

# The Ontology of Naproche

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## Abstract

This document provides an overview of the object-level ontology of Naproche. Moreover, it introduces common (symbolic and verbal) macros for the notions provided by this ontology which extend the core language of ForTheL.

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## 1 Introduction

This document describes the notions and axioms hard-coded in Naproche. Furthermore, it presents some additional symbolic and verbal macros that extend the core language of ForTheL. These macros are provided by the library `libraries/meta` which contains a module `preliminaries.ftl` that exports all

these macros. Every library module (excluding those in the library `meta`) is expected to import the module `preliminaries.ftl`. So by importing any library module to your formalization, you automatically have access to all macros described in this document.

## 2 Propositions

ForTheL is essentially a surface language for (unsorted) first-order logic. We therefore provide some symbolic notations for logical connectives and quantifiers that are supported by ForTheL:

**Convention 1.** Let  $\top$  stand for truth.

**Convention 2.** Let  $\perp$  stand for falsity.

**Convention 3.** Let  $\neg \varphi$  stand for not  $\varphi$ .

**Convention 4.** Let  $\varphi \wedge \psi$  stand for  $\varphi$  and  $\psi$ .

**Convention 5.** Let  $\varphi \vee \psi$  stand for  $\varphi$  or  $\psi$ .

**Convention 6.** Let  $\varphi \implies \psi$  stand for  $\varphi$  implies  $\psi$ .

**Convention 7.** Let  $\varphi \iff \psi$  stand for  $\varphi$  iff  $\psi$ .

**Convention 8.** Let  $\forall x R y : \varphi$  stand for for all  $x R y$  we have  $\varphi$ .

**Convention 9.** Let  $\exists x R y : \varphi$  stand for there exists an  $x R y$  such that  $\varphi$ .

## 3 Notions

### 3.1 Objects

*Objects* are mathematical entities that are “small” enough to be contained in classes.

**Signature 10.** A *mathematical object* is a notion. Let an *object* stand for a mathematical object. Let an *element* stand for a mathematical object.

### 3.2 Ordered Pairs

**Signature 11.** Let  $a, b$  be objects.  $(a, b)$  is an object.

**Convention 12.** Let the *ordered pair of  $a$  and  $b$*  stand for  $(a, b)$ . Let an *ordered pair* stand for an object  $x$  such that  $x = (a, b)$  for some objects  $a, b$ .

### 3.3 Classes and Sets

**Signature 13.** A *class* is a notion. Let a *collection* stand for a class.

Sets are considered to be classes that are “small” enough to be contained in other classes:

**Definition 14.** A *set* is a class that is an object.

**Signature 15.** Let  $A$  be a class. A *element of  $A$*  is an object. Let  $a \in A$  stand for  $a$  is an element of  $A$ . Let  $a \notin A$  stand for  $a$  is not an element of  $A$ .

**Convention 16.** Let a *member of  $X$*  stand for an element of  $X$ . Let  $x$  *belongs to  $X$*  stand for  $x$  is an element of  $X$ . Let  $X$  *contains  $x$*  stand for  $x$  is an element of  $X$ . Let  $x$  is *contained in  $X$*  stand for  $x$  is an element of  $X$ . Let  $x$  *lies in  $X$*  stand for  $x$  is an element of  $X$ . Let  $x$  is *in  $X$*  stand for  $x$  is an element of  $X$ .

### 3.4 Comprehension Terms

In ForTheL we have three types of comprehension terms:

- “Enumeration classes”, i.e. terms of the form  $\{a_1, \dots, a_n\}$ , that represent the class that contains precisely the objects  $a_1, \dots, a_n$ . Typical examples are  $\{4, 7, 11\}$  (i.e. the class consisting of the numbers 4, 7 and 11) or  $\{x, y + 1, z - x\}$  (i.e. the class consisting of the elements  $x$ ,  $y + 1$  and  $z - x$  for some fixed objects  $x$ ,  $y$  and  $z$ ).
- “Comprehension classes”, i.e. terms of the form  $\{\tau \mid \varphi\}$ , that represent the class of all objects of the form  $\tau$  such that the components of  $\tau$  satisfy the condition  $\varphi$ . Typical examples are  $\{p \mid p \text{ is an odd prime number}\}$  (i.e the class of all odd prime numbers) or  $\{(p, n) \mid n \in \mathbb{N} \text{ and } p \text{ is a prime divisor of } n\}$  (i.e. the class of all pairs  $(p, n)$  of natural numbers such that  $p$  is a prime divisor of  $n$ ).
- “Separation classes”, i.e. terms of the form  $\{\tau \in A \mid \varphi\}$ , that represent the subclass of a class  $A$  of all objects of the form  $\tau$  such that the components of  $\tau$  satisfy the condition  $\varphi$ . Typical examples are  $\{p \in \mathbb{P} \mid p \text{ is odd}\}$  (i.e the class of all odd prime numbers) or  $\{(p, n) \in \mathbb{N} \times \mathbb{N} \mid p \text{ is a prime divisor of } n\}$  (i.e. the class of all pairs  $(p, n)$  of natural numbers such that  $p$  is a prime divisor of  $n$ ).

**Signature 17.** Let  $a_1, \dots, a_n$  be objects.  $\{a_1, \dots, a_n\}$  is a class.

**Signature 18.** Let  $\tau$  be a term and  $\varphi$  be a formula.  $\{\tau \mid \varphi\}$  is a class.

**Signature 19.** Let  $\tau$  be a term,  $A$  be a class and  $\varphi$  be a formula.  $\{\tau \in A \mid \varphi\}$  is a class.

### 3.5 Maps and Functions

**Signature 20.** A *map* is a notion.

Similar to the notion of sets which are considered to be “small” classes, we consider *functions* to be “small” maps.

**Definition 21.** A *function* is a map that is an object.

**Signature 22.** Let  $f$  be a map.  $\text{dom}(f)$  is a class. Let the *domain of  $f$*  stand for  $\text{dom}(f)$ .

**Convention 23.** Let  $f$  is *defined on  $X$*  stand for  $\text{dom}(f) = X$ .

**Signature 24.** Let  $f$  be a map and  $a \in \text{dom}(f)$ .  $f(a)$  is an object.

**Convention 25.** Let the *value of  $f$  at  $x$*  stand for  $f(x)$ . Let  $f(x, y)$  stand for  $f((x, y))$ .

### 3.6 $\lambda$ -Terms

In low-level map definitions, ForTheL allows to define maps via  $\lambda$ -terms:

**Signature 26.** Let  $v$  be a variable,  $A$  be a class and  $\tau$  be a term.  $\lambda v \in A. \tau$  is a map.

## 4 Relations

**Convention 27.** Let  $x = y$  stand for  $x$  is equal to  $y$ .

**Convention 28.** Let  $x \neq y$  stand for  $x$  is not equal to  $y$ .

**Convention 29.** Let  $x$  and  $y$  *agree* stand for  $x = y$ . Let  $x$  *agrees with  $y$*  stand for  $x = y$ .

**Convention 30.** Let  $x$  and  $y$  are *distinct* stand for  $x \neq y$ .

Naproche has a hard-coded induction proof tactic (cf. section 5.4) that relies on an “induction ordering”  $\prec$ . An important special case of this tactic is the following induction principle: If, for all  $x$ , we can show that a predicate  $\varphi[y]$  holds for every  $y \prec x$ , then  $\varphi[x]$  holds for every  $x$ . Typical instances of this principle are the (strong) induction principle from natural number arithmetics (with  $\prec$  being the  $<$ -relation) or the  $\in$ -induction principle from set theory (with  $\prec$  being the  $\in$ -relation).

**Convention 31.** Let  $x \prec y$  stand for  $x$  is inductively less than  $y$ .

## 5 Axioms

### 5.1 Extensionality Axioms

To provide equality conditions for ordered pairs, classes and maps, the following extensionality axioms are hard-coded in Naproche.

**Axiom 32.** Let  $a, a', b, b'$  be objects. If  $(a, b) = (a', b')$  then  $a = a'$  and  $b = b'$ .

**Axiom 33.** Let  $A, B$  be classes. If every element of  $A$  is an element of  $B$  and every element of  $B$  is an element of  $A$  then  $A = B$ .

**Axiom 34.** Let  $f, g$  be maps. If  $\text{dom}(f) = \text{dom}(g)$  and  $f(a) = g(a)$  for all  $a \in \text{dom}(f)$  then  $f = g$ .

### 5.2 Comprehension Term Axioms

The following axioms specify the elements of comprehension terms (cf. section 3.4).

**Axiom 35.** Let  $\tau$  be a term,  $A$  be a class and  $\varphi$  be a formula. If  $A$  is a set then  $\{\tau \in A \mid \varphi\}$  is a set.

**Corollary 36.** Let  $A$  be a set and  $B$  be a class. Assume that every element of  $B$  is an element of  $A$ . Then  $B$  is a set.

*Proof.* Define  $C := \{x \in A \mid x \in B\}$ . Then  $C$  is a set and  $C = B$ . ■

**Axiom 37.** Let  $a_1, \dots, a_n$  be objects. Then  $x \in \{a_1, \dots, a_n\}$  iff  $x = a_1$  or  $\dots$  or  $x = a_n$ .

**Axiom 38.** Let  $v_1, \dots, v_n$  be variables,  $\tau[v_1, \dots, v_n]$  be a term and  $\varphi[v_1, \dots, v_n]$  be a formula. Then  $x \in \{\tau[v_1, \dots, v_n] \mid \varphi[v_1, \dots, v_n]\}$  iff  $x$  is an object and there exist  $y_1, \dots, y_n$  such that  $\varphi[y_1, \dots, y_n]$  and  $x = \tau[y_1, \dots, y_n]$ .

**Corollary 39.** Let  $v$  be a variable and  $\varphi[v]$  be a formula. Then  $x \in \{v \mid \varphi[v]\}$  iff  $x$  is an object and  $\varphi[x]$ .

### 5.3 $\lambda$ -Term Axioms

The following axioms specify the domain and values of  $\lambda$ -terms:

**Axiom 40.** Let  $v$  be a variable,  $A$  be a class and  $\tau[v]$  be a term. Then  $\text{dom}(\lambda v \in A. \tau) = A$ .

**Axiom 41.** Let  $v$  be a variable,  $A$  be a class and  $\tau[v]$  be a term. Then  $(\lambda v \in A. \tau)(x) = \tau[x]$  for all  $x \in A$ .

### 5.4 Miscellaneous Axioms

Subclasses of sets that are formed via separation class terms (cf. 3.4) are again sets:

**Axiom 42.** Let  $v_1, \dots, v_n$  be variables,  $\tau[v_1, \dots, v_n]$  be a term,  $A$  be a class and  $\varphi[v_1, \dots, v_n]$  be a formula. Then  $x \in \{\tau[v_1, \dots, v_n] \in A \mid \varphi[v_1, \dots, v_n]\}$  iff  $x \in A$  and there exist  $y_1, \dots, y_n$  such that  $\varphi[y_1, \dots, y_n]$  and  $x = \tau[y_1, \dots, y_n]$ .

**Corollary 43.** Let  $v$  be a variable,  $A$  be a class and  $\varphi[v]$  be a formula. Then  $x \in \{v \in A \mid \varphi[v]\}$  iff  $x \in A$  and  $\varphi[x]$ .

As noted in section 3.5, functions are considered to be “small” maps. Whether a map is small or not is determined by the “size” of its domain:

**Axiom 44.** Let  $f$  be a map.  $f$  is a function iff  $\text{dom}(f)$  is a set.

To facilitate proofs by induction, Naproche provides a hard-coded induction proof tactic which corresponds to the following induction principle.

**Axiom 45.** Let  $v_1, \dots, v_n$  be variables,  $\tau[v_1, \dots, v_n]$  be a term and  $\psi_1[v_1]$ ,  $\psi_2[v_1, v_2]$ ,  $\dots$ ,  $\psi_n[v_1, \dots, v_n]$  and  $\varphi[v_1, \dots, v_n]$  be formulas. Assume that the following holds for all  $x_1, \dots, x_n$  with  $\psi_i[x_1, \dots, x_i]$  for all  $i \in \{1, \dots, n\}$ : If  $\tau[y_1, \dots, y_n] \prec \tau[x_1, \dots, x_n]$  for all  $y_1, \dots, y_n$  with  $\psi_i[y_1, \dots, y_i]$  for all  $i \in \{1, \dots, n\}$  then  $\varphi[x_1, \dots, x_n]$ . Then  $\varphi[x_1, \dots, x_n]$  for all  $x_1, \dots, x_n$  with  $\psi_i[x_1, \dots, x_i]$  for all  $i \in \{1, \dots, n\}$ .

**Corollary 46.** Let  $A$  be a class,  $v$  be a variable and  $\varphi[v]$  be a formula. Assume that the following holds for all  $x \in A$ : If  $\varphi[y]$  for for all  $y \in A$  with  $y \prec x$  then  $\varphi[x]$ . Then  $\varphi[x]$  for all  $x \in A$ .

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