

In ML,
A ⊔ B type

(α, β) disjoint union

= one of α | two of β

Disjoint unions

For a set X, {l} × X is a set of elements of X tagged by l

Definition 75 The disjoint union $A \uplus B$ of two sets A and B is the set

$$A \uplus B = (\{1\} \times A) \cup (\{2\} \times B) .$$

Thus,

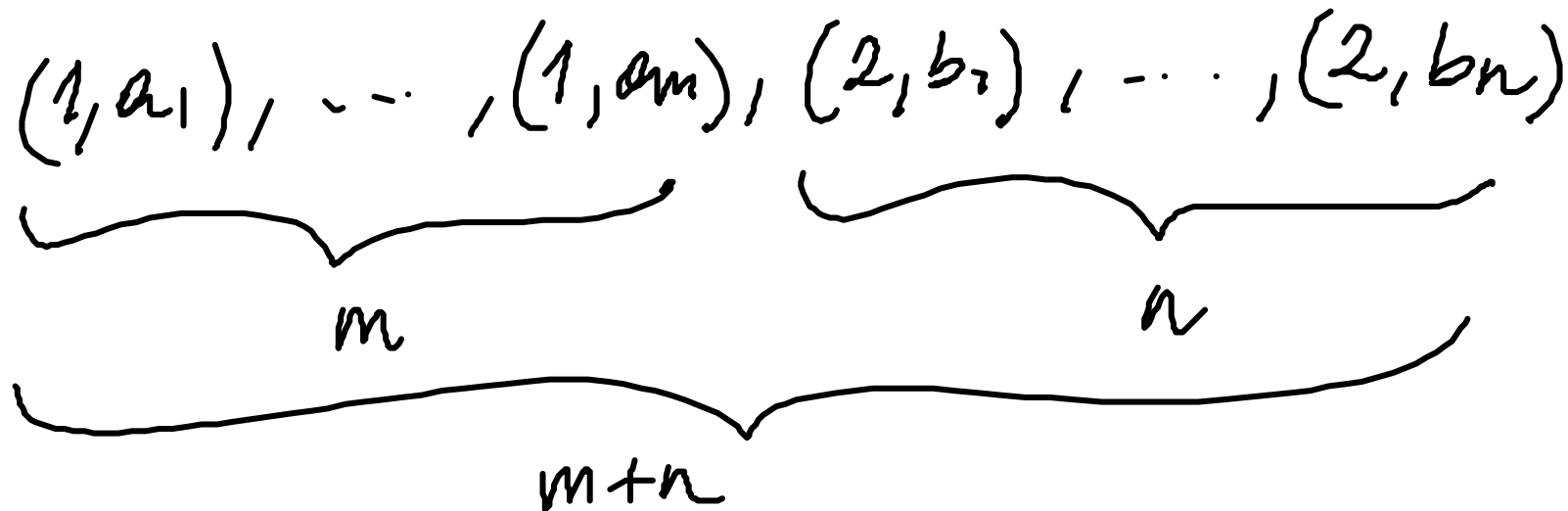
$$\forall x. x \in (A \uplus B) \iff (\exists a \in A. x = (1, a)) \vee (\exists b \in B. x = (2, b)) .$$

Proposition 76 For all finite sets A and B ,

$$\#(A \uplus B) = \#A + \#B .$$

PROOF IDEA: $\#A = m$, $A = \{a_1, \dots, a_m\}$
 $\#B = n$, $B = \{b_1, \dots, b_n\}$

$A \uplus B$



sets of ordered pairs.

Relations

domain
(or source)

Definition 77 A (binary) relation R from a set A to a set B

$$R : A \rightarrow B \quad \text{or} \quad R \in \text{Rel}(A, B) ,$$

is

$$\begin{array}{c} \Downarrow \\ \mathcal{P}(A \times B) \end{array}$$

$$R \subseteq A \times B$$

codomain
(a target)

Notation 78 One typically writes $a R b$ for $(a, b) \in R$.

Informal examples:

- ▶ Computation.
- ▶ Typing.
- ▶ Program equivalence.
- ▶ Networks.
- ▶ Databases.

Examples:

- ▶ Empty relation.

$$\emptyset : A \dashrightarrow B$$

$$(a \emptyset b \iff \text{false})$$

- ▶ Full relation.

$$(A \times B) : A \dashrightarrow B$$

$$(a (A \times B) b \iff \text{true})$$

- ▶ Identity (or equality) relation.

$$I_A = \{ (a, a) \mid a \in A \} : A \dashrightarrow A$$

$$(a I_A a' \iff a = a')$$

- ▶ Integer square root.

$$R_2 = \{ (m, n) \mid m = n^2 \} : \mathbb{N} \dashrightarrow \mathbb{Z}$$

$$(m R_2 n \iff m = n^2)$$

eg. $4 R_2 2, 4 R_2 (-2)$

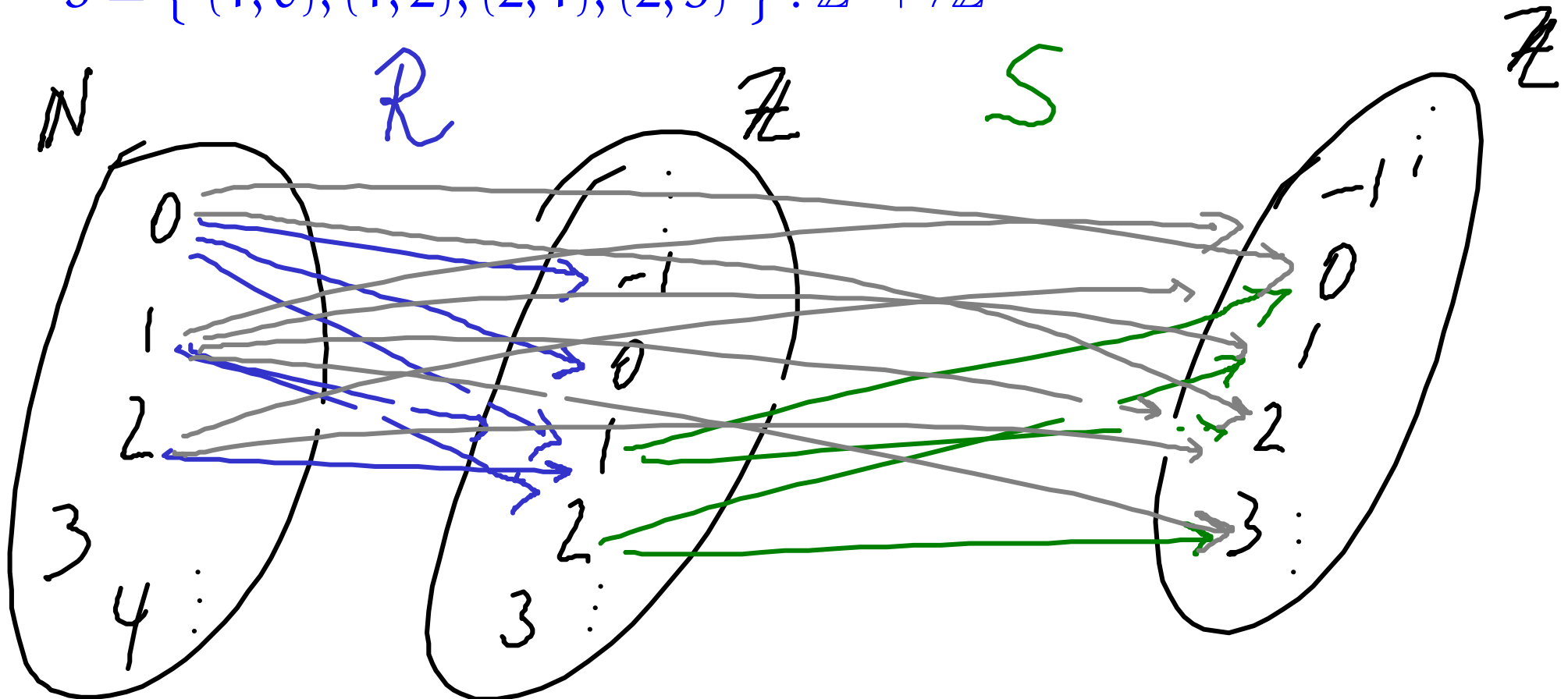
Internal diagrams

SoR

Example:

$$R = \{ (0, 0), (0, -1), (0, 1), (1, 2), (1, 1), (2, 1) \} : \mathbb{N} \rightarrow \mathbb{Z}$$

$$S = \{ (1, 0), (1, 2), (2, 1), (2, 3) \} : \mathbb{Z} \rightarrow \mathbb{Z}$$



Relational composition

Given

$$R: A \rightarrow B, \quad S: B \rightarrow C$$

we define

$$S \circ R: A \rightarrow C$$

$$\left[\begin{array}{c} \text{notation} \\ \hline R; S: A \rightarrow C \end{array} \right]$$

by

$$a(S \circ R)c \iff \exists b \in B. aRb \ \& \ bSc$$

It is unambiguous to write
 $T \circ S \circ R$

Theorem 79 Relational composition is associative and has the identity relation as neutral element.

► Associativity.

For all $R : A \rightarrow B$, $S : B \rightarrow C$, and $T : C \rightarrow D$,

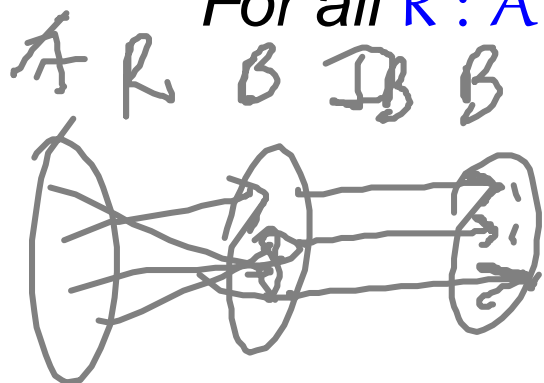
$$(T \circ S) \circ R = T \circ (S \circ R)$$

Exercise

► Neutral element.

For all $R : A \rightarrow B$,

$$R \circ I_A = R = I_B \circ R$$



Show
 $a [(T \circ S) \circ R] d$
 iff
 $a [T \circ (S \circ R)] d$

Relations and matrices

Definition 80

1. For positive integers m and n , an $(m \times n)$ -matrix M over a semiring $(S, 0, \oplus, 1, \odot)$ is given by entries $M_{i,j} \in S$ for all $0 \leq i < m$ and $0 \leq j < n$.

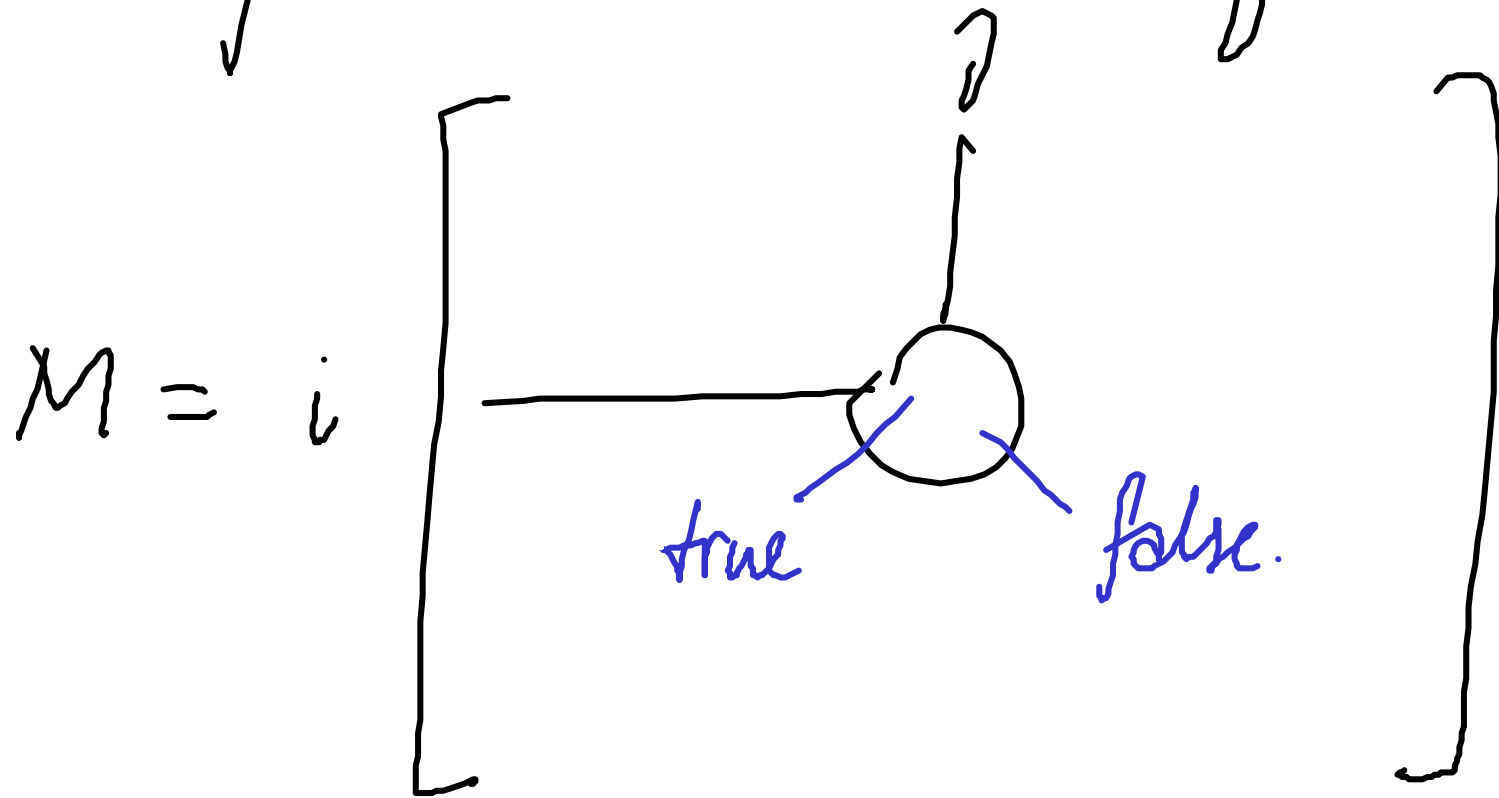
$(m \times n)$ -matrix M $(n \times l)$ -matrix L

\rightsquigarrow $(m \times l)$ -matrix $(L \cdot M)$

$$(L \cdot M)_{i,j} = \bigoplus_{k=0}^{n-1} L_{i,k} \odot M_{k,j}$$

Theorem 81 Matrix multiplication is associative and has the identity matrix as neutral element.

What if S is the Boolean semiring.



Every M gives us a relation $\underline{\text{rel}}(M) : [m] \rightarrow [n]$

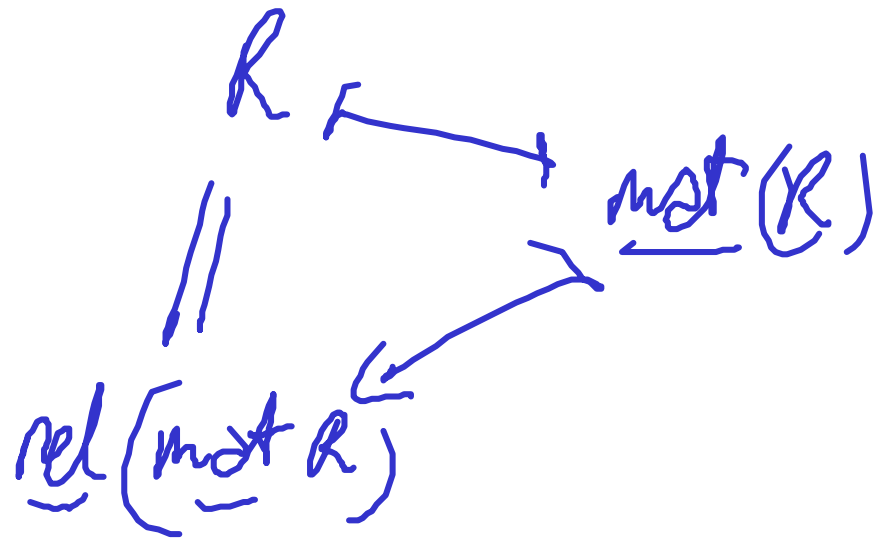
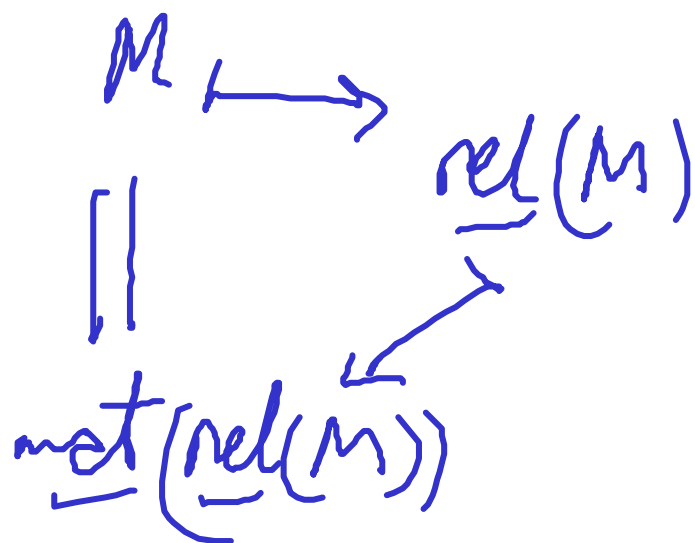
where $[k] = \{0, 1, \dots, k-1\}$ defined by

$$i (\underline{\text{rel}} M) j \iff M_{ij} = \text{true}$$

Given a relation $R: [m] \rightarrow [n]$

defines $\underline{\text{mat}}(R)$ by

$$\underline{\text{mat}}(R)_{ij} = \begin{cases} \text{true} & \text{if } iRj \\ \text{false} & \text{otherwise} \end{cases}$$



$$R, S \longmapsto \underline{\text{mat}}(R), \underline{\text{mat}}(S)$$

$$S \circ R \longmapsto \underline{\text{mat}}(S \circ R)$$

|| fact

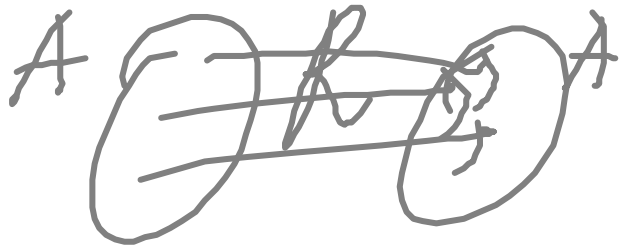
$$\underline{\text{mat}}(S) \cdot \underline{\text{mat}}(R)$$

$$\text{Id}_A \longmapsto \underline{\text{mat}}(\text{Id}_A) = \begin{bmatrix} \text{true} & \text{false} \\ \text{false} & \text{true} \end{bmatrix}$$

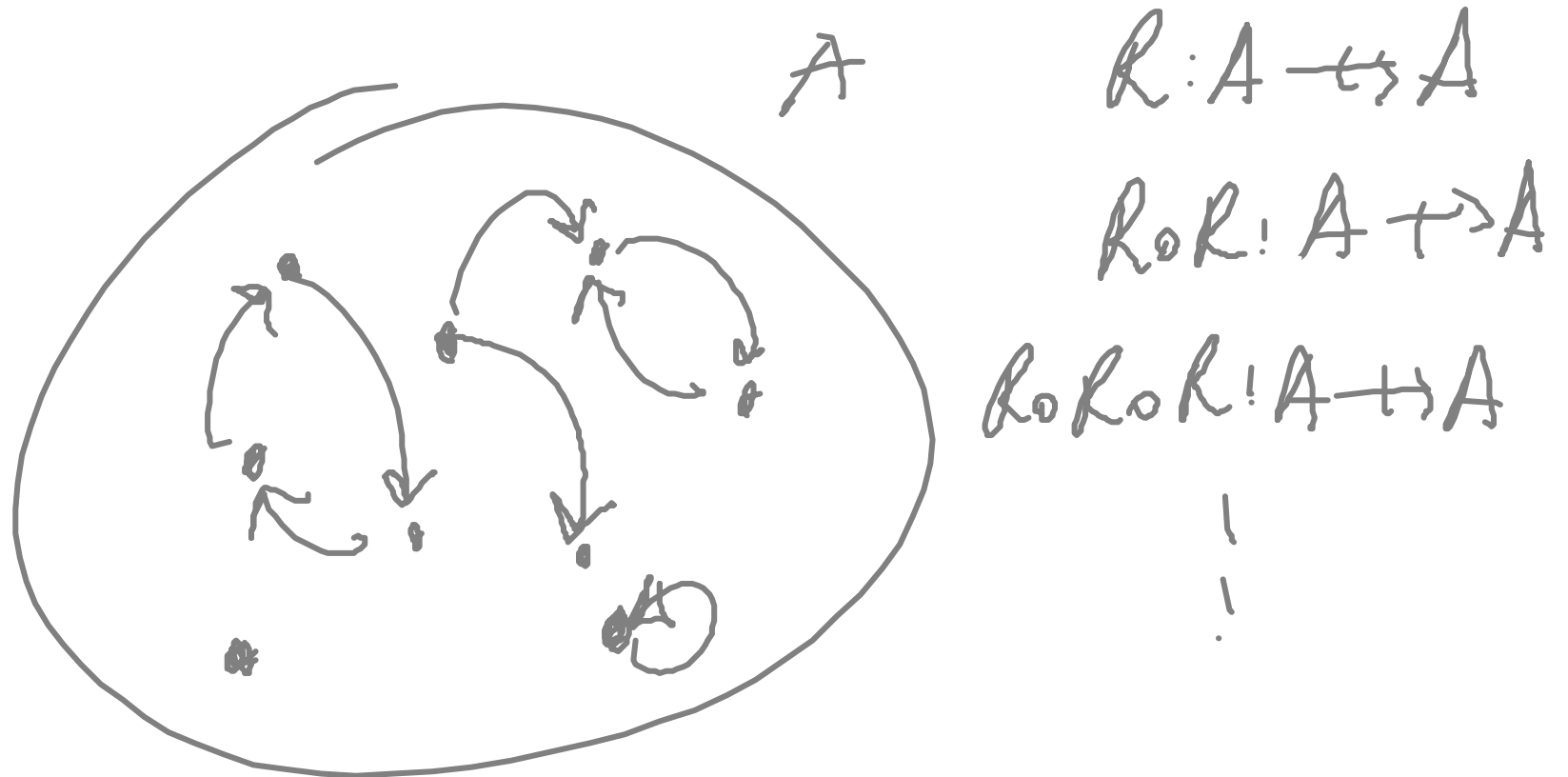
Relations from $[m]$ to $[n]$ and $(m \times n)$ -matrices over Booleans provide two alternative views of the same structure.

This carries over to identities and to composition/multiplication .

Directed graphs A



Definition 82 A directed graph (A, R) consists of a set A and a relation R on A (i.e. a relation from A to A).



Corollary 84 For every set A , the structure

$$(\text{Rel}(A), I_A, \circ)$$

is a monoid.

Definition 85 For $R \in \text{Rel}(A)$ and $n \in \mathbb{N}$, we let

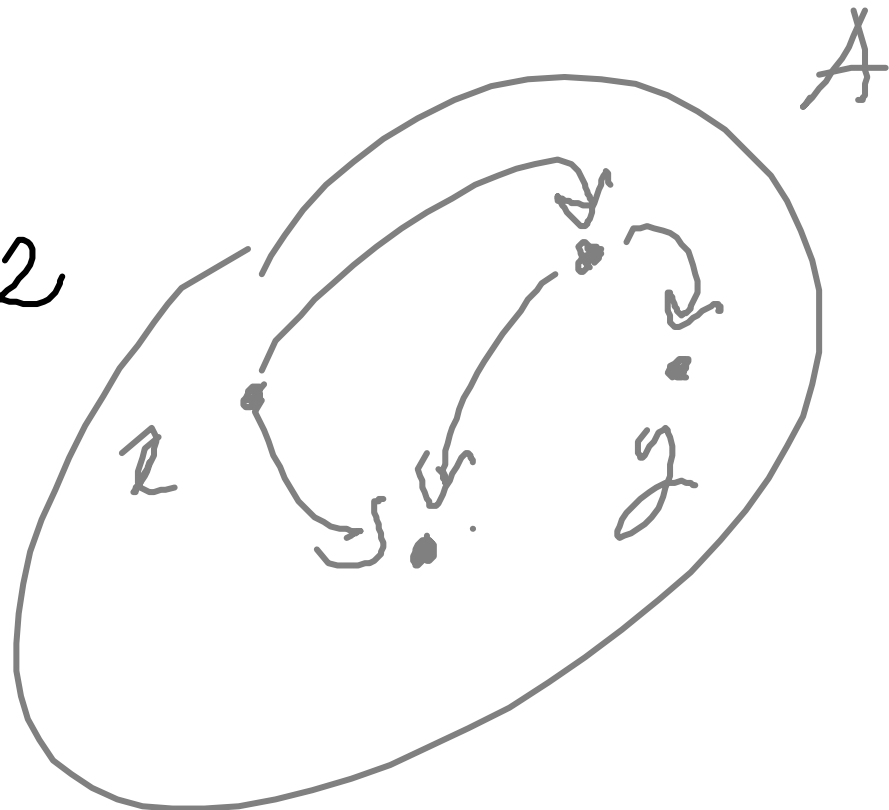
$$R^{\circ n} = \underbrace{R \circ \dots \circ R}_{n \text{ times}} \in \text{Rel}(A)$$

be defined as I_A for $n = 0$, and as $R \circ R^{\circ m}$ for $n = m + 1$.

$R \circ R: A \rightarrow A$

$x (R \circ R) y \iff \exists z: x R z \ \& \ z R y$

\iff There is a path of length 2 in the directed graph



Claim $x R^n y \iff$ There is a path of length n in the directed graph.

Paths

Proposition 86 *Let (A, R) be a directed graph. For all $n \in \mathbb{N}$ and $s, t \in A$, $s R^{o^n} t$ iff there exists a path of length n in R with source s and target t .*

PROOF:

EXERCISE (by induction)

Definition 87 For $R \in \text{Rel}(A)$, let

$$R^{o*} = \bigcup \{ R^{on} \in \text{Rel}(A) \mid n \in \mathbb{N} \} = \bigcup_{n \in \mathbb{N}} R^{on} .$$

Corollary 88 Let (A, R) be a directed graph. For all $s, t \in A$, $s R^{o*} t$ iff there exists a path with source s and target t in R .