

Completely Positive Maps for Mixed Unitary Categories

Robin Cockett, Cole Comfort, and Priyaa Srinivasan

University of Calgary



Linearly distributive categories

A **linearly distributive category (LDC)** has two monoidal structures $(\otimes, \top, a_\otimes, u_\otimes^L, u_\otimes^R)$ and $(\oplus, \perp, a_\oplus, u_\oplus^L, u_\oplus^R)$ linked by natural transformations called the linear distributors:

$$\partial_L : A \otimes (B \oplus C) \rightarrow (A \otimes B) \oplus C$$

$$\partial_R : (A \oplus B) \otimes C \rightarrow A \oplus (B \otimes C)$$

LDCs are equipped with a graphical calculus.

LDCs provide a categorical semantics for multiplicative linear logic.

Mix categories

A **mix category** is a LDC with a mix map $m : \perp \rightarrow \top$ such that

$$\text{mx}_{A,B} : A \otimes B \rightarrow A \oplus B := \text{diagram} = \text{diagram}$$

$$(1 \oplus (u_{\oplus}^L)^{-1})(1 \otimes (m \oplus 1))\delta^L(u_{\otimes}^R \oplus 1) = ((u_{\oplus}^R)^{-1} \oplus 1)((1 \oplus m) \otimes 1)\delta^R(1 \oplus u_{\otimes}^R)$$

mx is called a **mixin**. The mixin is a natural transformation.

It is an **isomix** category if m is an isomorphism.

m being an isomorphism does not make the mixin an isomorphism. However, it does make the above coherence automatic.

A **compact LDC** is an LDC in which every mixin is an isomorphism i.e., in a compact LDC $\otimes \simeq \oplus$

†-isomix category

A **†-isomix category** is an isomix category equipped with a $\dagger : \mathbb{X}^{\text{op}} \rightarrow \mathbb{X}$ functor and the following natural isomorphisms:

$$\text{tensor laxors: } \lambda_{\oplus} : A^{\dagger} \oplus B^{\dagger} \rightarrow (A \otimes B)^{\dagger}$$

$$\lambda_{\otimes} : A^{\dagger} \otimes B^{\dagger} \rightarrow (A \oplus B)^{\dagger}$$

$$\text{unit laxors: } \lambda_{\top} : \top \rightarrow \perp^{\dagger}$$

$$\lambda_{\perp} : \perp \rightarrow \top^{\dagger}$$

$$\text{involutor: } \iota : A \rightarrow A^{\dagger\dagger}$$

such that certain coherence conditions hold.

Unitary category

A **unitary category** is a compact \dagger -isomix category in which every object has an isomorphism

$$A \xrightarrow{\varphi_A} A^\dagger$$

called the **unitary structure map** for A , such that φ_A satisfies certain coherence conditions.

An isomorphism $A \xrightarrow{f} B \in \mathbb{U}$ is said to be a **unitary isomorphism** if the following diagram commutes:

$$\begin{array}{ccc} A & \xrightarrow{\varphi_A} & A^\dagger \\ f \downarrow & & \uparrow f^\dagger \\ B & \xrightarrow{\varphi_B} & B^\dagger \end{array}$$

Mixed unitary category

A **mixed unitary category**, $M : \mathbb{U} \rightarrow \mathbb{C}$ consists of a

- \dagger -isomix category \mathbb{C}
- unitary category \mathbb{U}
- a strong \dagger - isomix functor $M : \mathbb{U} \rightarrow \mathbb{C}$
 - (M, m_\otimes, m_\top) is strong monoidal on \otimes
 - (M, n_\oplus, n_\perp) is strong comonoidal on \oplus
 - M is a linear functor
 - M preserves mix map: $n_\perp M(m) m_\top = m$
 - $(\rho, \rho^{-1}) : M((-)^\dagger) \rightarrow M(-)^\dagger$ is a linear natural isomorphism

An example of a MUC: $\mathbb{R} \subset \mathbb{C}$

Consider the discrete monoidal category \mathbb{C} :

Objects: $a + ib \in \mathbb{C}$

Maps: Identity maps only $c = c$

Tensor: multiplication

$$(a + ib) \otimes (x + iy) := (ax - by) + i(ay + bx)$$

Unit: 1

Dagger: $(a + ib)^\dagger := a - ib$

\mathbb{C} is a compact LDC ($\otimes \simeq \oplus$) with a non-stationary dagger functor.

The subcategory \mathbb{R} is a unitary category with the unitary structure map being the identity map.

$\mathbb{R} \subset \mathbb{C}$ is a mixed unitary category.

An example of a MUC: $\text{Mat}(\mathbb{C}) \subset \text{FMat}(\mathbb{C})$

A **finiteness space**, $(X, \mathcal{A}, \mathcal{B})$, consists of a set X and a subset $\mathcal{A}, \mathcal{B} \subseteq P(X)$ such that $\mathcal{B} = \mathcal{A}^\perp$, that is

$$\mathcal{B} = \{b \mid b \in P(X) \text{ with for all } a \in \mathcal{A}, |a \cap b| < \infty\},$$

and $\mathcal{A} = \mathcal{B}^\perp$.

A **finiteness relation**, $(X, \mathcal{A}, \mathcal{B}) \xrightarrow{R} (Y, \mathcal{A}', \mathcal{B}')$ is relation $X \xrightarrow{R} Y$ such that

$$\forall A \in \mathcal{A}. AR \in \mathcal{A}' \quad \text{and} \quad \forall B' \in \mathcal{B}'. RB' \in \mathcal{B}$$

Finiteness spaces with finiteness relation form a *-autonomous category.

An example of a MUC: $\text{Mat}(\mathbb{C}) \subset \text{FMat}(\mathbb{C})$

$\text{FMat}(\mathbb{C})$ is defined as follows:

Objects: Finiteness spaces $(X, \mathcal{A}, \mathcal{B})$

Maps: $(X, \mathcal{A}, \mathcal{B}) \xrightarrow{M} (Y, \mathcal{A}', \mathcal{B}')$ is a matrix $X \times Y \xrightarrow{M} \mathbb{C}$
such that

$$\text{supp}(M) := \{(x, y) \mid x \in X, y \in Y \text{ and } M(x, y) \neq 0\}$$

is a finiteness relation from $(X, \mathcal{A}, \mathcal{B})$ to $(Y, \mathcal{A}', \mathcal{B}')$.

Dagger: $(X, \mathcal{A}, \mathcal{B})^\dagger := (X, \mathcal{B}, \mathcal{A})$

M^\dagger is the complex conjugate of M .

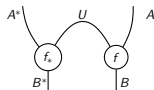
$\text{Mat}(\mathbb{C})$ is a full subcategory of $\text{FMat}(\mathbb{C})$ which is determined by the objects, $(X, P(X), P(X))$, where X is a finite set.

$\text{Mat}(\mathbb{C})$ is a unitary category, indeed a well-known \dagger -compact closed category.

The inclusion $\text{Mat}(\mathbb{C}) \subset \text{FMat}(\mathbb{C})$ is a mixed unitary category.

Goal

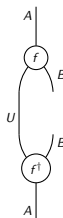
CPM construction



in \dagger - compact
closed categories

Selinger, 2007

CP[∞] construction



in \dagger -SMCs

Coecke and
Heunen, 2011

?

in MUCs

Kraus maps

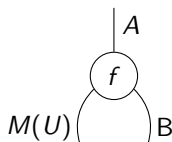
A Kraus map

$$(f, U) : A \rightarrow B$$

in a mixed unitary category, $M : \mathbb{U} \rightarrow \mathbb{C}$, is a map

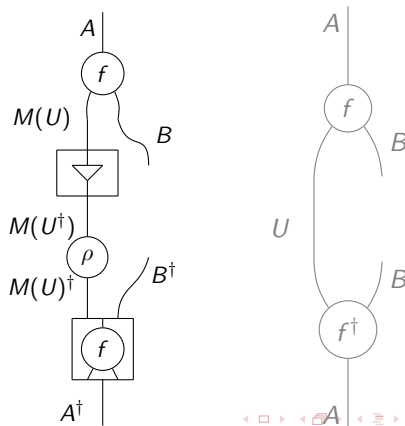
$$f : A \rightarrow M(U) \oplus B \in \mathbb{C}$$

for some $U \in \mathbb{U}$ and $M(U)$ is called the ancillary system of f .



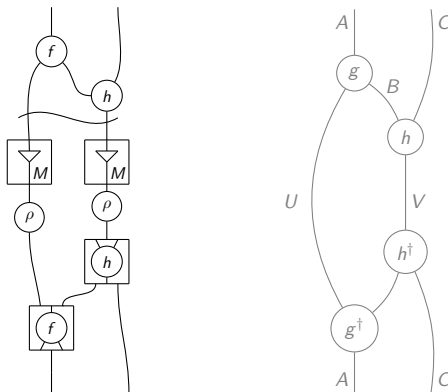
Combinator

A Kraus map can be glued along the unitary structure map with the dagger of itself to get a combinator which takes positive maps to positive maps:



Test maps

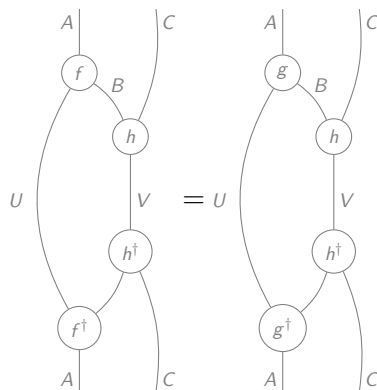
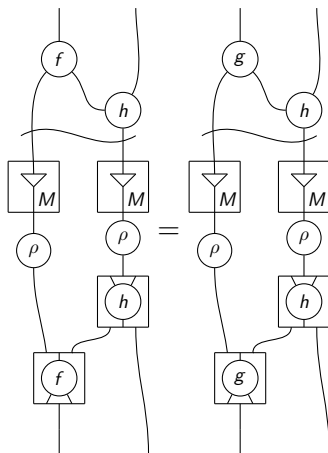
A combinator built from a Kraus map $(f, U) : A \rightarrow B$ acts on **test maps**, $h : B \otimes C \rightarrow M(V)$ as follows:



Test maps are glued along unitary structure map with its dagger to give a **positive map**.

Equivalence relation

$(f, U) \sim (g, V) : A \rightarrow B$, if for all test maps maps $h : B \otimes C \rightarrow V$, the following holds:



Unitarily isomorphic \Rightarrow equivalence

Lemma: For any two Kraus morphisms $(f, U), (g, V) : A \rightarrow B$,

$$(f, U) \sim (g, V)$$

if

- $U \xrightarrow{\alpha} V$ is a unitary isomorphism, and,
- $f(M(\alpha) \oplus 1) = f'$

CP[∞] construction

CP[∞]($M : \mathbb{U} \rightarrow \mathbb{C}$) is given as follows:

Objects: Same as \mathbb{C}

Maps: A map

$$[(f, U)] : A \rightarrow B \in \text{CP}^\infty(M : \mathbb{U} \rightarrow \mathbb{C})$$

is an equivalence class of Kraus maps

$$(f, U) : A \rightarrow B \in \mathbb{C} / \sim$$

CP[∞] construction cont...

Composition:

$$[(f, U)][(g, V)] := \left[\begin{array}{c} \text{Diagram showing the composition of } (f, U) \text{ and } (g, V) \text{ using a box } M \text{ and a direct sum operation } \oplus. \end{array} \right]$$

$$\begin{aligned} [A \xrightarrow{f} M(U) \oplus B \xrightarrow{1 \oplus g} M(U) \oplus (M(V) \oplus C) \xrightarrow{a \oplus} \\ (M(U) \oplus M(V) \oplus C \xrightarrow{n_{\oplus}^{-1} \oplus 1} M(U \oplus V) \oplus C] \end{aligned}$$

Identity: $1_A := [A \xrightarrow{(u_{\oplus}^L)^{-1}} \perp \oplus A \xrightarrow{n_{\perp}^{-1} \oplus 1} M(\perp) \oplus A]$

CP[∞] construction cont...

CP[∞](M : U → C) has two tensor products:

$$[(f, U)] \otimes [(g, V)] := \left[\begin{array}{c} \text{Diagram 1} \end{array} \right]$$

$$[(f, U)] \oplus [(g, V)] := \left[\begin{array}{c} \text{Diagram 2} \end{array} \right]$$

Unit of \otimes is \top and the unit of \oplus is \perp .

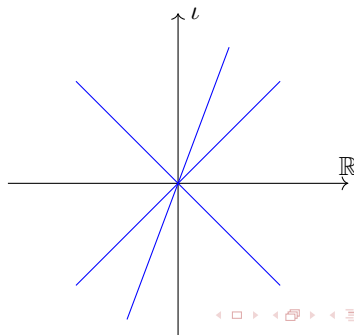
Lemma: CP[∞](M : U → C) is an isomix category.

$\mathbb{CP}^\infty(\mathbb{R} \subset \mathbb{C})$

Kraus maps in \mathbb{C} are $(=, r) : c \rightarrow c'$ such that $c = rc'$.

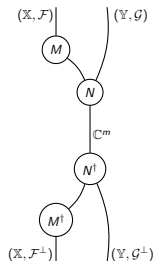
$c' \neq 0$ then there is at most one Kraus map $(=, r) : c \rightarrow c'$,
for all $c \in \mathbb{C}$ when $c = rc'$.

$c' = 0$ then $c = 0$ and for all $r' \in \mathbb{R}$,
 $(=, r) \sim (=, r') : c \rightarrow c'$.



$\text{CP}^\infty(\text{Mat}(\mathbb{C}) \subset \text{FMat}(\mathbb{C}))$

A Kraus map $(M, \mathbb{C}) : A \rightarrow B$ gives a **pure completely positive map**:



Choi's theorem for $\text{FMat}(\mathbb{C})$:

Every map in $\text{CP}^\infty(\text{Mat}(\mathbb{C}) \subset \text{FMat}(\mathbb{C}))$ can be written as a sum of pure completely positive maps.

Environment structure

An **environment structure** for a mixed unitary category, $M : \mathbb{U} \rightarrow \mathbb{C}$, is a pair

$$(F : \mathbb{C} \rightarrow \mathbb{D}, \perp)$$

where,

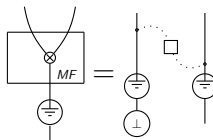
- \mathbb{D} is any isomix category,
- F is a strict Frobenius isomix functor, and
- $\perp : MF(U) \rightarrow \perp$ is a family of maps indexed by the objects $U \in \mathbb{U}$ such that the following conditions hold:

Environment structure cont...

For all unitary objects $U, V \in \mathbb{U}$, the following diagrams commute:

$$\begin{array}{ccccc}
 MF(U) \otimes MF(V) & \xrightarrow{m \times} & M(U) \oplus M(V) & \xrightarrow{\perp \oplus \perp} & \perp \oplus \perp \\
 \downarrow m_{\otimes} & & & & \downarrow u_{\oplus}^L \\
 MF(U \otimes V) & \xrightarrow{\quad \quad \quad} & & & \perp
 \end{array}$$

[Env.1a]



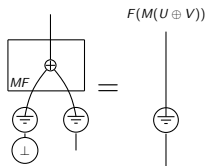
Environment structure cont...

For all unitary objects $U, V \in \mathbb{U}$, the following diagrams commute:

$$\begin{array}{ccc}
 MF(U \oplus V) & \xrightarrow{\quad \perp \quad} & \perp \\
 \downarrow n_{\oplus} & & \searrow \\
 MF(U) \oplus MF(V) & \xrightarrow{\quad \perp \oplus \perp \quad} \perp \oplus \perp \xrightarrow{\quad u_{\oplus} \quad} & \perp
 \end{array}$$

$\perp \oplus \perp$

[Env.1b]

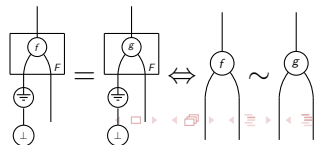


Environment structure cont...

$(f, U) \sim (g, V) \in \mathbb{C}$ if and only if the following equation holds:

$$\begin{array}{ccc}
 & MF(A) & \\
 MF(f) \swarrow & & \searrow MF(g) \\
 MF(U \oplus B) & & MF(U \oplus B) \\
 \downarrow F((n_{\oplus}^M)^{-1}) & & \downarrow F((n_{\oplus}^M)^{-1}) \\
 MF(U) \oplus MF(B) & & MF(U) \oplus MF(B) \\
 \downarrow \perp \oplus 1 & & \downarrow \perp \oplus 1 \\
 \perp \oplus MF(B) & \text{=====} & \perp \oplus MF(B)
 \end{array}$$

[Env.2]

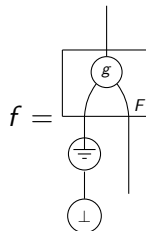


The purification axiom

An environment structure $(F : \mathbb{U}\mathbb{C}_M \rightarrow \mathbb{D}, \perp)$ satisfies **purification** axiom if

- F is injective on objects, and
- for all $f : A \rightarrow B \in \mathbb{D}$, there exists a Kraus map $(g, U) : X \rightarrow Y \in \mathbb{C}$ such that

[Env.3]



Equationally,

$$A \xrightarrow{f} B = F(A) \xrightarrow{F(g)} F(M(U) \oplus Y) \xrightarrow{n_{\oplus}} M(F(U)) \oplus F(Y) \\ \xrightarrow{\perp \oplus 1} \perp \oplus F(Y) \xrightarrow{u_{\oplus}} F(Y)$$

Examples

Lemma: For any mixed unitary category, $M : \mathbb{U} \rightarrow \mathbb{C}$, there exists an environment structure $(\mathbb{C} \xrightarrow{Q} CP^\infty(M : \mathbb{U} \rightarrow \mathbb{C}), \perp)$ satisfying purification axiom where,

$$Q(A) := A$$

$$Q(f) := [(f(u_\oplus^L)^{-1}(n_\oplus)^{-1}, \perp)]$$

$$\perp : F(M(U)) \rightarrow \perp := [((u_\oplus^R)^{-1}, U)]$$

Examples cont...

- Consider the MUC, $\mathbb{R} \subset \mathbb{C}$. Then,

$$(\mathbb{R} \xrightarrow{Q} \mathbb{CP}^\infty(\mathbb{R} \subset \mathbb{C}), \perp_r : r \rightarrow 1)$$

is an environment structure where,

$$\perp_r := (=, 1/r) : r \rightarrow 1$$

- Consider the MUC, $\text{Mat}_{\mathbb{C}} \rightarrow \text{FMat}(\mathbb{C})$. Then,

$$\text{Mat}(\mathbb{C}) \xrightarrow{Q} \mathbb{CP}^\infty(\text{Mat}(\mathbb{C}) \subset \text{FMat}(\mathbb{C}))$$

is an environment structure where,

$$\perp_{\mathbb{C}^n} : \mathbb{C}^n \rightarrow \mathbb{C}; \rho \mapsto \text{Tr}(\rho)$$

Axiomatization of CP^∞ construction

Proposition: Let \mathbb{D} be any isomix category. If $(F : \mathbb{C} \rightarrow \mathbb{D}, \perp)$ is an environment structure for the mixed unitary category $M : \mathbb{U} \rightarrow \mathbb{C}$ which satisfies the purification axiom, then

$$\mathbb{D} \simeq \text{CP}^\infty(M : \mathbb{U} \rightarrow \mathbb{C})$$

Robin Cockett, Cole Comfort, and Priyaa Srinivasan.

Dagger linear logic for categorical quantum mechanics.

ArXiv e-prints, September 2018.

Coming soon on arXiv....

Completely positive maps for mixed unitary categories.