## Implications of Rent's Rule for Network-on-Chip Design and Its Fault-Tolerance

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#### **Introduction: NoC Evaluation**

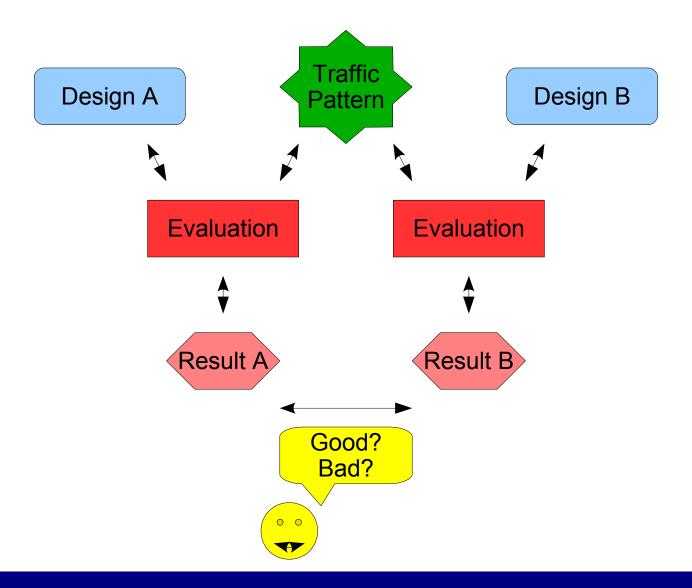
Design A

Design B

#### **Introduction: NoC Evaluation**



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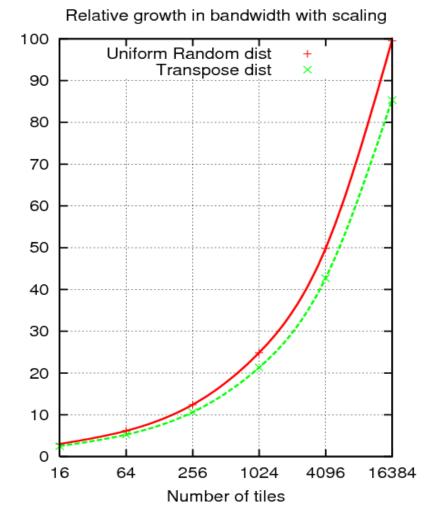


### Introduction

- When we build a new Network-on-Chip
   Is it better than existing networks on chip?
- Traffic models are used to evaluate and compare Networks-on-Chip
  - What does traffic look like?
  - Designing for the future, so what is its scaling behaviour?
  - Commonly use very simple non-local models. Is this wise?

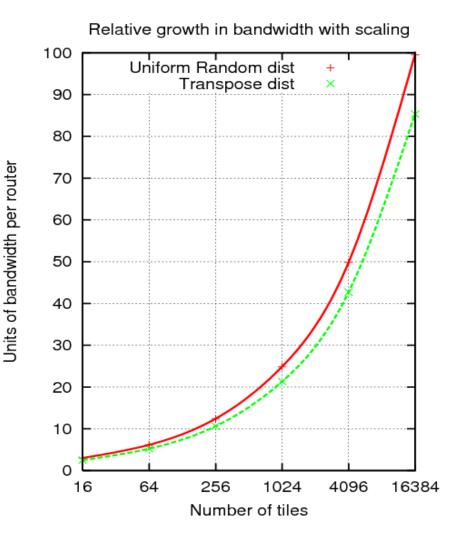
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  - Commonly use very simple non-local models.
    Is this wise?
- How important is locality and how can we model it?

# Scaling of per-tile BW



- Each tile generates one unit of BW
- BW through/routed by each tile grows exponentially with technology scaling
- Doubling per quadrupling of tiles
- # Metal layers, repeaters, power grows exponentially

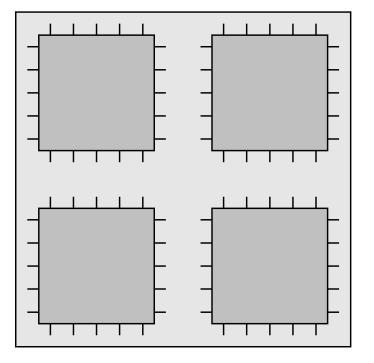
# Scaling of per-tile BW



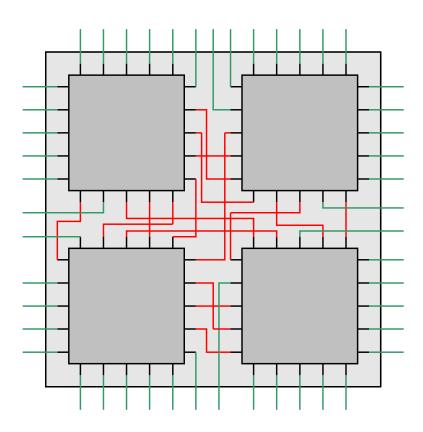
- Each tile generates one unit of BW
- BW through/routed by each tile grows exponentially with technology scaling
- Doubling per quadrupling of tiles
- # Metal layers, repeaters, power grows exponentially <u>for</u> <u>all on-chip topologies</u>

### **Classical Rent's Rule**





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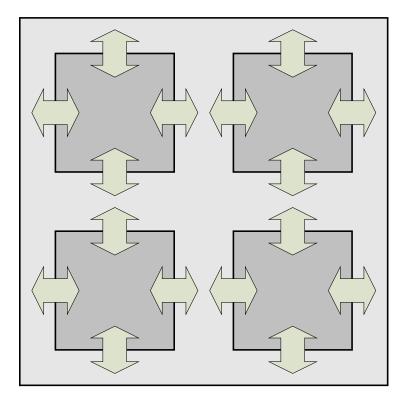
- A description of locality
- Rent's exponent characterises the split between internal and external nets

### **Reasons for Placement Locality**

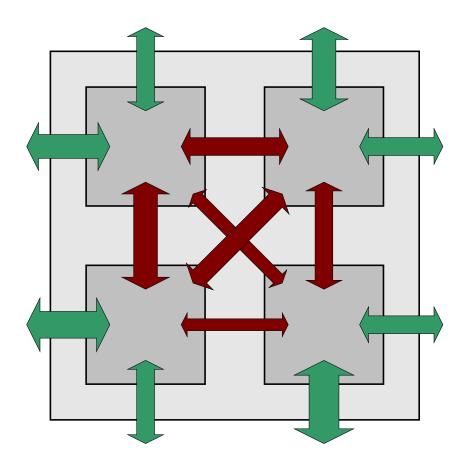
Domain to minimise	Wires	NoC
Delay	Wire delay	NoC latency
		(&  congestion)
Congestion	Wire-density	Cross-sectional
		bandwidth
Power	Wire buffering &	Hop-length & router
	$\operatorname{length}$	utilisation

### **Bandwidth Version**

Bandwidth locality



## **Bandwidth Version**

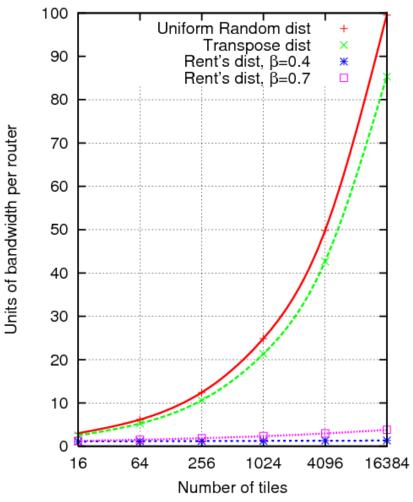


- Bandwidth locality
- Rent's bandwidthexponent characterises split between internal and external traffic

# Scaling of per-tile BW (comparison)

 Rent's-type locality allows graceful increase in per-tile bandwidth

Relative growth in bandwidth with scaling

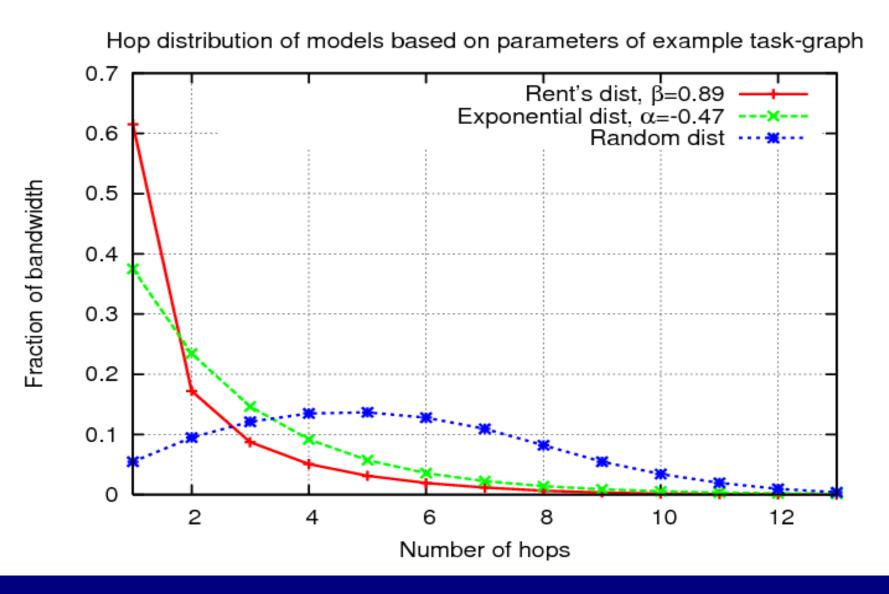


# Scaling of per-tile BW (Rent's)

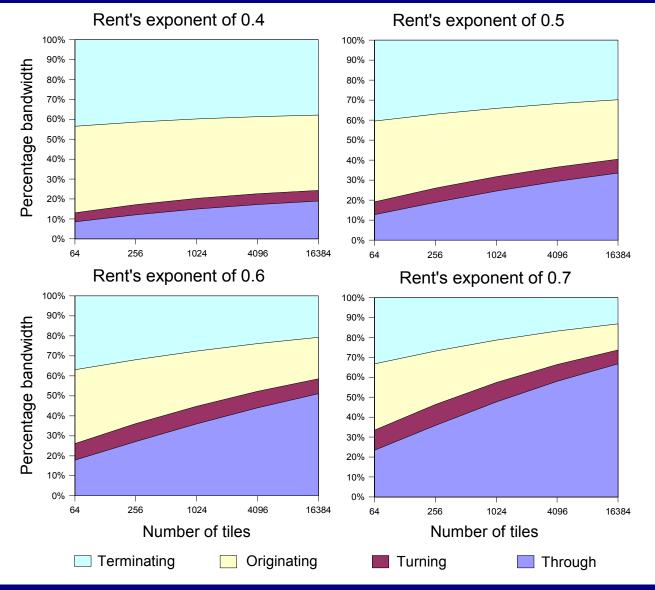
- Many-fold increase with scaling for high versus low Rent's bandwidth-exponent
- This is in line with ordinary VLSI expectations

Relative growth in bandwidth with scaling 4 Rent's dist, β=0.4 Rent's dist, β=0.5 Rent's dist, β=0.6 3.5 Rent's dist. B=0.7 З Units of bandwidth per router 2.5 2 1.5 0.5 0 16 64 256 1024 4096 16384 Number of tiles

## **Example Hop-length Distribution (8x8)**

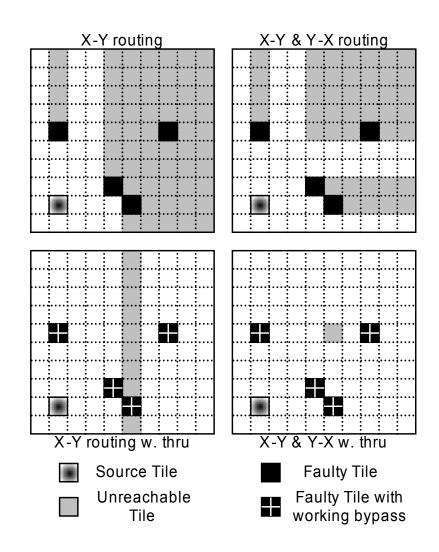


# **Traffic Types**

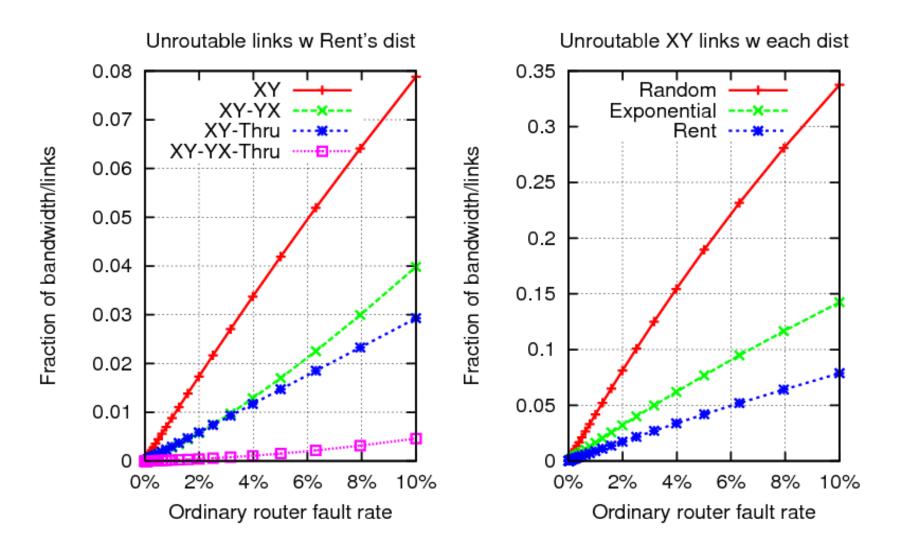


#### **Case-study: A Fault Tolerance Approach**

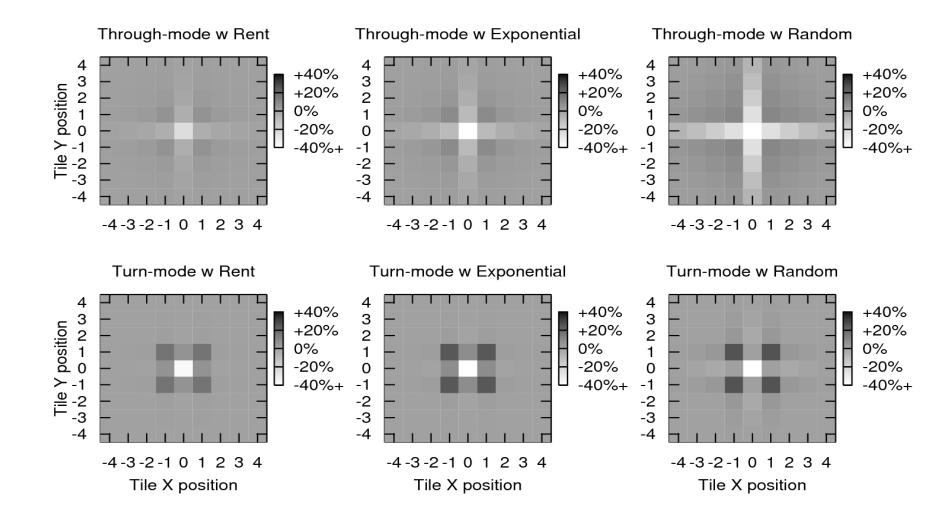
- Add bypass of faulty router logic
- Effectively a pipelined interconnect
- Implicit creditscheme preserves flow-control



## **Reachability Analysis**



## **Congestion Surrounding Faulty Tile**



## Conclusions

- Locality is essential. Non-local traffic models lead to unacceptable scaling behaviour
- Reasons for locality in software-circuits are analogous to ones in hardware-circuits
- A bandwidth version of Rent's rule for NoC can be derived to describe locality in software circuits
  - Distinct from classical Rent's rule
  - Hop-length distribution versus wire-length distribution
- Can analyse router utilisation according to traffic type
- In fault-tolerance analysis:
  - The traffic model used is important
  - Different traffic models lead to different results

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