

# Cylinder Diagrams: Visualising Multi-parameter Natural Transformations

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European Union  
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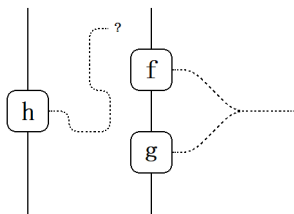
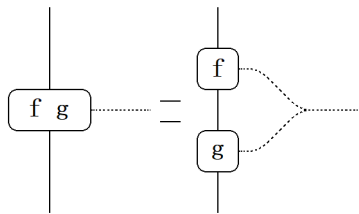
Things I like:

- ▶ Functional programs and composition.
- ▶ Effects like global store and nondeterminism.
- ▶ Parallelism and concurrency.

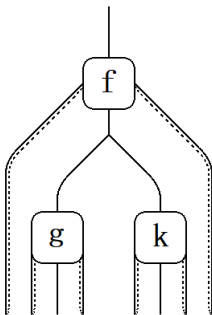
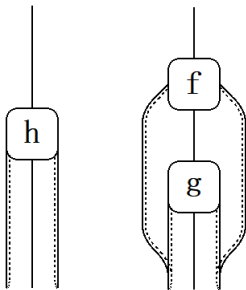
Dimensions of a page:

- ▶ Height.
- ▶ Width.
- ▶ ...

## Effectful composition



## How to add parallelism?



# Motivation

Possible solutions:

- ▶ Surface diagrams, sheet diagrams.
- ▶ Ribbon diagrams, tape diagrams.
- ▶ String diagrams with special operations.

Solution of this talk: Cylinder diagrams

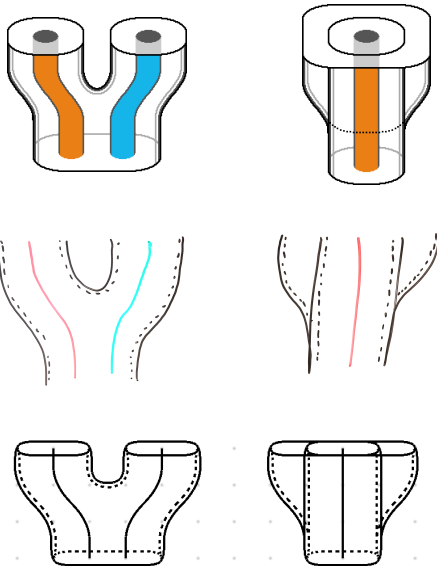
Inflate the strings and allow them to contain diagrams.

- ▶ Allows for arbitrary nesting.
- ▶ Can be neatly flattened to 2 dimensions.

Methods for drawing:

- ▶ This presentation will use drawing assembled using tiles with premade illustrations.
- ▶ Can also be illustrated by 2-dimensional string diagrams using *bracket-strings*.

# Illustration: Parallel and sequential composition



# Core Categorical structure

Given a monoidal category  $\mathcal{C}$ ,  
we describe the category consisting of:

- ▶ As objects multi-endofunctors on  $\mathcal{C}$ .
- ▶ As morphisms natural transformations between them.

The resulting category is also monoidal.

It contains  $\mathcal{C}$  as a full subcategory of constant functors.

Main example is taking the category of sets  $\mathcal{C} = \mathbf{Set}$ ,  
with the Cartesian product as monoidal product.

# Objects and morphisms of $\mathcal{C}$

Objects are represented (solid) cylinders, and functions by blocks between them:

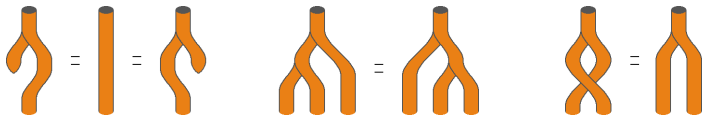
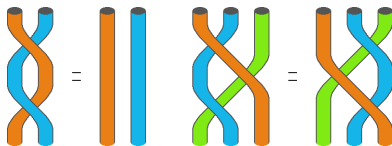
$$\text{Object: } A = \text{cylinder} \qquad \text{Morphism: } \left( f : \frac{A}{B} \right) = \text{block with } f \text{ between two cylinders}$$

Special morphisms which exist in **Set** ( $\times$  is the product):

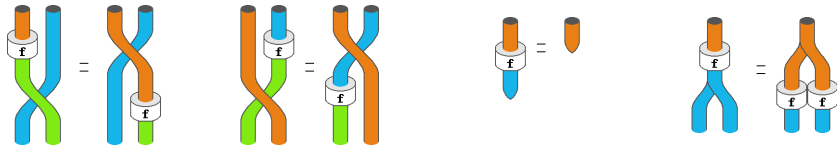
$$\left( s : \frac{X \times Y}{Y \times X} \right) = \text{crossing} \qquad c : \frac{X}{X \times X} = \text{join} \qquad D : \frac{X}{1} = \text{point}$$

# Objects and morphisms of $\mathcal{C}$

Special equations:



Naturality:





# Endofunctors of $\mathcal{C}$

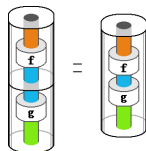
Functors are represented by hollow cylinders, within which lies the object or morphism to which it is applied.

$$A = \text{orange cylinder} \mapsto TA = \text{orange cylinder inside a hollow cylinder}$$

$$\left(f : \frac{A}{B}\right) = \text{orange cylinder on blue cylinder with label } f \mapsto \left(Tf : \frac{TA}{TB}\right) = \text{orange cylinder inside cylinder with } Tf \text{ label on blue cylinder} = \text{orange cylinder inside cylinder with } f \text{ label on blue cylinder}$$


Functor law for morphism composition:

$$\left(f : \frac{X}{Y}, g : \frac{Y}{Z}\right) \mapsto \left(Tg \circ Tf = T(g \circ f) : \frac{TX}{TZ}\right)$$



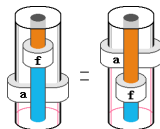
# Natural Transformations of $\mathcal{C}$

Natural transformations, a parametrized morphisms:

$$\left( \forall X. a_X : \frac{TX}{RX} \right) =$$


The object  $X$  goes through the natural transformation block, illustrating naturality:



$$\left( \forall f : \frac{X}{Y}. a_Y \circ Tf = Rf \circ a_X \right)$$



# Let there be monads

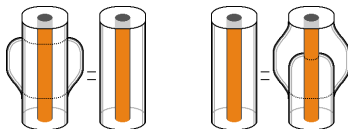
## Definition

A *monad* is a functor  $T$  with two natural transformations

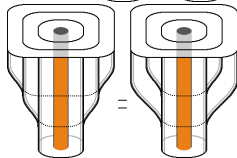
$(\eta_X : \frac{X}{TX}) =$   and  $(\mu_X : \frac{TTX}{TX}) =$  , satisfying the

following two equations:

(Left and right-unit)

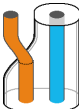


(Associativity)

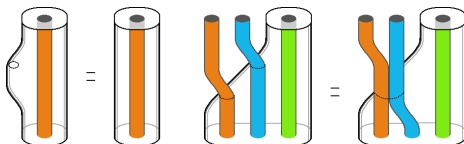


# Strength

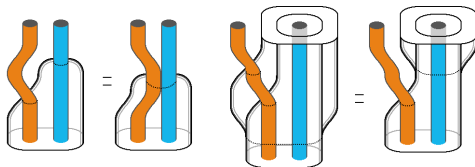
Strength is illustrated by pulling in the parameter:

$$\left( \sigma_{X,Y} : \frac{X \times TY}{T(X \times Y)} \right) =$$


Monoidal strength equations:

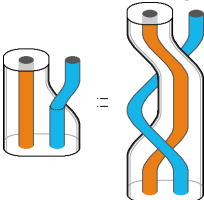


Strong monad equations:



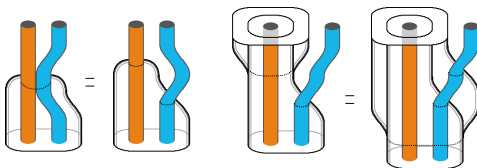
# Costrength

Using swap, we define costrength as

$$\left( \tau_{X,Y} : \frac{TX \times Y}{T(X \times Y)} \right) = \text{diagram} = \text{diagram}.$$


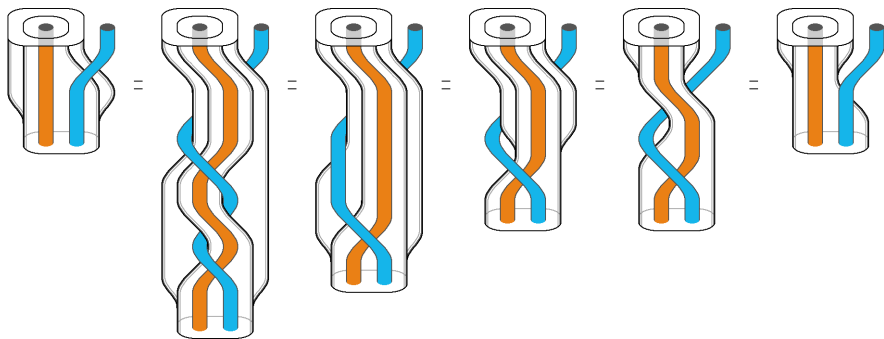
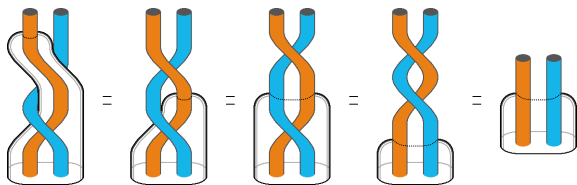
## Lemma

*A strong monad is a costrong monad.*



# Costrength

Proof.



# Commutativity

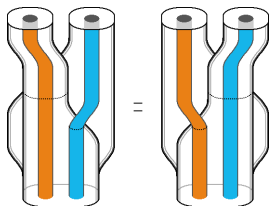
How do we *merge* monads?

One method is to use strength and costrength.

Commutativity says the order doesn't matter.

## Definition

A monad is commutative if the following equation holds:



# Monoidal Monad

## Definition

$T$  is a *monoidal monad* if there is a natural transformation

$$\left(m : \frac{TX \times TY}{T(X \times Y)}\right) = \text{diagram}$$

The diagram shows a container with two inputs at the top. An orange strand enters from the left and a blue strand enters from the right. They both curve downwards and then merge into a single output at the bottom.

equations:

$$(\eta) \text{ diagram} = \text{diagram}$$

The first diagram shows a container with two inputs at the top. An orange strand enters from the left and a blue strand enters from the right. They both curve downwards and then merge into a single output at the bottom. The second diagram shows a container with two inputs at the top. An orange strand enters from the left and a blue strand enters from the right. They both go straight down and then merge into a single output at the bottom.

$$(\mu) \text{ diagram} = \text{diagram}$$

The first diagram shows a container with two inputs at the top. An orange strand enters from the left and a blue strand enters from the right. They both curve downwards and then merge into a single output at the bottom. The second diagram shows a container with two inputs at the top. An orange strand enters from the left and a blue strand enters from the right. They both go straight down and then merge into a single output at the bottom.

## Lemma



*Any strong commutative monad is a monoidal monad.*



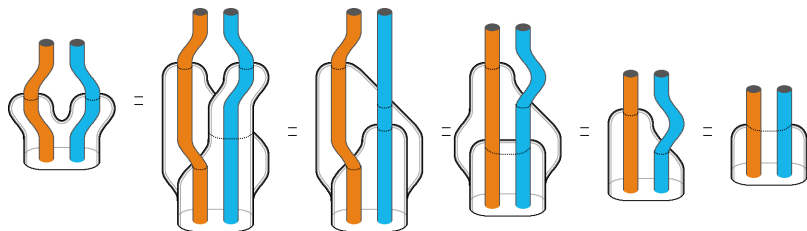
# Monoidal Monad

Proof.

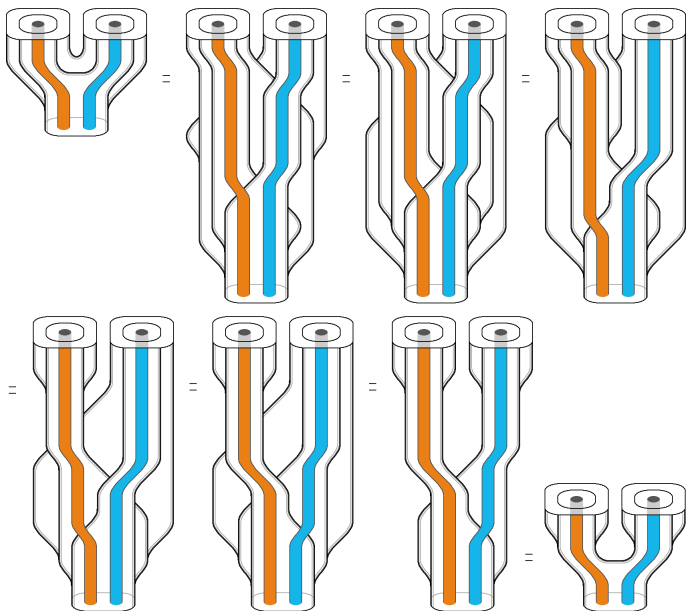
We define the following natural transformation:

$$\left(m : \frac{TX \times TY}{T(X \times Y)}\right) = \text{[Diagram 1]} = \text{[Diagram 2]}.$$


Using commutativity, we can prove both monoidal monad properties:



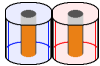
# Monoidal Monad



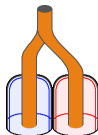
# Product of Monads

$T$  and  $R$  both monads.

Assume we have copy and delete.

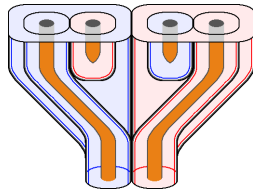
$$(T \times R)(X) = (TX \times RX) =$$


► The monad unit  $\left( \eta_X^{T \times R} : \frac{X}{TX \times RX} \right) =$



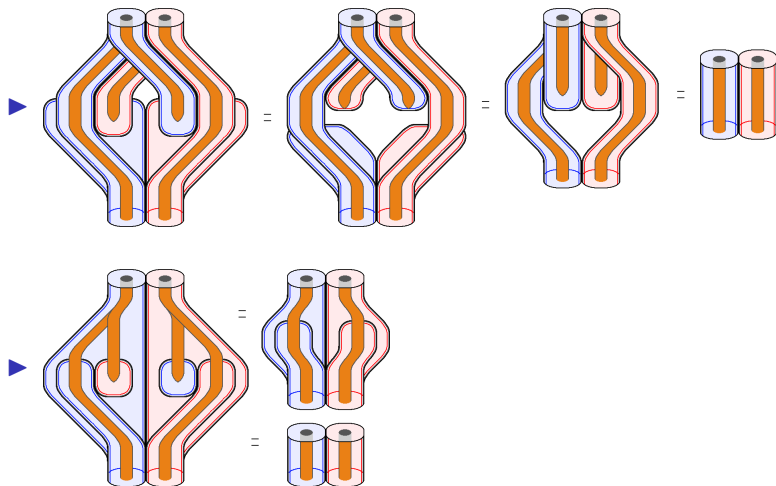
► The monad multiplication

$$\left( \mu_X^{T \times R} : \frac{T(TX \times RX) \times R(TX \times RX)}{TX \times RX} \right) =$$

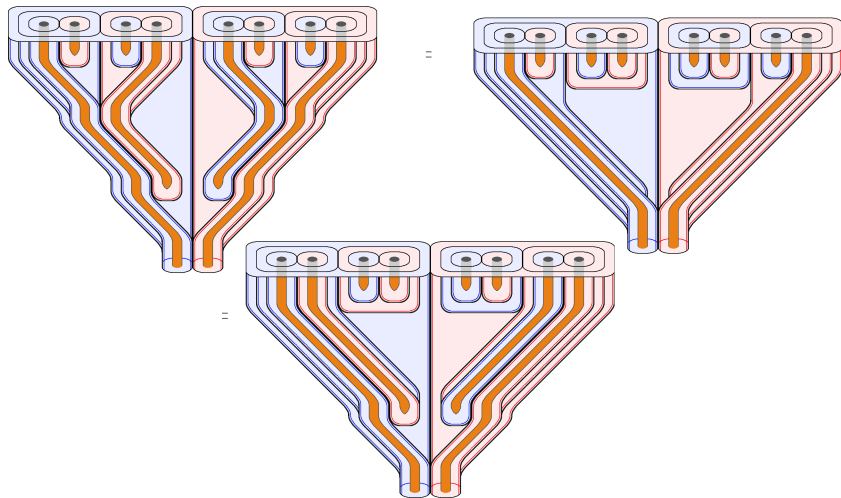


# Product of Monads

This is a monad. Proof:

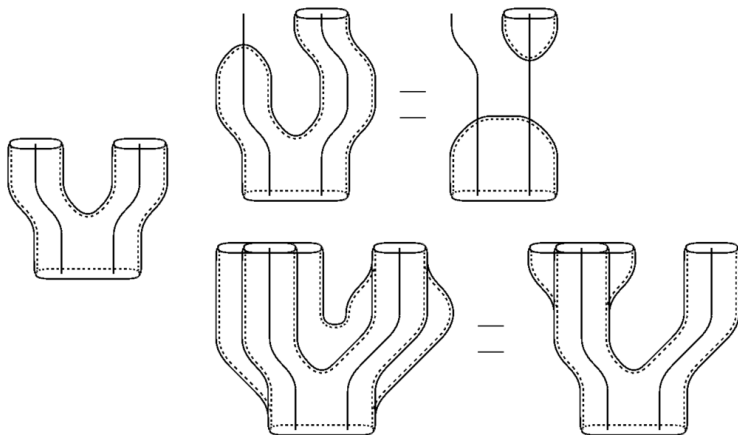


# Product of Monads

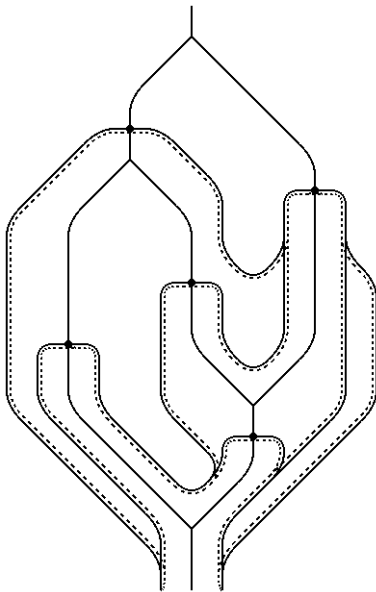
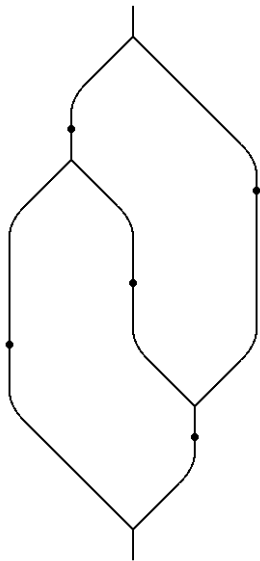


## Personal motivation: Comonad interactions

- ▶ Models with both leading/active monadic structure, and following/passive comonad structure.
- ▶ Interaction between programs as lead and follow, using interaction laws.



## Personal motivation: Comomonad interactions



## Multi-functors generically

Functors like  $X \times TY$  have multiple arguments.

Generically, this can be illustrated by glueing cylinders.

Natural transformations can be illustrated by connecting the appropriate arguments.

For example, coproduct in **Set** can be represented by glued cylinders. Below is an illustration of distributivity.

$$\left( \delta_{X,Y,Z} : \frac{X \times (Y + Z)}{(X \times Y) + (X \times Z)} \right) =$$



## Conclusions: Open questions

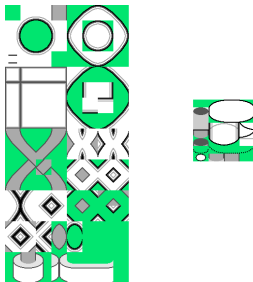
How to do cylinder diagrams monoidally?

Geometric interpretation of cylinder diagrams.

Considering other generic features.

## Conclusions: Pictures

Most images were composed by combining tiles in layers, using the following tilesets:



Working on more automatized methods too.

String drawing method on grid: [Link](#)