Number systems

Objectives

- Get an appreciation for the abstract notion of number system, considering four examples: natural numbers, integers, rationals, and modular integers.
- Prove the correctness of three basic algorithms in the theory of numbers: the division algorithm, Euclid's algorithm, and the Extended Euclid's algorithm.
- Exemplify the use of the mathematical theory surrounding Euclid's Theorem and Fermat's Little Theorem in the context of public-key cryptography.

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Natural numbers

In the beginning there were the *natural numbers*

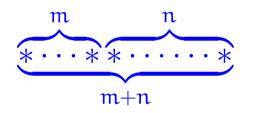
 \mathbb{N} : 0, 1, ..., n, n+1, ...

generated from zero by successive increment; that is, put in ML:

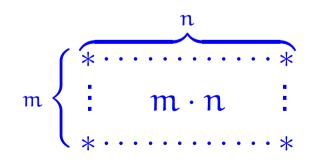
datatype
N = zero | succ of N

The basic operations of this number system are:

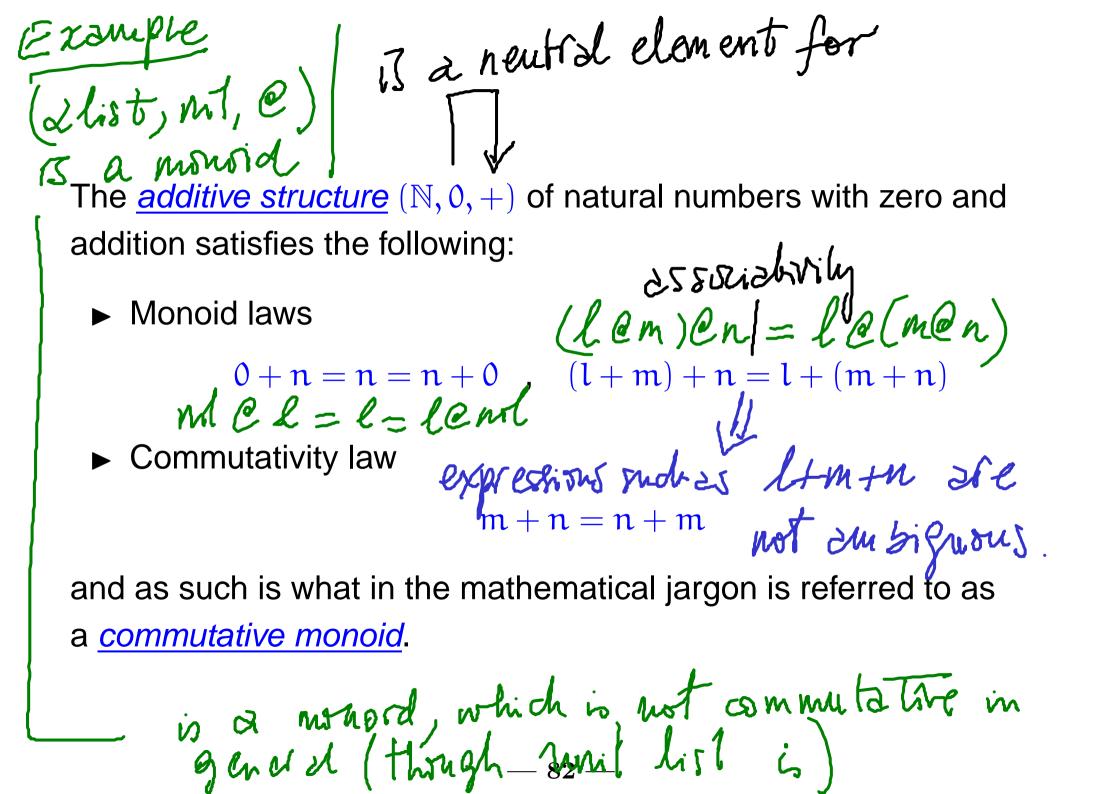








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Also the *multiplicative structure* $(\mathbb{N}, 1, \cdot)$ of natural numbers with one and multiplication is a commutative monoid:

Monoid laws

$$1 \cdot n = n = n \cdot 1$$
, $(l \cdot m) \cdot n = l \cdot (m \cdot n)$

Commutativity law

 $\mathbf{m} \cdot \mathbf{n} = \mathbf{n} \cdot \mathbf{m}$

The additive and multiplicative structures interact nicely in that they satisfy the

► Distributive law

 $l \cdot (m+n) = l \cdot m + l \cdot n$ $l \left[l \cdot (mn) \\ m + n \right] \qquad l \left[l \cdot m \\ m \\ n \right]$

and make the overall structure $(\mathbb{N}, 0, +, 1, \cdot)$ into what in the mathematical jargon is referred to as a *commutative semiring*.

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Cancellation

The additive and multiplicative structures of natural numbers further satisfy the following laws.

► Additive cancellation

For all natural numbers k, m, n,

$$k+m=k+n \implies m=n$$

Multiplicative cancellation

For all natural numbers k, m, n,

if $k \neq 0$ then $k \cdot m = k \cdot n \implies m = n$.

In general, let (M, e, *) be à commutative A An element zim Mass on inverse monsid if There is a gin M such That X * y= e. Inverses z 2 * z * y = (z * 2) * y = e * y = y **Definition 37** 1. A number x is said to admit an <u>additive inverse</u> whenever there exists a number y such that x + y = 0. Propositions. The inverse of on element is unique. PROOF Let y and z be inverses for z; That is, $\widehat{T} \times y = e$ and $\overline{z} \times z = e$. We show y = z. By () 2# X* 7= 2* C= 2:

there exists How do ve prore unique $\exists z \cdot P(z)$ $(\Rightarrow (\exists z, P(z)) \&$ $\left[\forall z_1, x_2, P(z_1) \& P(x_2) \Rightarrow (z_1 = x_2) \right]$

Inverses

Definition 37

- 1. A number x is said to admit an additive inverse whenever there exists a number y such that x + y = 0.
- 2. A number x is said to admit a multiplicative inverse whenever there exists a number y such that $x \cdot y = 1$.

Extending the system of natural numbers (i) to admit all additive inverses and then (ii) to also admit all multiplicative inverses for non-zero numbers yields two very interesting results:

(i) the *integers*

 \mathbb{Z} : ... - n, ..., -1, 0, 1, ..., n, ...

which then form what in the mathematical jargon is referred to as a *commutative ring*, and

(ii) the <u>rationals</u> \mathbb{Q} which then form what in the mathematical jargon is referred to as a <u>field</u>.

The division theorem and algorithm

7 Idea Understand m/n but within Z.

Theorem 38 (Division Theorem) For every natural number m and positive natural number n, there exists a unique pair of integers q and r such that $q \ge 0$, $0 \le r < n$, and $m = q \cdot n + r$.

An Argument of with a hidden use of mothemotical induction in the form of the well-ordering principle Genren m and n, collect I the natural numbers of The form m- た n for k in The integers. In parbiaular, m=m-O.n so The above makes sense. From This collection, let r be the smallest such, Cloin r is as described in the Thm.