### Lecture 4

Scott Induction

# Scott's Fixed Point Induction Principle

Let  $f:D \to D$  be a continuous function on a domain D.

fixed point of f is in S, *i.e.* that For any admissible subset  $S\subseteq D$ , to prove that the least

$$fx(f) \in S$$
,

it suffices to prove

$$\forall d \in D \ (d \in S \Rightarrow f(d) \in S)$$
.

# Chain-closed and admissible subsets

for all chains  $d_0 \sqsubseteq d_1 \sqsubseteq d_2 \sqsubseteq \ldots$  in DLet D be a cpo. A subset  $S\subseteq D$  is called chain-closed iff

$$(\forall n \ge 0 \cdot d_n \in S) \Rightarrow \left(\bigsqcup_{n \ge 0} d_n\right) \in S$$

chain-closed subset of D and  $\bot \in S$ . If D is a domain,  $S\subseteq D$  is called admissible iff it is a

(resp. admissible) subset of D. A property  $\Phi(d)$  of elements  $d \in D$  is called *chain-closed* (resp. admissible) iff  $\{d\in D\mid \Phi(d)\}$  is a chain-closed

## **Building chain-closed subsets (I)**

Let D, E be cpos.

#### **Basic relations:**

• For every  $d \in D$ , the subset

$$\downarrow\!(d) \stackrel{\mathrm{def}}{=} \{ \, x \in D \mid x \sqsubseteq d \, \}$$

of *D* is chain-closed.

The subsets

$$\{(x,y) \in D \times D \mid x \sqsubseteq y\}$$
$$\{(x,y) \in D \times D \mid x = y\}$$

of  $D \times D$  are chain-closed.

# Example (I): Least pre-fixed point property

Let D be a domain and let f:D o D be a continuous function.

$$\forall d \in D. f(d) \sqsubseteq d \implies fx(f) \sqsubseteq d$$

Proof by Scott induction.

Let  $d \in D$  be a pre-fixed point of f. Then,

$$x \in \downarrow(d) \implies x \sqsubseteq d$$

$$\implies f(x) \sqsubseteq f(d)$$

$$\implies f(x) \sqsubseteq d$$

$$\implies f(x) \subseteq d$$

Hence,

$$fx(f) \in \downarrow(d)$$
.

## Building chain-closed subsets (II)

#### Inverse image:

Let  $f:D \to E$  be a continuous function.

If S is a chain-closed subset of E then the inverse image

$$f^{-1}S = \{x \in D \mid f(x) \in S\}$$

is an chain-closed subset of D.

#### Example (II)

functions such that  $f \circ g \sqsubseteq g \circ f$ . Then, Let D be a domain and let  $f,g:D\to D$  be continuous

$$f(\bot) \sqsubseteq g(\bot) \implies fx(f) \sqsubseteq fx(g)$$
.

Proof by Scott induction.

Consider the admissible property  $\Phi(x) \equiv (f(x) \sqsubseteq g(x))$ 

Since

$$f(x) \sqsubseteq g(x) \Rightarrow g(f(x)) \sqsubseteq g(g(x)) \Rightarrow f(g(x)) \sqsubseteq g(g(x))$$

we have that

$$f(fx(g)) \sqsubseteq g(fx(g))$$
.

## Building chain-closed subsets (III)

### Logical operations:

If  $S,T\subseteq D$  are chain-closed subsets of D then

 $S \cup T$  and  $S \cap S$ 

are chain-closed subsets of D.

- If  $\{S_i\}_{i\in I}$  is a family of chain-closed subsets of Dsubset of *D*. indexed by a set I, then  $\bigcap_{i\in I} S_i$  is a chain-closed
- If a property P(x,y) determines a chain-closed subset of a chain-closed subset of E $D \times E$ , then the property  $\forall x \in D.\ P(x,y)$  determines

### **Example (III): Partial correctness**

Let  $\mathcal{F}: State \longrightarrow State$  be the denotation of

while 
$$X > 0$$
 do  $(Y := X * Y; X := X - 1)$ .

For all 
$$x,y\geq 0$$
, 
$$\mathcal{F}[X\mapsto x,Y\mapsto y]\downarrow \\ \Rightarrow \mathcal{F}[X\mapsto x,Y\mapsto y] = [X\mapsto 0,Y\mapsto!x\cdot y].$$

#### Recall that

$$\mathcal{F}=fx(f)$$
 where  $f:(State \to State \to State) \to (State \to State)$  is given by 
$$f(w)=\lambda(x,y) \in State. \left\{ \begin{array}{ll} (x,y) & \text{if } x \leq 0 \\ w(x-1,x\cdot y) & \text{if } x>0 \end{array} \right.$$

### Proof by Scott induction.

We consider the admissible subset of (State 
ightharpoonup State) given by

$$S = \begin{cases} w & \forall x, y \ge 0. \\ w[X \mapsto x, Y \mapsto y] \downarrow \\ \Rightarrow w[X \mapsto x, Y \mapsto y] = [X \mapsto 0, Y \mapsto !x \cdot y] \end{cases}$$

and show that

$$w \in S \implies f(w) \in S$$
.