

Monoidal Grothendieck Construction

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Motivation

$$k \in \text{Ring}$$

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$$k \in \text{Ring} \rightsquigarrow \text{Mod}_k \in \text{Cat}$$

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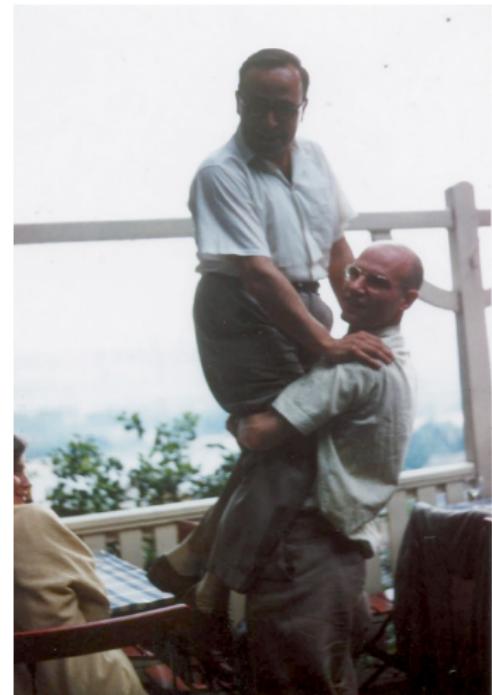
$$k \in \text{Ring} \rightsquigarrow \text{Mod}_k \in \text{Cat}$$

$\text{Mod}_{all} ???$

Motivation

Grothendieck: Yes!

- ▶ objects (k, M) , where $M \in \text{Mod}_k$
- ▶ maps $(f, g): (k, M) \rightarrow (k', M')$
where $f: k \rightarrow k'$ and
 $g: M \rightarrow f^*(M')$



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$$f: k \rightarrow k' \rightsquigarrow f^*: \text{Mod}_{k'} \rightarrow \text{Mod}_k$$

Motivation

$$k \in \text{Ring} \rightsquigarrow \text{Mod}_k \in \text{Cat}$$

$$\text{Mod}: \text{Ring}^{\text{op}} \rightarrow \text{Cat}$$

$$f: k \rightarrow k' \rightsquigarrow f^*: \text{Mod}_{k'} \rightarrow \text{Mod}_k$$

Motivation

Given

$$\text{Mod}: \text{Ring}^{\text{op}} \rightarrow \text{Cat}$$

we defined Mod_{all} to have

- ▶ objects (k, M) , where
 $M \in \text{Mod}_k$
- ▶ maps
 $(f, g): (k, M) \rightarrow (k', M')$
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Given

$$\mathcal{F}: \mathcal{X}^{\text{op}} \rightarrow \text{Cat}$$

we define $\int \mathcal{F}$ to have

- ▶ objects (x, a) , where $x \in \mathcal{X}$, $a \in \mathcal{F}(x)$
- ▶ maps $(f, g): (x, a) \rightarrow (x', a')$
where $f: x \rightarrow x'$ and
 $g: a \rightarrow \mathcal{F}f(a')$

Indexed Categories

2-category $\text{ICat}(\mathcal{X})$:

- ▶ an **indexed category** is a pseudofunctor $\mathcal{F}: \mathcal{X}^{\text{op}} \rightarrow \text{Cat}$.
- ▶ an **indexed functor** is a pseudonatural transformation
 $\alpha: \mathcal{F} \Rightarrow G$
- ▶ an **indexed natural transformation** is a modification
 $m: \alpha \Rrightarrow \beta$

Fibrations

$$\begin{array}{ccc} \mathcal{A} & & b \\ P \downarrow & & \\ \mathcal{X} & \xrightarrow{f} & y \end{array}$$

Fibrations

cartesian lift

$$\begin{array}{ccc} \mathcal{A} & & \\ P \downarrow & & \\ \mathcal{X} & & \end{array} \quad \begin{array}{ccc} a & \xrightarrow{\phi} & b \\ & & \\ x & \xrightarrow{f} & y \end{array}$$

Fibrations

pullback

$$\begin{array}{ccc} \mathcal{A} & & f^*(b) \xrightarrow{\phi} b \\ P \downarrow & & \\ \mathcal{X} & & x \xrightarrow{f} y \end{array}$$

Fibrations

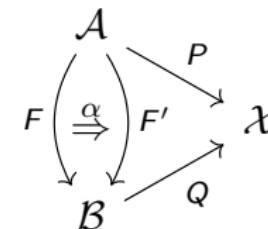
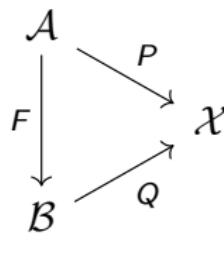
reindexing functor

$$\mathcal{A}_y \xrightarrow{f^*} \mathcal{A}_x$$

Fibrations

2-category $\text{Fib}(\mathcal{X})$:

- ▶ an object is a fibration
 $P: \mathcal{A} \rightarrow \mathcal{X}$
- ▶ a 1-morphism is a functor F
- ▶ a 2-morphism is a natural transformation



which preserves cartesian
liftings

The Grothendieck Construction

For indexed category $F: \mathcal{X}^{\text{op}} \rightarrow \text{Cat}$, $\int F$ is naturally fibred over \mathcal{X} :

$$\begin{aligned} P_F: \int F &\rightarrow \mathcal{X} \\ (x, a) &\mapsto x \\ (f, k) &\mapsto f \end{aligned}$$

2-Equivalence

Theorem

The Grothendieck construction gives an equivalence:

$$\mathrm{ICat}(\mathcal{X}) \cong \mathrm{Fib}(\mathcal{X})$$

Fibre-wise Monoidal Grothendieck Construction

A **(fibre-wise) monoidal indexed category** is

- ▶ a pseudofunctor $\mathcal{F}: \mathcal{X}^{\text{op}} \rightarrow \text{MonCat}$

Let $f\text{MonICat}(\mathcal{X})$ denote the 2-category of fibre-wise monoidal indexed categories

Fibre-wise Monoidal Grothendieck Construction

A **(fibre-wise) monoidal fibration** is

- ▶ fibration $P: \mathcal{A} \rightarrow \mathcal{X}$
- ▶ the fibres \mathcal{A}_x are monoidal
- ▶ the reindexing functors are monoidal

Let $f\text{MonFib}(\mathcal{X})$ denote the 2-category of fibre-wise monoidal fibrations.

Fibre-wise Monoidal Grothendieck Construction

Theorem (Vasilakopoulou, M)

The Grothendieck construction lifts to an equivalence:

$$f\text{MonFib}(\mathcal{X}) \simeq f\text{MonICat}(\mathcal{X})$$

Global Monoidal Grothendieck Construction

A **(global) monoidal indexed category** is

- ▶ an indexed category $\mathcal{F}: \mathcal{X}^{\text{op}} \rightarrow \text{Cat}$
- ▶ \mathcal{X} is monoidal
- ▶ \mathcal{F} is lax monoidal $(\mathcal{F}, \phi): (\mathcal{X}^{\text{op}}, \otimes) \rightarrow (\text{Cat}, \times)$

Let $g\text{MonICat}$ denote the 2-category of global monoidal indexed categories.

Global Monoidal Grothendieck Construction

A **(global) monoidal fibration** is a fibration $P: \mathcal{A} \rightarrow \mathcal{X}$

- ▶ \mathcal{A} and \mathcal{X} are monoidal
- ▶ P is a strict monoidal functor
- ▶ $\otimes_{\mathcal{A}}$ preserves cartesian liftings.

Let $g\text{MonFib}(\mathcal{X})$ denote the 2-category of global monoidal fibrations.

Global Monoidal Grothendieck Construction

Theorem (Vasilakopoulou, M)

The Grothendieck construction lifts to an equivalence:

$$g\text{MonFib}(\mathcal{X}) \simeq g\text{MonICat}(\mathcal{X})$$

Monoidal structure on the total category

Given a lax monoidal functor

$$(\mathcal{F}, \phi): (\mathcal{X}^{\text{op}}, \otimes) \rightarrow (\text{Cat}, \times)$$

$$\phi: \mathcal{F}x \times \mathcal{F}y \rightarrow \mathcal{F}(x \otimes y)$$

$$(x, a) \otimes (y, b) =$$

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$$(x, a) \otimes (y, b) = (x \otimes y, \phi_{x,y}(a, b))$$

$$\begin{aligned}
 & F(x) \times F(y) \\
 & \text{(c,d)} \quad \phi_{xy} \\
 & \int F = c^\circ + d^\circ + \dots + F(x \otimes y)
 \end{aligned}$$

The diagram illustrates the decomposition of a product object $F(x) \times F(y)$ into simpler components. At the top, a circle labeled (c,d) is connected by an arrow labeled ϕ_{xy} to a sum of objects below. The sum consists of three circles: one labeled c° , one labeled d° , and a third circle containing a dot with an arrow pointing to it, labeled \dots . Below this sum is the expression $F(x \otimes y)$.

Credit: Bartosz Milewski

Shulman's Monoidal Grothendieck Construction

Theorem (Shulman)

If \mathcal{X} is cartesian monoidal, then

$$g\text{MonFib}(\mathcal{X}) \simeq f\text{MonICat}(\mathcal{X})$$

Cartesian Case

Theorem (Vasilakopoulou, M)

If \mathcal{X} is a cartesian monoidal category, then

$$\begin{array}{ccc} g\text{MonFib}(\mathcal{X}) & \xrightarrow{\cong} & g\text{MonICat}(\mathcal{X}) \\ \downarrow \varphi & & \downarrow \varrho \\ f\text{MonFib}(\mathcal{X}) & \xrightarrow{\cong} & f\text{MonICat}(\mathcal{X}) \end{array}$$

Example: Modules

$\text{Mod}: \text{Ring}^{\text{op}} \rightarrow \text{Cat}$

$(f: k \rightarrow k') \mapsto (f^*: \text{Mod}_{k'} \rightarrow \text{Mod}_k)$

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$\int \text{Mod}:$

$(k, M) \xrightarrow{(f,g)} (k', N)$

Example: Modules

$$(\text{Mod}, \mu): (\text{Ring}^{\text{op}}, \otimes) \rightarrow (\text{Cat}, \times)$$

$$\begin{aligned}\mu: \text{Mod}_k \times \text{Mod}_{k'} &\rightarrow \text{Mod}_{k \otimes k'} \\ (M, N) &\mapsto M \otimes_{\mathbb{Z}} N\end{aligned}$$

Example: Modules

$$(\text{Mod}, \mu): (\text{Ring}^{\text{op}}, \otimes) \rightarrow (\text{Cat}, \times)$$

$$\begin{aligned}\mu: \text{Mod}_k \times \text{Mod}_{k'} &\rightarrow \text{Mod}_{k \otimes k'} \\ (M, N) &\mapsto M \otimes_{\mathbb{Z}} N\end{aligned}$$

$$(\int \text{Mod}, \otimes_{\mu}):$$

$$(k, M) \otimes_{\mu} (k', N) = (k \otimes k', M \otimes_{\mathbb{Z}} N)$$

Applications

- ▶ Zunino & Turaev modules
- ▶ (Co)modules of (co)algebras
- ▶ Operad of wiring diagrams/discrete dynamical systems
- ▶ Structured cospans
- ▶ Catalysts in Petri nets

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Theory Appl. Categ., 20:No. 18, 650–738, 2008.