

Databases 2011

Lectures 04 – 05

Timothy G. Griffin

Computer Laboratory
University of Cambridge, UK

Databases, Easter 2011

Lectures 04 and 5 : Modeling with Entities and Relationships

Outline

- Entities
- Relationships
- Their relational implementations
- On the importance of SCOPE
- n-ary relationships
- Weak entity sets
- Generalization

Some real-world data ...

... from the Internet Movie Database (IMDb).

| Title | Year | Actor |
|---|------|------------|
| Austin Powers: International Man of Mystery | 1997 | Mike Myers |
| Austin Powers: The Spy Who Shagged Me | 1999 | Mike Myers |
| Dude, Where's My Car? | 2000 | Bill Chott |
| Dude, Where's My Car? | 2000 | Marc Lynn |

Entities diagrams and Relational Schema



These diagrams represent relational schema

Movie(MovieID, Title, Year)

Person(PersonID, FirstName, LastName)

Yes, this ignores types ...

Entity sets (relational instances)

Movie

| <u>MovieID</u> | Title | Year |
|----------------|---|------|
| 55871 | Austin Powers: International Man of Mystery | 1997 |
| 55873 | Austin Powers: The Spy Who Shagged Me | 1999 |
| 171771 | Dude, Where's My Car? | 2000 |

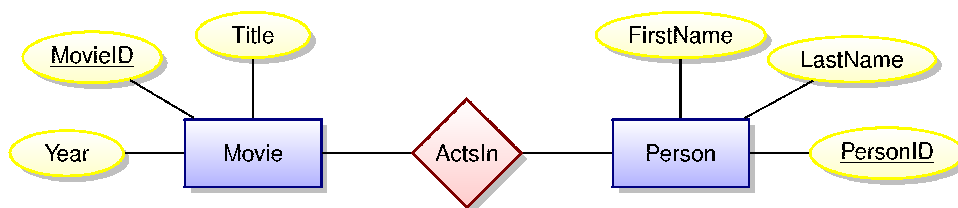
(I used line number from IMDb raw file movies.list as MovieID.)

Person

| <u>PersonID</u> | FirstName | LastName |
|-----------------|-----------|----------|
| 6902836 | Mike | Myers |
| 1757556 | Bill | Chott |
| 5882058 | Marc | Lynn |

(I used line number from IMDb raw file actors.list as PersonID)

Relationships



Foreign Keys and Referential Integrity

Foreign Key

Suppose we have $R(\underline{\mathbf{Z}}, \mathbf{Y})$. Furthermore, let $S(\mathbf{W})$ be a relational schema with $\mathbf{Z} \subseteq \mathbf{W}$. We say that \mathbf{Z} represents a **Foreign Key in S for R** if for any instance we have $\pi_{\mathbf{Z}}(S) \subseteq \pi_{\mathbf{Z}}(R)$. This is a semantic assertion.

Referential integrity

A database is said to have **referential integrity** when all foreign key constraints are satisfied.

A relational representation

A relational schema

ActsIn(MovieID, PersonID)

With **referential integrity constraints**

$$\pi_{\text{MovieID}}(\text{ActsIn}) \subseteq \pi_{\text{MovieID}}(\text{Movie})$$

$$\pi_{\text{PersonID}}(\text{ActsIn}) \subseteq \pi_{\text{PersonID}}(\text{Person})$$

ActsIn

| <u>PersonID</u> | <u>MovieID</u> |
|-----------------|----------------|
| 6902836 | 55871 |
| 6902836 | 55873 |
| 1757556 | 171771 |
| 5882058 | 171771 |

Foreign Keys in SQL

```
create table ActsIn
(  MovieID int not NULL,
   PersonID int not NULL,
   primary key (MovieID, PersonID),
   constraint actsin_movie
       foreign key (MovieID)
       references Movie(MovieID),
   constraint actsin_person
       foreign key (PersonID)
       references Person(PersonID))
```

Yes, we could do a better job on the date type ...

Relational representation of relationships, in general?


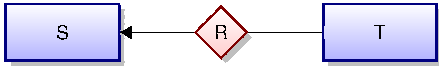

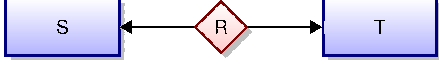
That depends ...

Mapping Cardinalities for binary relations, $R \subseteq S \times T$

| Relation R is | meaning |
|-----------------|--|
| many to many | no constraints |
| one to many | $\forall t \in T, s_1, s_2 \in S. (R(s_1, t) \wedge R(s_2, t)) \implies s_1 = s_2$ |
| many to one | $\forall s \in S, t_1, t_2 \in T. (R(s, t_1) \wedge R(s, t_2)) \implies t_1 = t_2$ |
| one to one | one to many and many to one |

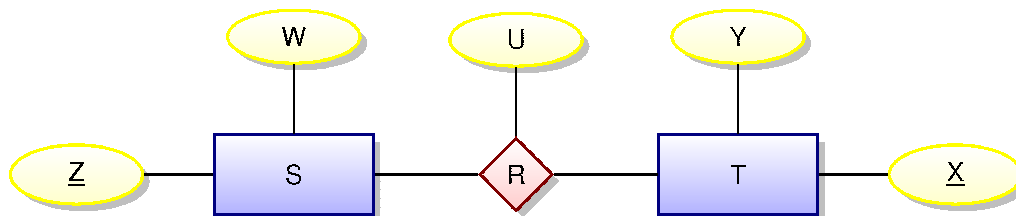
Note that the database terminology differs slightly from standard mathematical terminology.

Diagrams for Mapping Cardinalities

| ER diagram | Relation R is |
|---|--------------------------|
|  | many to many ($M : N$) |
|  | one to many ($1 : M$) |
|  | many to one ($M : 1$) |
|  | one to one ($1 : 1$) |

Navigation icons: back, forward, search, etc.

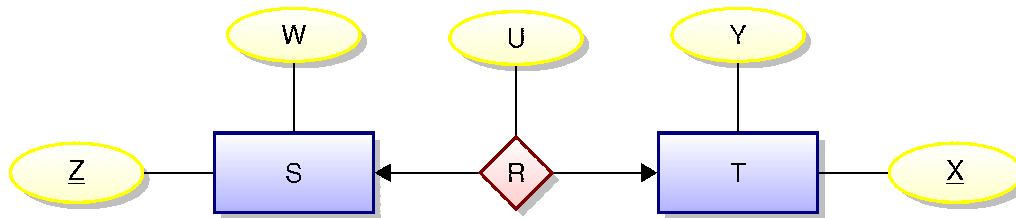
Relationships to Relational Schema



| Relation R is | Schema |
|--------------------------|--|
| many to many ($M : N$) | $R(\underline{X}, \underline{Z}, U)$ |
| one to many ($1 : M$) | $R(X, \underline{Z}, U)$ |
| many to one ($M : 1$) | $R(\underline{X}, Z, U)$ |
| one to one ($1 : 1$) | $R(\underline{X}, Z, U)$ or $R(X, \underline{Z}, U)$ |

Navigation icons: back, forward, search, etc.

“one to one” does not mean a “1-to-1 correspondence”



This database instance is OK

| S | | R | | | T | |
|----------------|----------------|----------------|----------------|----------------|----------------|----------------|
| Z | W | Z | X | U | X | Y |
| z ₁ | w ₁ | z ₁ | x ₂ | u ₁ | x ₁ | y ₁ |
| z ₂ | w ₂ | | | | x ₂ | y ₂ |
| z ₃ | w ₃ | | | | x ₃ | y ₃ |
| | | | | | x ₄ | y ₄ |

Navigation icons: back, forward, search, etc.

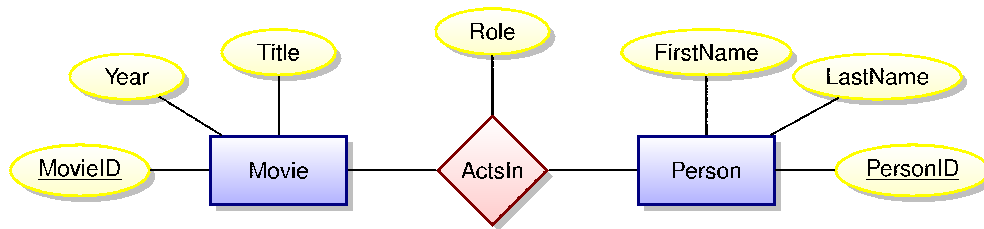
Some more real-world data ... (a slight change of SCOPE)

| Title | Year | Actor | Role |
|---|------|------------|------------------|
| Austin Powers: International Man of Mystery | 1997 | Mike Myers | Austin Powers |
| Austin Powers: International Man of Mystery | 1997 | Mike Myers | Dr. Evil |
| Austin Powers: The Spy Who Shagged Me | 1999 | Mike Myers | Austin Powers |
| Austin Powers: The Spy Who Shagged Me | 1999 | Mike Myers | Dr. Evil |
| Austin Powers: The Spy Who Shagged Me | 1999 | Mike Myers | Fat Bastard |
| Dude, Where's My Car? | 2000 | Bill Chott | Big Cult Guard 1 |
| Dude, Where's My Car? | 2000 | Marc Lynn | Cop with Whips |

How will this change our model?

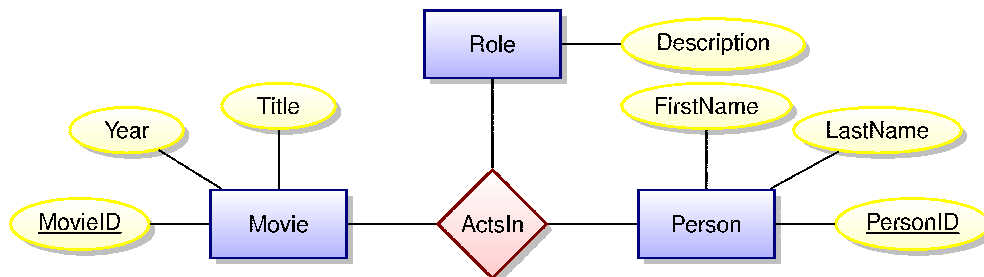
Navigation icons: back, forward, search, etc.

Will **ActsIn** remain a binary Relationship?



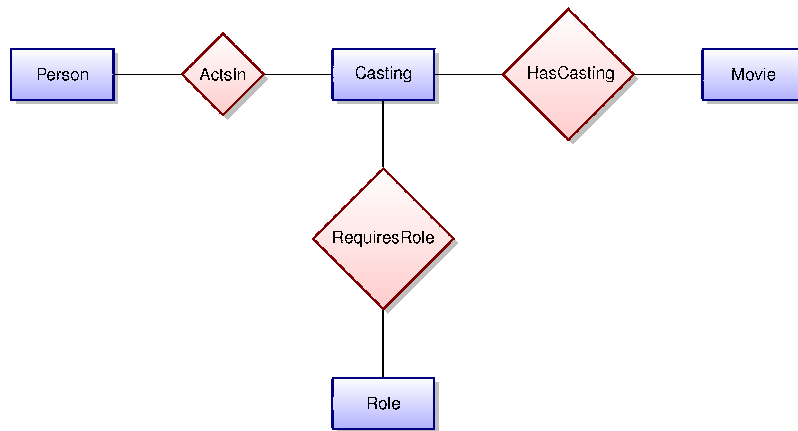
No! An actor can have many roles in the same movie!

Could **ActsIn** be modeled as a Ternary Relationship?



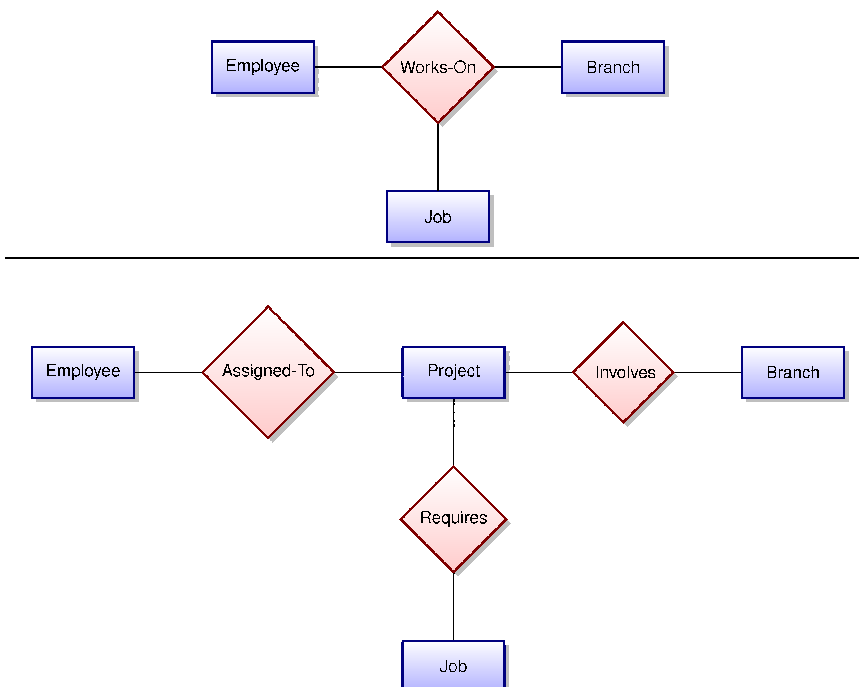
Yes, this works!

Can a ternary relationship be modeled with multiple binary relationships?

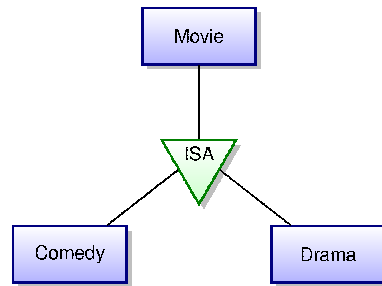


The **Casting** entity seems artificial. What attributes would it have?

Sometimes ternary to multiple binary makes more sense ...



Generalization



Questions

- Is every movie either comedy or a drama?
- Can a movie be a comedy and a drama?

But perhaps this isn't a good model ...

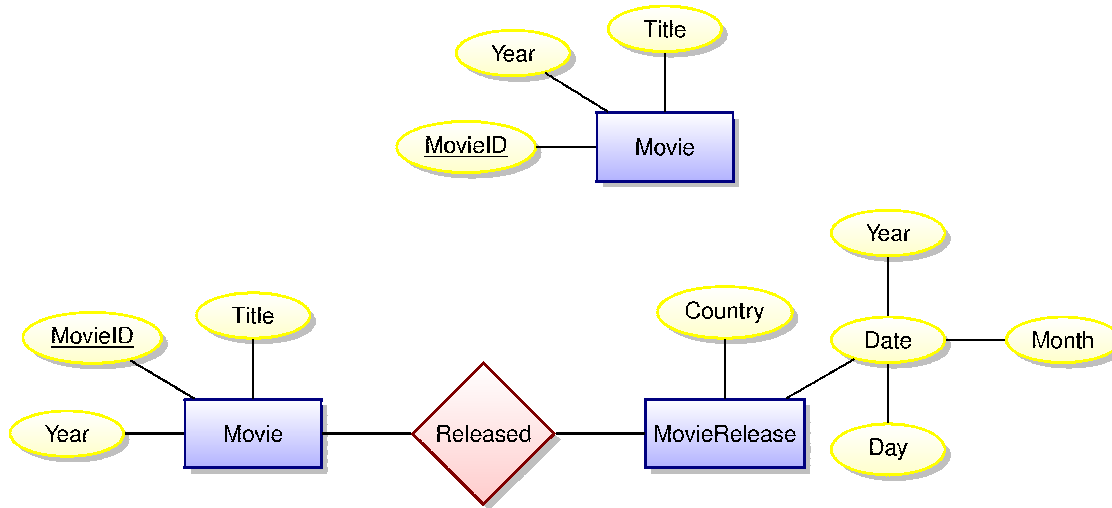
- What attributes would distinguish Drama and Comedy entities?
- What about **Science Fiction**?
- Perhaps **Genre** would make a nice entity, which could have a relationship with **Movie**.

Another change of SCOPE ...

Movies with detailed release dates

| Title | Country | Day | Month | Year |
|---|---------|-----|-------|------|
| Austin Powers: International Man of Mystery | USA | 02 | 05 | 1997 |
| Austin Powers: International Man of Mystery | Iceland | 24 | 10 | 1997 |
| Austin Powers: International Man of Mystery | UK | 05 | 09 | 1997 |
| Austin Powers: International Man of Mystery | Brazil | 13 | 02 | 1998 |
| Austin Powers: The Spy Who Shagged Me | USA | 08 | 06 | 1999 |
| Austin Powers: The Spy Who Shagged Me | Iceland | 02 | 07 | 1999 |
| Austin Powers: The Spy Who Shagged Me | UK | 30 | 07 | 1999 |
| Austin Powers: The Spy Who Shagged Me | Brazil | 08 | 10 | 1999 |
| Dude, Where's My Car? | USA | 10 | 12 | 2000 |
| Dude, Where's My Car? | Iceland | 9 | 02 | 2001 |
| Dude, Where's My Car? | UK | 9 | 02 | 2001 |
| Dude, Where's My Car? | Brazil | 9 | 03 | 2001 |
| Dude, Where's My Car? | Russia | 18 | 09 | 2001 |

... and an attribute becomes an entity with a connecting relation.



Is this really a good model?

- What is (natural) key of MovieRelease?
- Would a ternary relationship be better?

Question: What is the right model?

Answer: The question doesn't make sense!

- There is no “right” model ...
- It depends on the intended use of the database.
- What activity will the DBMS support?
- What data is needed to support that activity?

The issue of SCOPE is missing from most textbooks

- **Suppose** that all databases begin life with beautifully designed schemas.
- **Observe** that many operational databases are in a sorry state.
- **Conclude** that the **scope and goals** of a database continually change, and that **schema evolution** is a difficult problem to solve, in practice.