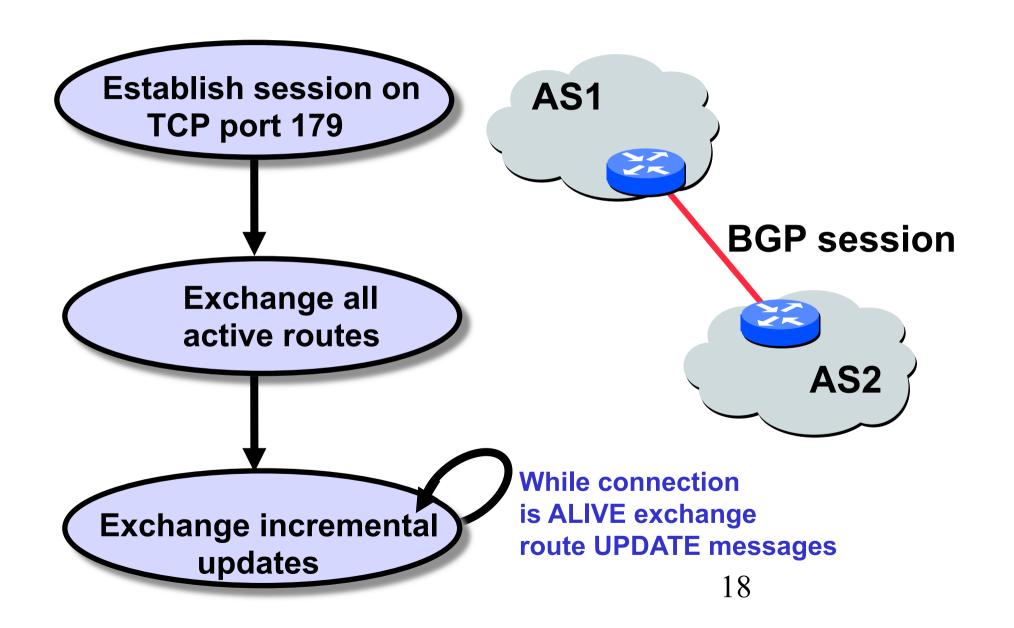
Peering also allows connectivity between the customers of "Tier 1" providers.



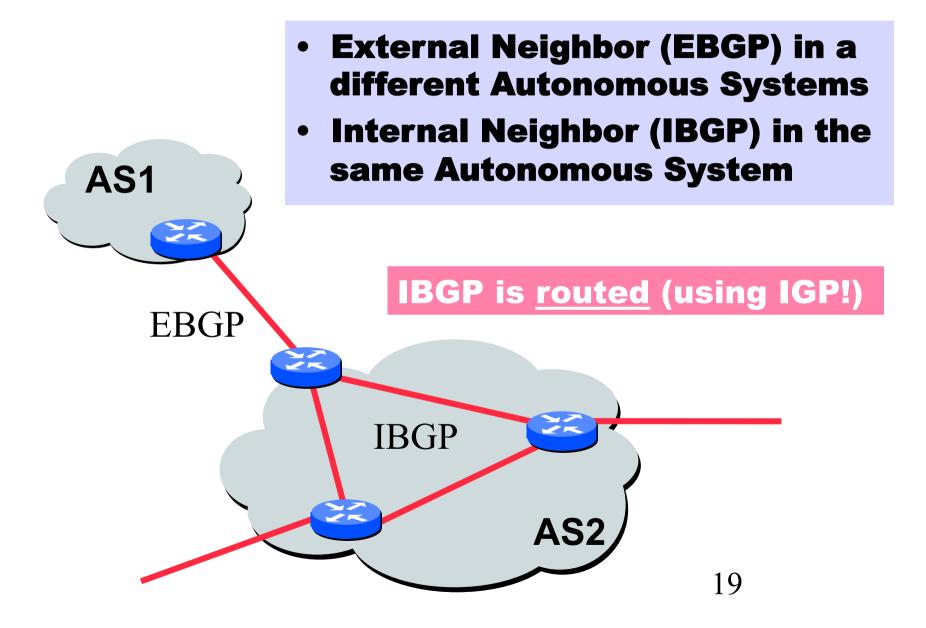
## BGP-4

- **BGP** = **B**order **G**ateway **P**rotocol
- Is a <u>Policy-Based</u> routing protocol
- Is the <u>de facto EGP</u> of today's global Internet
- Relatively simple protocol, but configuration is complex and the entire world can see, and be impacted by, your mistakes.
  - 1989 : BGP-1 [RFC 1105]
    - Replacement for EGP (1984, RFC 904)
  - 1990 : BGP-2 [RFC 1163]
  - 1991 : BGP-3 [RFC 1267]
  - 1995 : BGP-4 [RFC 1771]
    - Support for Classless Interdomain Routing (CIDR)
  - 2006 : BGP-4 [RFC 4271]

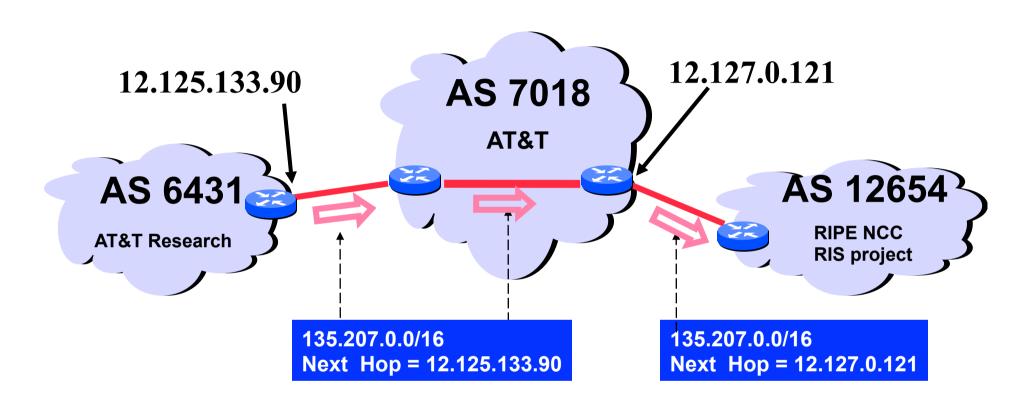
## **BGP Operations (Simplified)**



#### **Two Types of BGP Sessions**

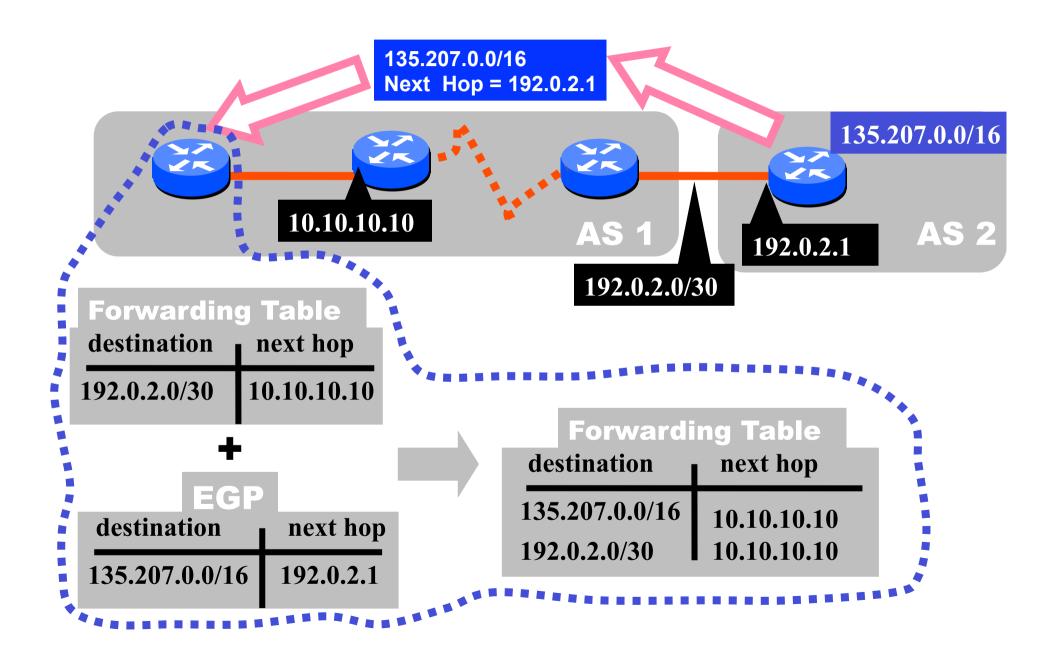


## **BGP Next Hop Attribute**



Every time a route announcement crosses an AS boundary, the Next Hop attribute is changed to the IP address of the border router that announced the route.

#### Join EGP with IGP For Connectivity



## Four Types of BGP Messages

- Open: Establish a peering session.
- **Keep Alive**: Handshake at regular intervals.
- Notification : Shuts down a peering session.
- **Update**: <u>Announcing</u> new routes or <u>withdrawing</u> previously announced routes.

announcement = prefix + attributes values

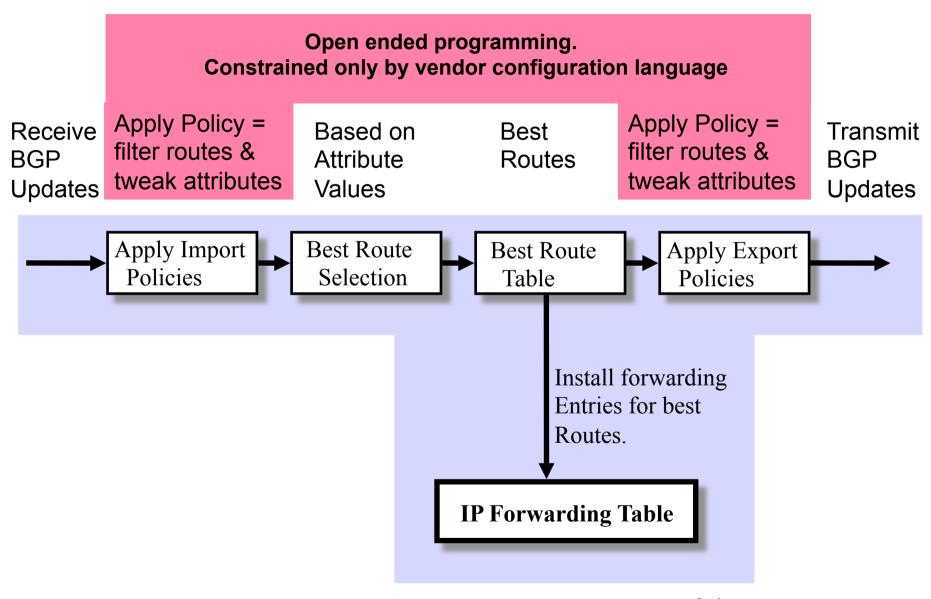
#### **BGP Attributes**

ue 1 2 3 4 5 6 7 8 9 L0 L1 L2 L3 L4 L5 L6 55	Code	Reference [RFC1771] [RFC1771] [RFC1771] [RFC1771] [RFC1771] [RFC1771] [RFC1771] [RFC2796] [RFC2796] [Chen] [RFC1863] [RFC1863] [RFC283] [RFC2283] [RFC2283] [RFC2283]	Most importan attributes
--	------	---	--------------------------

From IANA: http://www.iana.org/assignments/bgp-parameters

Not all attributes need to be present in every announcement

## **BGP Route Processing**



## **Route Selection Summary**

**Highest Local Preference** 

**Enforce relationships** 

**Shortest ASPATH** 

**Lowest MED** 

i-BGP < e-BGP

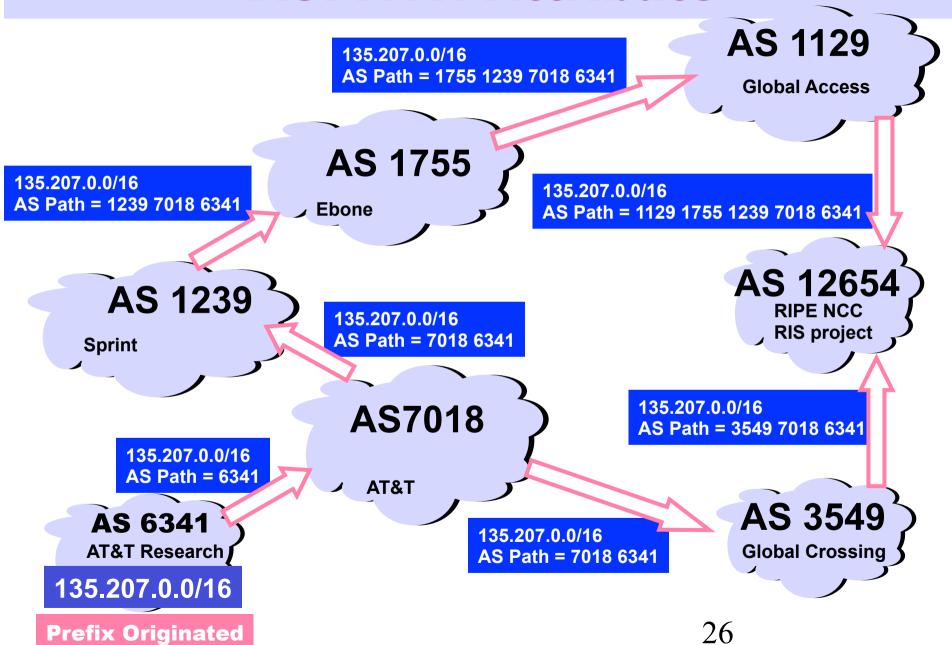
traffic engineering

Lowest IGP cost to BGP egress

**Lowest router ID** 

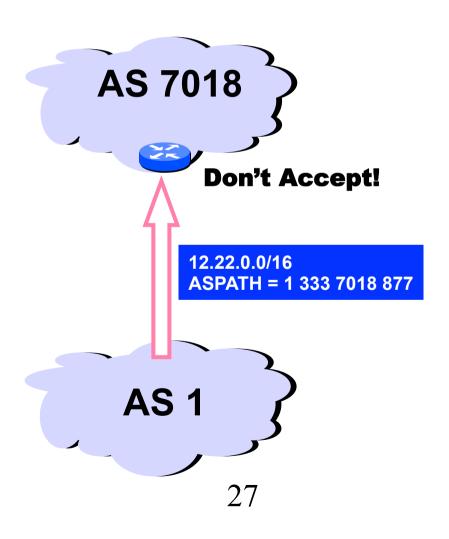
Throw up hands and break ties

#### **ASPATH Attribute**

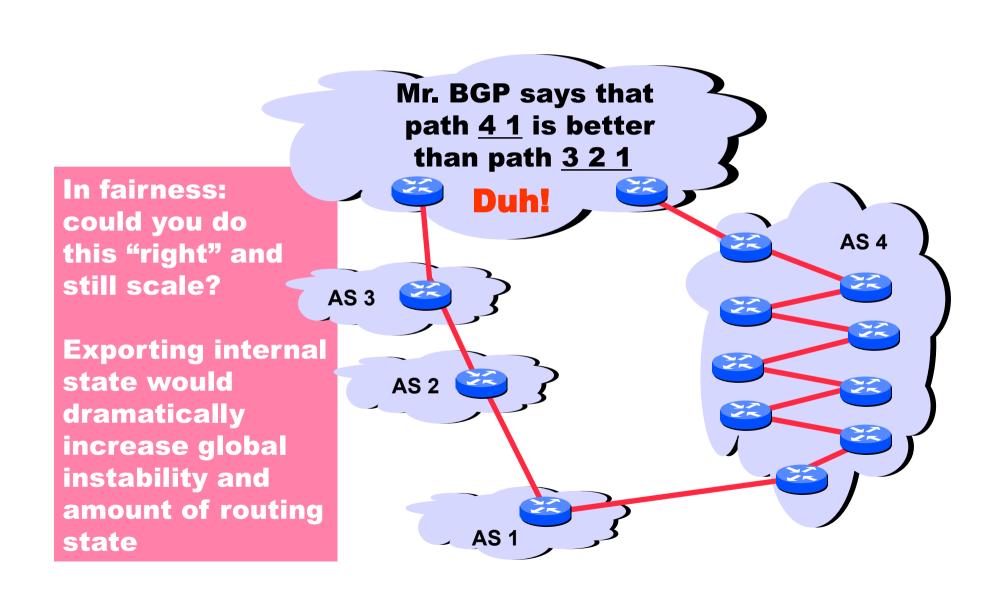


## **Interdomain Loop Prevention**

BGP at AS YYY will never accept a route with ASPATH containing YYY.



## Shorter Doesn't Always Mean Shorter



## **BGP Routing Tables**

0 65056 4637 174 21889 i

0 12654 7018 174 21889 i

0 65056 4637 174 21889 i

0 12654 7018 174 21889 i

0 65056 4637 174 21889 i 0 4608 1221 4637 174 21889 i

0 12654 7018 174 21889 i

0 65056 4637 174 21889 i

0 12654 7018 174 21889 i

0 65056 4637 174 21889 i

0 12654 7018 174 21889 i

0 65056 4637 174 21889 i

0 4608 1221 4637 174 21889 i

0 12654 3741 10310 14780 i

0 65056 4637 10310 14780 i

0 4608 1221 4637 10310 14780 i

Network Next Hop 193.0.4.28 \*> 0.0.0.0 3.0.0.0 193.0.4.28 \*> 203.50.0.33 202.12.29.79 193.0.4.28 4.0.0.0 \*> 203.50.0.33 202.12.29.79 193.0.4.28 4.0.0.0/9 \*> 203.50.0.33 202.12.29.79 193.0.4.28 4.23.112.0/24 203.50.0.33 202.12.29.79 4.23.113.0/24 193.0.4.28 203.50.0.33 202.12.29.79 4.23.114.0/24 193.0.4.28 203.50.0.33 202.12.29.79 4.36.116.0/23 193.0.4.28 203.50.0.33 202.12.29.79 193.0.4.28 4.36.116.0/24 \*> 203.50.0.33 202.12.29.79 4.36.117.0/24 193.0.4.28 203.50.0.33 202.12.29.79 4.36.118.0/24 193.0.4.28 203.50.0.33 202.12.29.79 \*> 4.78.22.0/23 193.0.4.28 203.50.0.33 202.12.29.79 \*> 4.78.56.0/23 193.0.4.28 203.50.0.33 202.12.29.79 4.79.181.0/24 193.0.4.28 \*> 203.50.0.33 202.12.29.79

```
Thanks to Geoff Huston.
Metric LocPrf Weight Path
                                            http://bgp.potaroo.net on Feb 1, 2008
                 0 12654 34225 1299 i
                 0 12654 7018 701 703 80 i
                 0 65056 4637 703 80 i
                 0 4608 1221 4637 703 80 i
                 0 12654 7018 3356 i
                 0 65056 4637 3356 i
                 0 4608 1221 4637 3356 i
                 0 12654 7018 3356 i
                 0 65056 4637 3356 i
                 0 4608 1221 4637 3356 i
                 0 12654 7018 174 21889 i
                 0 65056 4637 174 21889 i
                 0 4608 1221 4637 174 21889 i
                 0 12654 7018 174 21889 i
```

0 12654 3257 19151 13909 13909 13909 13909 13909 13909 13909 13909 13909 13909 13909 i

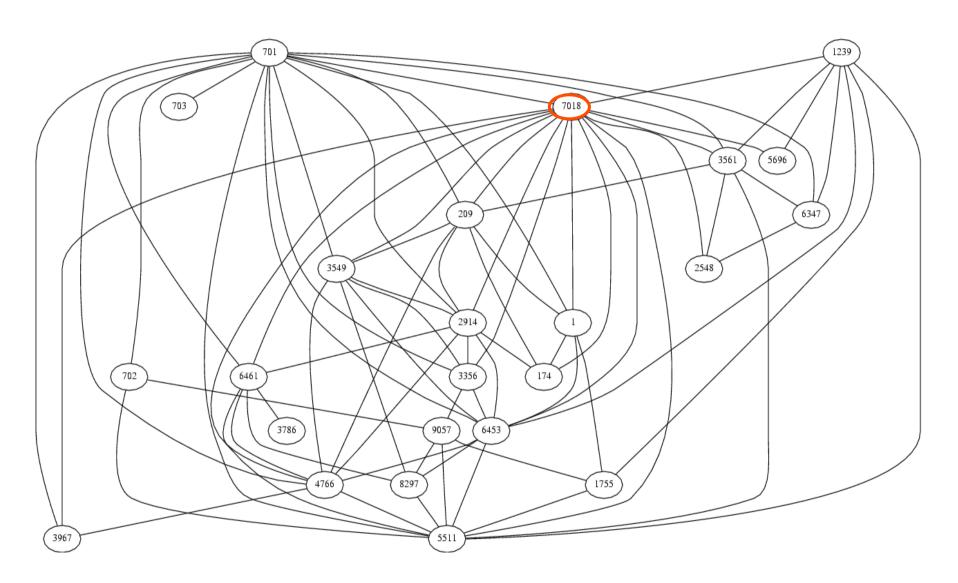
0 12654 3257 19151 13909 13909 13909 13909 13909 13909 13909 13909 13909 13909 i

0 65056 4637 1299 1239 19151 13909 13909 13909 13909 13909 13909 13909 13909 13909 13909 13909

0 4608 1221 4637 1299 1239 19151 13909 13909 13909 13909 13909 13909 13909 13909 13909 13909 13

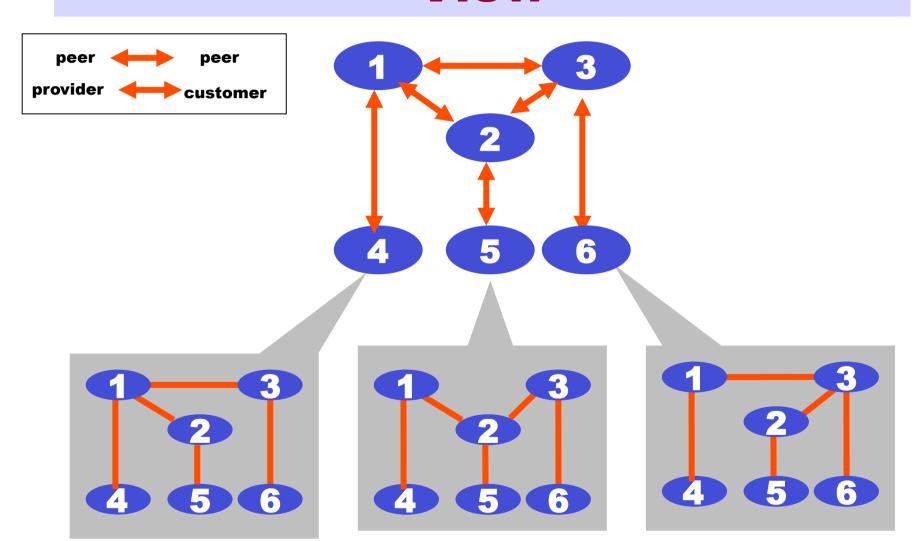
0 65056 4637 1299 1239 19151 13909 13909 13909 13909 13909 13909 13909 13909 13909 13909 13909 0 4608 1221 4637 1299 1239 19151 13909 1390

## **AS Graphs Can Be Fun**



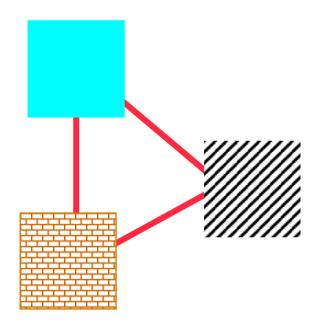
The <u>subgraph</u> showing all ASes that have more than 100 neighbors in full graph of 11,158 nodes. July 6, 2001. Point of view: AT&T route-server

## AS Graphs Depend on Point of View

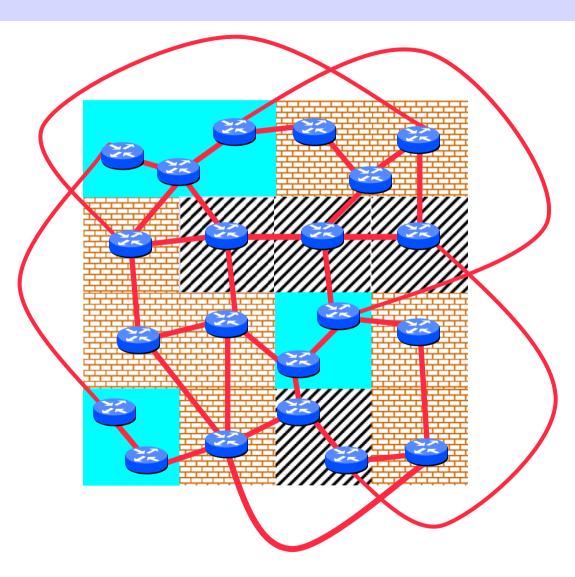


#### AS Graphs Do Not Show "Topology"!

**BGP** was designed to throw away information!



The AS graph may look like this.



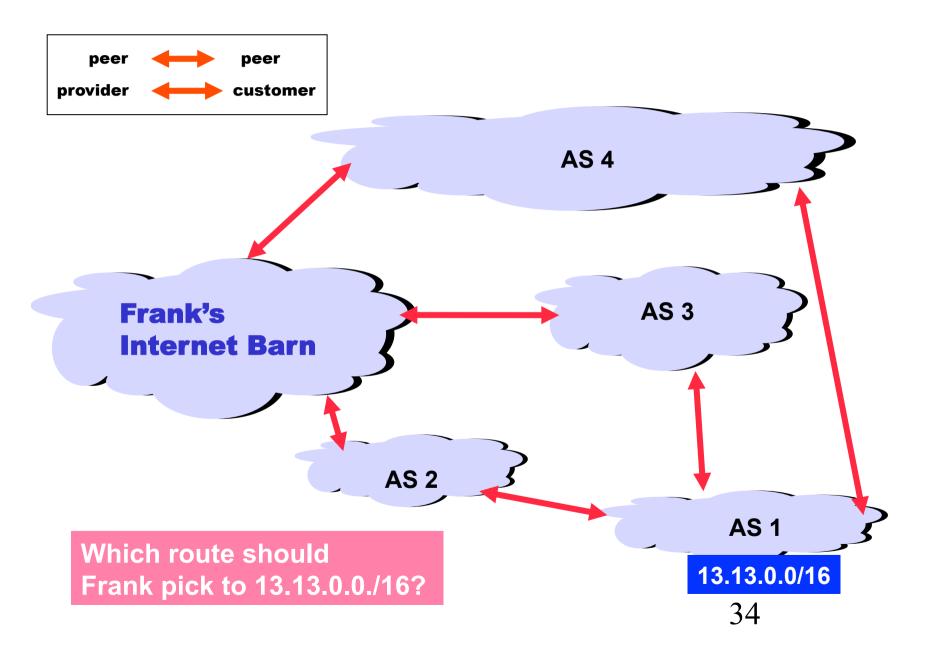
Reality may be closer to this...

## Implementing Customer/Provider and Peer/Peer relationships

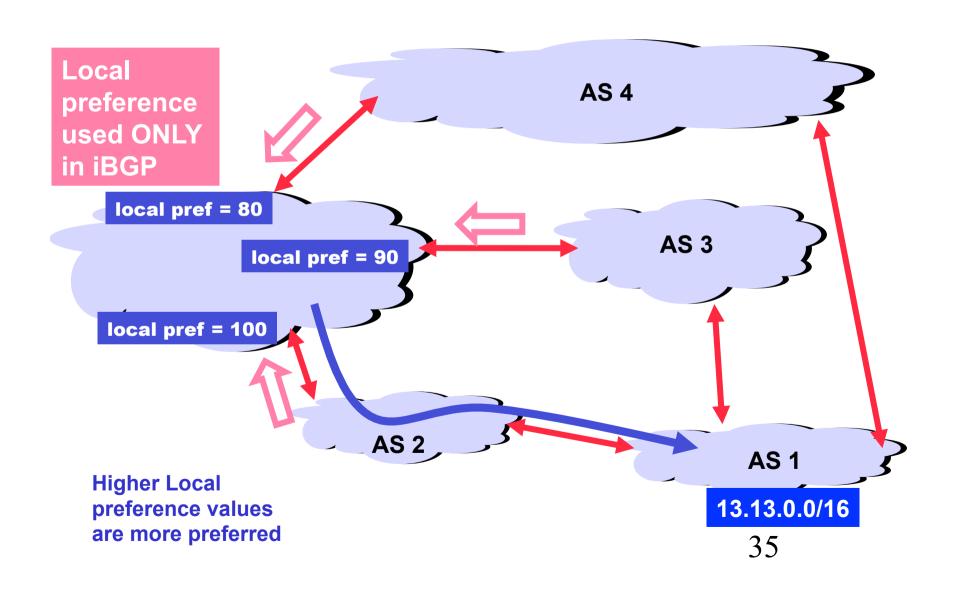
#### Two parts:

- Enforce transit relationships
  - Export all (best) routes to customers
  - Send only own and customer routes to all others
- Enforce order of route preference
  - provider < peer < customer</pre>

## So Many Choices

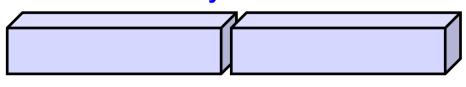


### LOCAL PREFERENCE



# How Can Routes be Classified? BGP Communities!

A community value is 32 bits



Used for signally within and between ASes

By convention, first 16 bits is ASN indicating who is giving it an interpretation

community

Very powerful BECAUSE it has no (predefined) meaning

Community Attribute = a list of community values. (So one route can belong to multiple communities)

**RFC 1997 (August 1996)** 

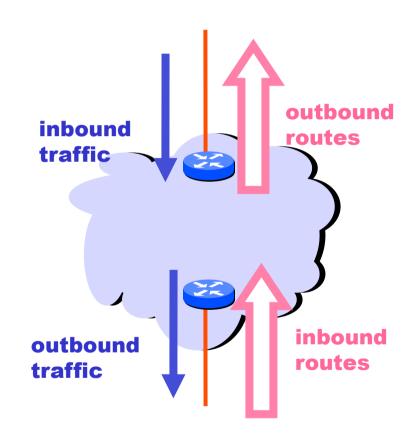
#### **Reserved communities**

no\_export = 0xFFFFFF01: don't export out of AS

no\_advertise 0xFFFFF02: don't pass to BGP neighbors

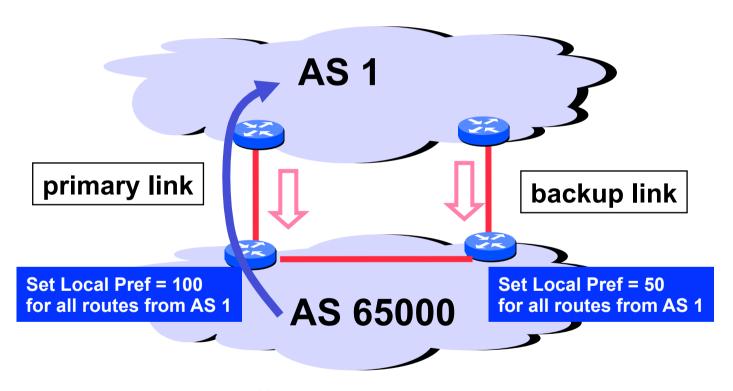
## Tweak Tweak (TE)

- For <u>inbound</u> traffic
  - Filter outbound routes
  - Tweak attributes on <u>outbound</u> routes in the hope of influencing your neighbor's best route selection
- For <u>outbound</u> traffic
  - Filter <u>inbound</u> routes
  - Tweak attributes on <u>inbound</u> routes to influence best route selection



In general, an AS has more control over outbound traffic

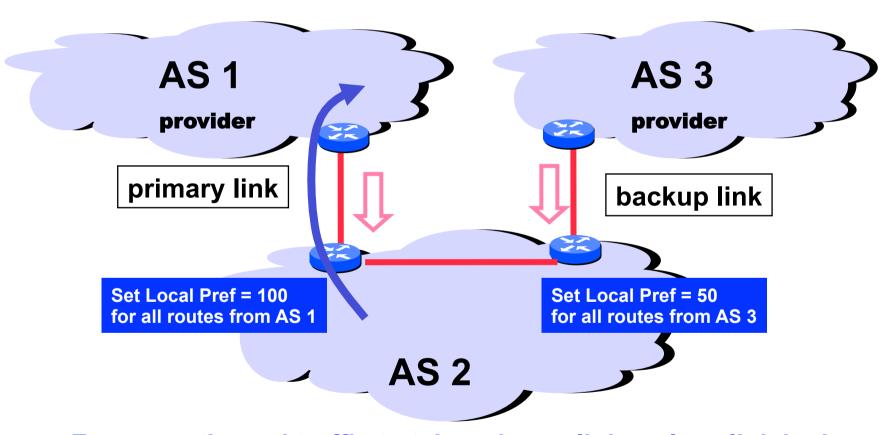
## Implementing Backup Links with Local Preference (Outbound Traffic)



Forces <u>outbound</u> traffic to take primary link, unless link is down.

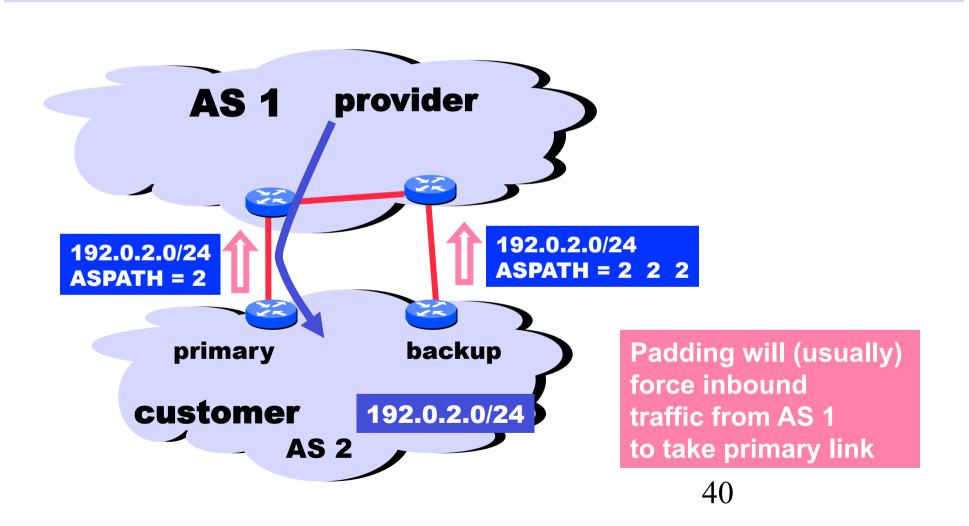
We'll talk about inbound traffic soon ...

# Multihomed Backups (Outbound Traffic)

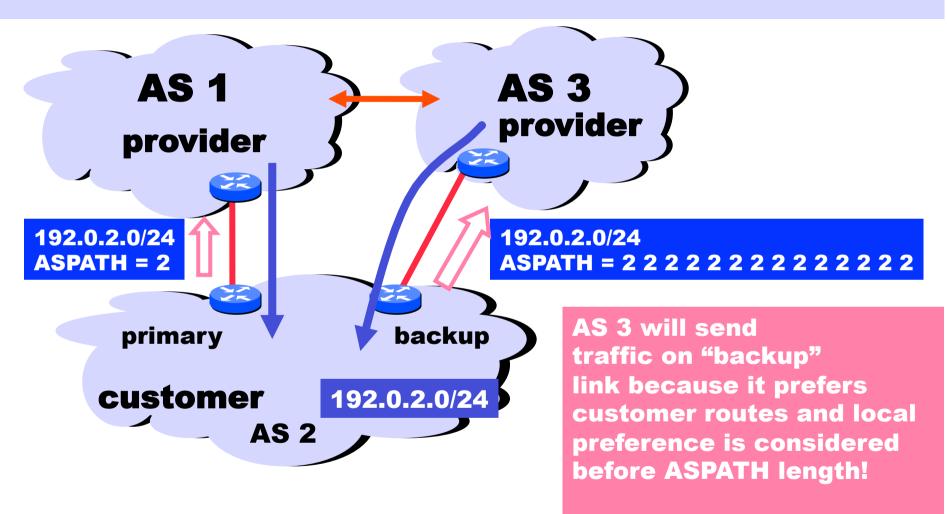


Forces <u>outbound</u> traffic to take primary link, unless link is down.

## Shedding Inbound Traffic with ASPATH Padding. Yes, this is a Glorious Hack ....

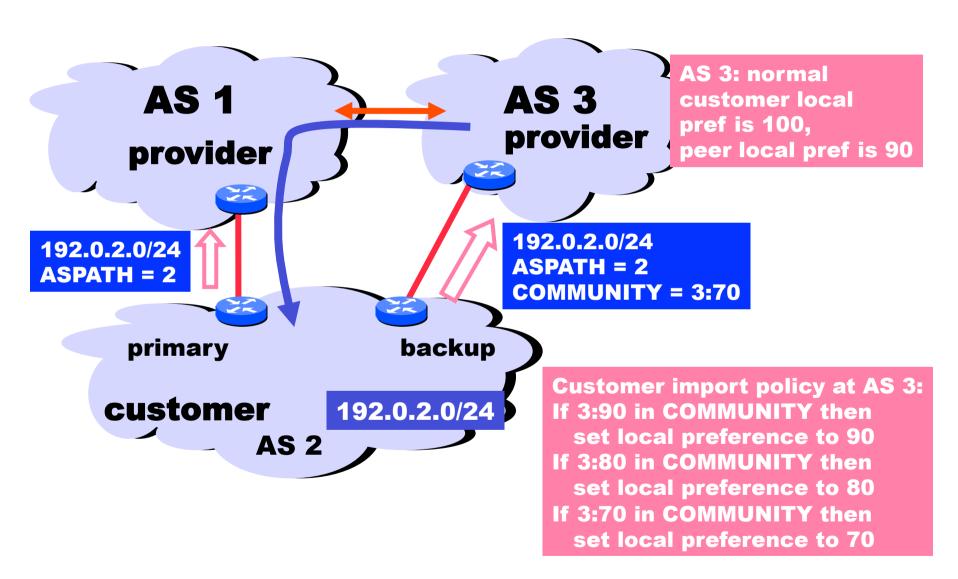


#### ... But Padding Does Not Always Work

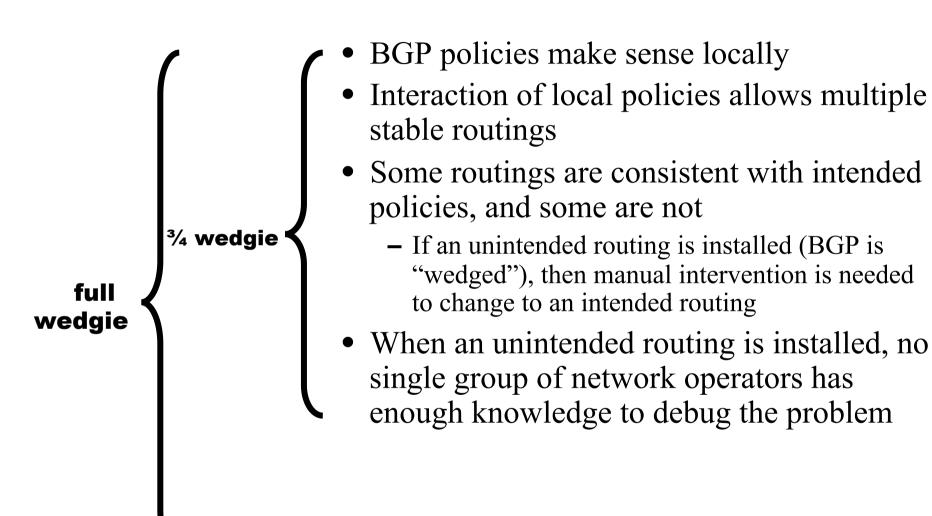


Padding in this way is often used as a form of load balancing

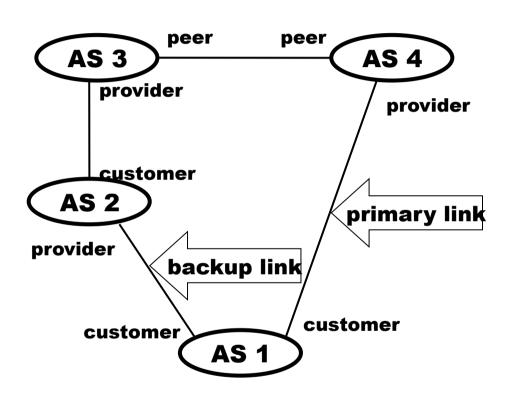
#### **COMMUNITY Attribute to the Rescue!**



### What is a BGP Wedgie (RFC 4264)?

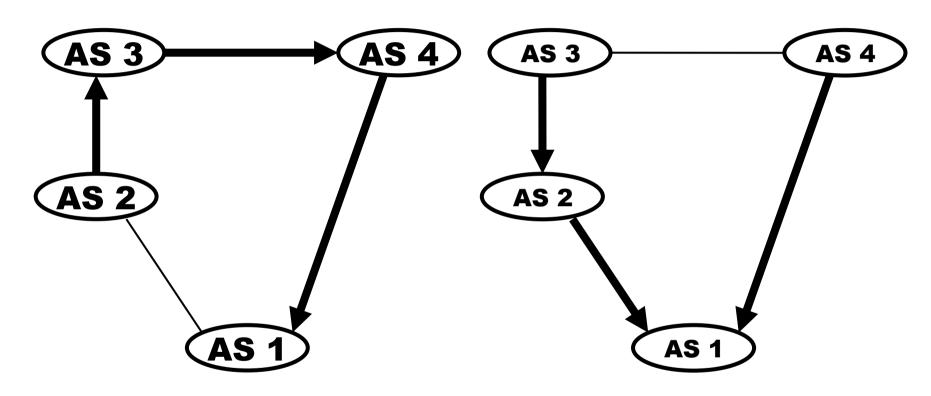


## 3/4 Wedgie Example



- AS 1 implements backup link by sending AS 2 a "depref me" community.
- AS 2 implements this community so that the resulting local pref is below that of routes from it's upstream provider (AS 3 routes)

### And the Routings are...



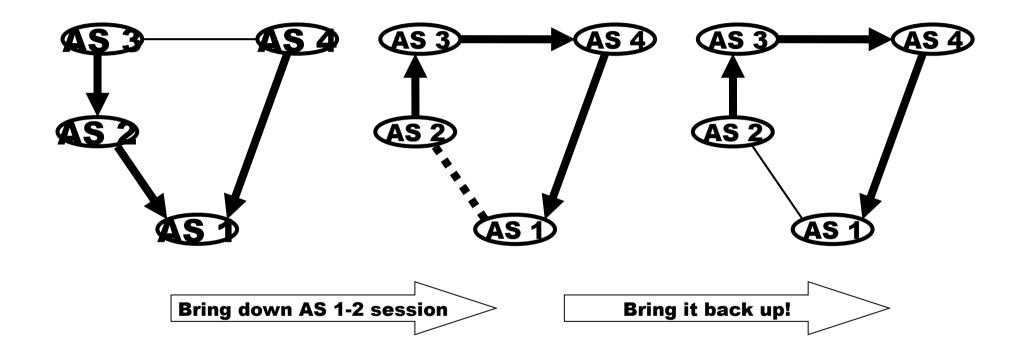
#### **Intended Routing**

Note: this would be the ONLY routing if AS2 translated its "depref me" community to a "depref me" community of AS 3

#### **Unintended Routing**

Note: This is easy to reach from the intended routing just by "bouncing" the BGP session on the primary link.

### Recovery

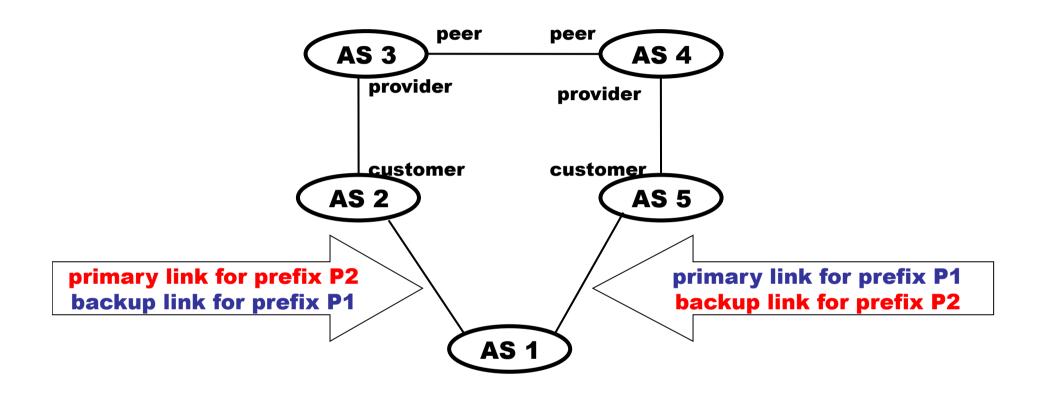


- Requires manual intervention
- Can be done in AS 1 or AS 2

## What the heck is going on?

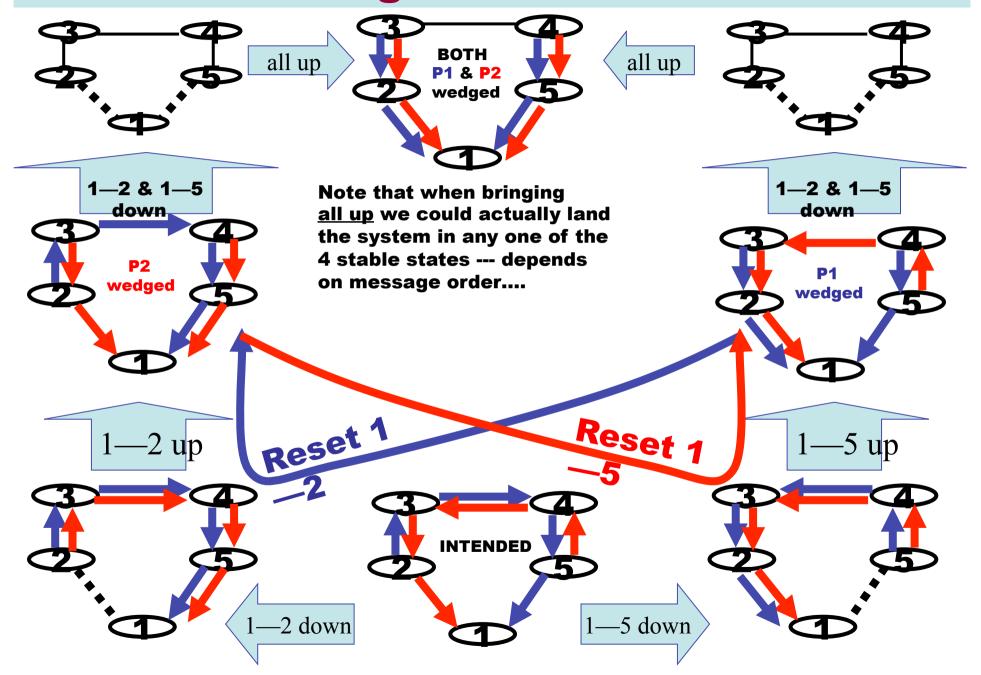
- There is no guarantee that a BGP configuration has a unique routing solution.
  - When multiple solutions exist, the (unpredictable) order of updates will determine which one is wins.
- There is no guarantee that a BGP configuration has any solution!
  - And checking configurations NP-Complete
  - Lab demonstrations of BGP configs never converging
- Complex policies (weights, communities setting preferences, and so on) increase chances of routing anomalies.
  - ... yet this is the current trend!

## **Load Balancing Example**

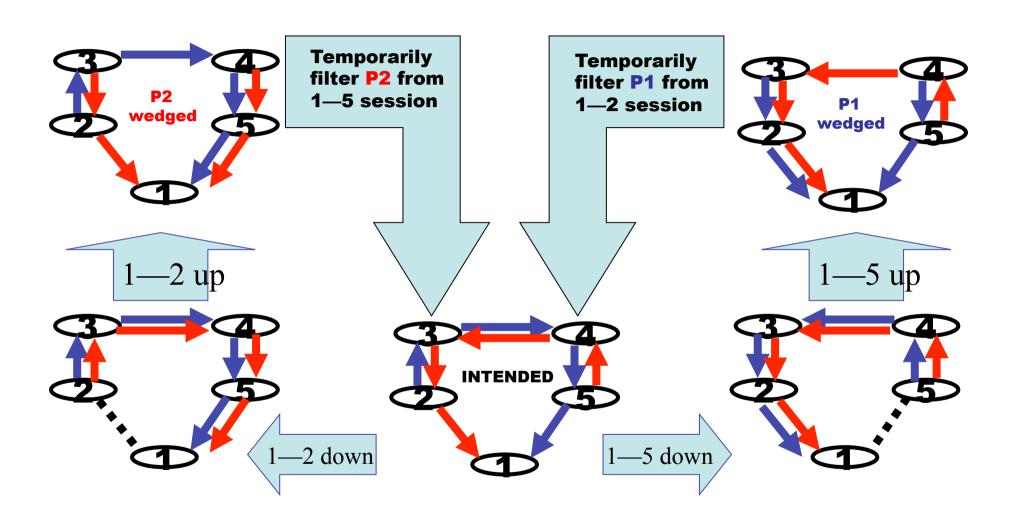


Simple session reset my not work!!

#### Can't un-wedge with session resets!

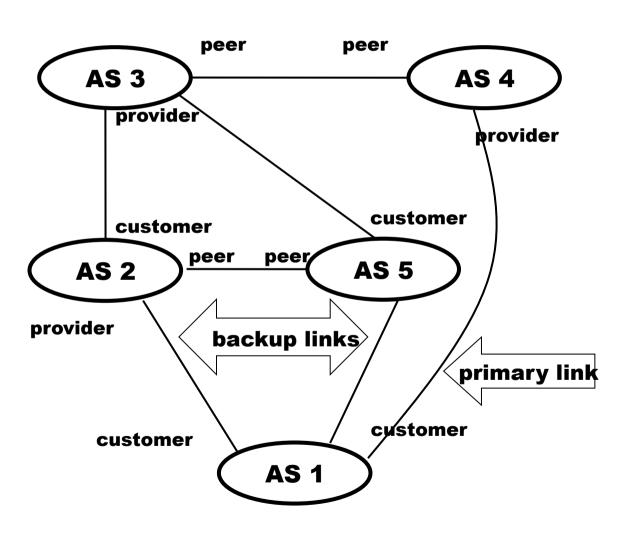


#### Recovery



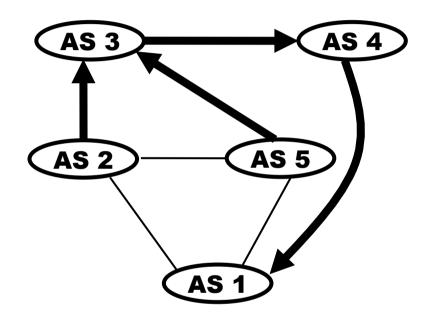
Who among us could figure this one out? When 1—2 is in New York and 1—5 is in Tokyo?

## Full Wedgie Example

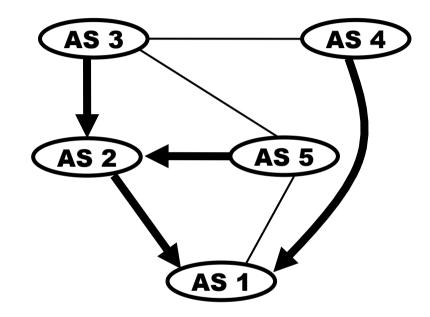


- AS 1 implements backup links by sending AS 2 and AS 3 a "depref me" communities.
- AS 2 implements its community so that the resulting local pref is below that of its upstream providers and it's peers (AS 3 and AS 5 routes)
- AS 5 implements its community so that the resulting local pref is below its peers (AS 2) but above that of its providers (AS 3)

## And the Routings are...

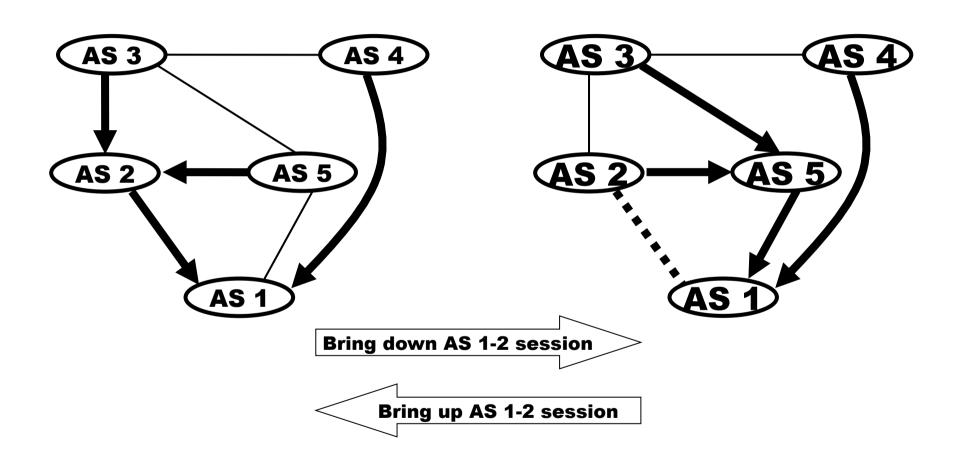


**Intended Routing** 

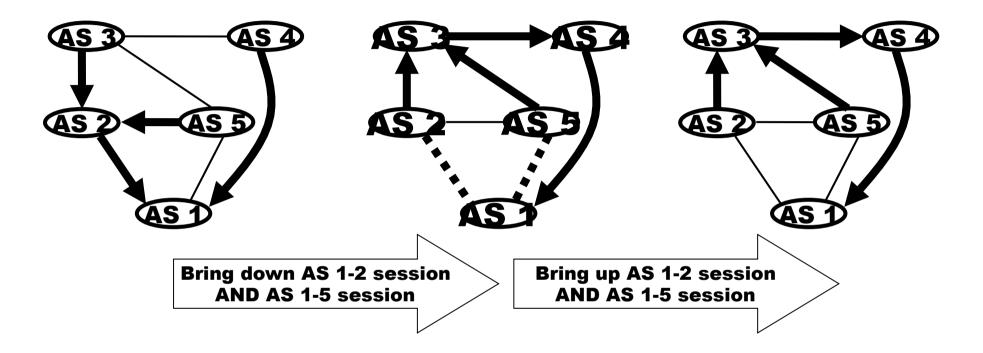


**Unintended Routing** 

#### Resetting 1—2 does not help!!



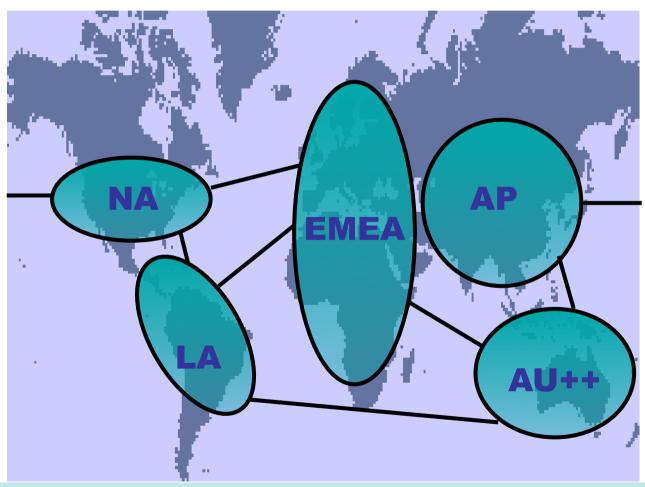
### Recovery



A lot of "non-local" knowledge is required to arrive at this recovery strategy!

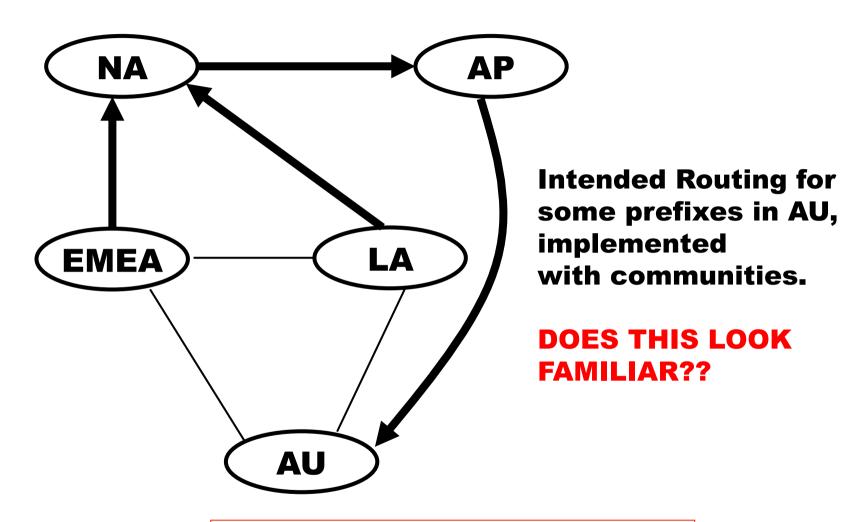
Try to convince AS 5 and AS 1 that their session has be reset (or filtered) even though it is not associated with an active route!

#### That Can't happen in MY network!!



An "normal" global global backbone (ISP or Corporate Intranet) implemented with 5 regional ASes

#### The Full Wedgie Example, in a new Guise



Message: Same problems can arise with "traffic engineering" across regional networks.