

# Examples

- For a monotonic function  $f : D \rightarrow C$  between posets, the free element over  $c \in C$  is the least element  $d \in D$  such that  $c \leq f(d)$ .
- $X^*$  is a free monoid over  $X$  wrt. the forgetful functor  $G : \mathbf{Mon} \rightarrow \mathbf{Sets}$ .
- The path category of a graph  $G$  is free over  $G$  wrt. the forgetful functor  $U : \mathbf{Cat} \rightarrow \mathbf{Graph}$ .

Make unit arrows explicit!

**Exercise:** What is the free poset over a set  $X$  wrt. the forgetful functor  $G : \mathbf{Pos} \rightarrow \mathbf{Sets}$ ?

**Exercise:** What is the free object over  $(A, B)$  wrt. the diagonal functor  $\Delta : \mathbf{C} \rightarrow \mathbf{C} \times \mathbf{C}$ ?

# Facts about free objects

**Fact:** For a functor  $G : \mathbf{D} \rightarrow \mathbf{C}$ , free objects over  $X \in |\mathbf{C}|$  are initial objects in the comma category  $K_X \downarrow G$  where  $K_X : \mathbf{1} \rightarrow \mathbf{C}$  is the functor constant at  $X$ .

**Corollary:** Free objects, if they exist, are unique up to isomorphism.

**Fact:** If  $A$  is free over  $X$  wrt.  $G : \mathbf{D} \rightarrow \mathbf{C}$  then for each  $B \in |\mathbf{D}|$  there is a bijection

$$(-)^{\#} : \text{hom}_{\mathbf{C}}(X, GB) \cong \text{hom}_{\mathbf{D}}(A, B)$$

# Free objects are functorial

Consider a functor  $G : \mathbf{D} \rightarrow \mathbf{C}$ .

If every  $C \in |\mathbf{C}|$  has a free object  $FC \in |\mathbf{D}|$  wrt.  $G$  then the mapping

$$\begin{array}{ccc} C & \mapsto & FC \\ f : C \rightarrow C' & \mapsto & (\eta_{C'} \circ f)^\# \end{array}$$

defines a **functor**  $F : \mathbf{C} \rightarrow \mathbf{D}$ .

Further,  $\eta : \text{Id}_{\mathbf{C}} \rightarrow GF$  is a natural transformation.

$$\begin{array}{ccc} \mathbf{C} & \xleftarrow{G} & \mathbf{D} \\ \\ \begin{array}{ccc} C & \xrightarrow{\eta_C} & GFC \\ \downarrow f & & \downarrow GFf \\ C' & \xrightarrow{\eta_{C'}} & GFC' \end{array} & & \begin{array}{ccc} FC & & \\ \downarrow Ff = (\eta_{C'} \circ f)^\# & & \\ FC' & & \end{array} \end{array}$$

# Left adjoints

**Defn.** A functor  $F : \mathbf{C} \rightarrow \mathbf{D}$  is **left adjoint to**  $G : \mathbf{D} \rightarrow \mathbf{C}$  **with unit**  $\eta : \text{Id}_{\mathbf{C}} \rightarrow GF$  if for every  $C \in |\mathbf{C}|$ ,  $FC$  with  $\eta_C$  is free over  $C$  wrt.  $G$ .

## Examples:

- the free monoid functor is left adjoint to  $G : \mathbf{Mon} \rightarrow \mathbf{Sets}$
- the path category functor is left adjoint to  $U : \mathbf{Cat} \rightarrow \mathbf{Graph}$
- “the” coproduct functor  $+$  :  $\mathbf{C} \times \mathbf{C} \rightarrow \mathbf{C}$  is left adjoint to the diagonal functor  $\Delta : \mathbf{C} \rightarrow \mathbf{C} \times \mathbf{C}$ .

**Exercise:** What is a left adjoint to a monotonic function between posets?

# Facts about left adjoints

**Theorem:** Left adjoints to any fixed  $G$ , if they exist, are unique up to **natural** isomorphism.

**Theorem:** If  $F$  is left adjoint to  $G$ , then:

- $F$  preserves colimits,
- $G$  preserves limits.

**Theorem:** Let  $\mathbf{D}$  be locally small & complete (ie. have all limits).

A functor  $G : \mathbf{D} \rightarrow \mathbf{C}$  has a left adjoint if and only if:

- $G$  preserves limits,
- for every  $C \in |\mathbf{C}|$  there exists a **set** (ie not a proper class)

$\{f_i : C \rightarrow GD_i \mid i \in \mathcal{I}\}$  of arrows such that

for each  $D \in \mathbf{D}$  and  $f : C \rightarrow GD$ ,

there exist  $i \in \mathcal{I}$  and  $g : D_i \rightarrow D$  such that  $f = Gg \circ f_i$ .