

Quantum Sets

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two categories

Fun: sets and functions

qFun: quantum sets and quantum functions

$$\text{Inc} \dashv \text{Pts}$$

$$\begin{array}{ccccc} \text{Fun} & \xleftarrow{\text{Inc}} & \text{qFun} & \xrightarrow{\text{Pts}} & \text{Fun} \\ & \searrow & & \nearrow & \\ & & \mathbb{R} & & \\ & & \text{Id} & & \end{array}$$

Inc is full and faithful

quantum sets

definition

A quantum set \mathcal{X} is a set of nonzero finite-dimensional Hilbert spaces.

$$\text{Inc}(\mathcal{S}) = \mathcal{S} = \{ \mathbb{C}^{\{s\}} \mid s \in \mathcal{S} \}$$

$$\text{Pts}(\mathcal{X}) = \{ X \in \mathcal{X} \mid \dim(X) = 1 \}$$

qFun has terminal object $\mathbf{1} = \{ \mathbb{C} \} \cong \{ * \}$.

$$\text{Pts}(\mathcal{X}) \cong \text{qFun}(\mathbf{1}, \mathcal{X})$$

+ and \times

Let \mathcal{X} and \mathcal{Y} be quantum sets.

definition

Cartesian product $\mathcal{X} \times \mathcal{Y} = \{X \otimes Y \mid X \in \mathcal{X}, Y \in \mathcal{Y}\}$

$\mathcal{X} \times \mathcal{Y}$ is **not** the product of \mathcal{X} and \mathcal{Y}

definition

disjoint union $\mathcal{X} + \mathcal{Y} = (\mathcal{X} \times \{1\}) \cup (\mathcal{Y} \times \{2\})$

$\mathcal{X} + \mathcal{Y}$ is the coproduct of \mathcal{X} and \mathcal{Y}

$$\text{Inc}(S + T) = \text{Inc}(S) + \text{Inc}(T)$$

$$\text{Pts}(S + T) = \text{Pts}(S) + \text{Pts}(T)$$

$$\text{Inc}(S \times T) = \text{Inc}(S) \times \text{Inc}(T)$$

$$\text{Pts}(S \times T) = \text{Pts}(S) \times \text{Pts}(T)$$

quantum functions

expository definition

A quantum function F from a quantum set \mathcal{X} to a quantum set \mathcal{Y} assigns to each element X of \mathcal{X} a unitary operator

$$X \cong (H^1 \otimes Y_1) \oplus (H^2 \otimes Y_2) \oplus \cdots \oplus (H^n \otimes Y_n),$$

up to unitary equivalence of the coefficients H^1, \dots, H^n .

example: qubit measurement

$$\mathcal{X} = \{\mathbb{C}^2\}$$

$$\mathcal{X} \xrightarrow{F} \mathcal{S}$$

$$\mathcal{S} = \{\frac{1}{2}, -\frac{1}{2}\}$$

$$\mathbb{C}^2 \cong (\mathbb{C} \otimes \mathbb{C}^{\{\frac{1}{2}\}}) \oplus (\mathbb{C} \otimes \mathbb{C}^{\{-\frac{1}{2}\}})$$

composing quantum functions

$$\mathcal{X} \xrightarrow{F} \mathcal{Y} \xrightarrow{G} \mathcal{Z}$$

$$X \cong \bigoplus_i H^i \otimes Y_i$$

$$Y_i \cong \bigoplus_j K_i^j \otimes Z_j$$

$$\implies X \cong \bigoplus_{i,j} H^i \otimes K_i^j \otimes Z_j \cong \bigoplus_j \left(\bigoplus_i H^i \otimes K_i^j \right) \otimes Z_j$$

qFun is like a topos

theorem (K)

The symmetric monoidal category (\mathbf{qFun}, \times)

- 1 has finite colimits,
- 2 has finite limits,
- 3 has a terminal monoidal unit,
- 4 is closed monoidal, and
- 5 classifies subobjects by “classical” quantum functions into $\mathbf{1} + \mathbf{1}$:

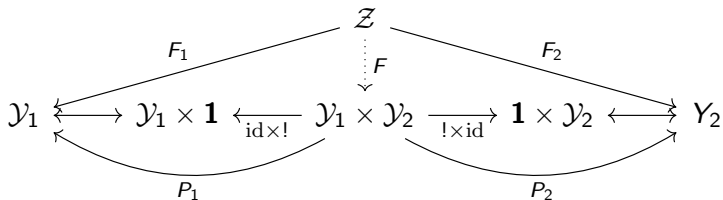
$$\begin{array}{ccc} \mathcal{Z} & \longrightarrow & \mathbf{1} \\ \downarrow & & \downarrow \tau \\ \mathcal{X} & \dashrightarrow^! & \mathbf{1} + \mathbf{1} \end{array}$$

For a symmetric monoidal category (\mathbf{C}, \times) satisfying (1) – (5):

$$(\mathbf{C}, \times) \text{ is a topos} \iff \times \text{ is a category-theoretic product}$$

compatible quantum functions

we say that F_1 and F_2 are compatible just in case



definition

A quantum function out of \mathcal{X} is classical iff it is compatible with every quantum function out of \mathcal{X} . A quantum set is classical iff $I_{\mathcal{X}}$ is classical.

classical quantum sets and classical quantum functions

proposition (K)

A quantum set \mathcal{X} is classical iff there is a set S such that $\mathcal{X} \cong 'S$.

proposition (K)

A quantum function $F: \mathcal{X} \rightarrow \mathcal{Y}$ iff there is a function $f: \mathcal{X} \rightarrow \text{Pts}(\mathcal{Y})$ with

$$\begin{array}{ccc} \mathcal{X} & \xrightarrow{F} & \mathcal{Y} \\ Q \downarrow & & \uparrow J \\ '\mathcal{X} & \xrightarrow{f} & '\text{Pts}(\mathcal{Y}) \end{array}$$

$$X \cong_Q X \otimes \mathbb{C}\{X\}$$

$$\mathbb{C}\{Y\} \cong_J \mathbb{C} \otimes Y$$

von Neumann algebras

proposition (K)

There is a full and faithful contravariant functor $\ell_q^\infty : \mathbf{qFun} \rightarrow \mathbf{vNalg}$.

$$\ell_q(\mathcal{X}) = \prod_{X \in \mathcal{X}} L(X) \qquad \ell_q^\infty(\mathcal{X}) = \left\{ a \in \ell_q(\mathcal{X}) \mid \sup_{X \in \mathcal{X}} \|a(X)\| < \infty \right\}$$

theorem (K)

Let A be a von Neumann algebra. The following are equivalent:

- 1 $A \cong \ell_q^\infty(\mathcal{X})$ for some quantum set \mathcal{X}
- 2 every von Neumann subalgebra of A is atomic
- 3 if $a^\dagger = a$, then there is an orthogonal family of projections $(p_\alpha \mid \alpha \in \mathbb{R})$

$$a = \sum_{\alpha \in \mathbb{R}} \alpha \cdot p_\alpha$$

internal ring of quantum complex numbers

Write $\mathcal{X} * \mathcal{Y}$ for category theoretic product of \mathcal{X} and \mathcal{Y} . Write $\mathcal{C} = \mathbb{R} * \mathbb{R}$.

There are quantum functions

$$\mathcal{C} * \mathcal{C} \xrightarrow{+} \mathcal{C} \quad \mathcal{C} * \mathcal{C} \xrightarrow{\cdot} \mathcal{C} \quad \mathcal{C} \xrightarrow{\dagger} \mathcal{C} \quad \mathbb{C} \hookrightarrow \mathcal{C}$$

such that the set $\text{qFun}(\mathcal{X}, \mathcal{C})$ has the structure of a \dagger -algebra over \mathbb{C} .

proposition (K)

We have a natural isomorphism $\text{qFun}(\mathcal{X}, \mathcal{C}) \cong \ell_q(\mathcal{X})$.

interlude: quantum relations

definition (essentially, Weaver)

A quantum relation R from a quantum set \mathcal{X} to a quantum set \mathcal{Y} assigns to each element X of \mathcal{X} and each element Y of \mathcal{Y} a subspace

$$R(X, Y) \leq L(X, Y)$$

Quantum relations correspond to quantum functions $\mathcal{X} \times \mathcal{Y}^* \rightarrow \mathbf{1} + \mathbf{1}$.

The category **qRel** of quantum sets and quantum relations is a dagger compact category enriched over ortholattices.

Definition

A quantum function from \mathcal{X} to \mathcal{Y} is a quantum relation such that

$$R^\dagger \circ R \geq I_{\mathcal{X}} \qquad R \circ R^\dagger \leq I_{\mathcal{Y}}$$

the graph coloring game

parameters: a finite simple graph G and a finite set S

players: Alice and Bob, cooperating blindly against a Referee

round 1: Referee plays a pair $(g_A, g_B) \in G \times G$

(Alice sees only g_A , and Bob sees only g_B)

round 2: Alice plays a color s_A and Bob plays a color s_B

(Alice sees only s_A and Bob sees only s_B)

scoring: Alice and Bob lose iff

$$(g_A = g_B \text{ and } s_A \neq s_B) \text{ or } (g_A \sim g_B \text{ and } s_A = s_B)$$

Alice and Bob have a winning strategy $\Leftrightarrow G$ can be properly colored by S

true if Alice and Bob share classical randomness

false if Alice and Bob share quantum randomness

from the graph coloring game to quantum functions

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quantum families of graph colorings

$$'G \times \mathcal{Z} \xrightarrow{F} 'S$$

$$F \circ (E_G \otimes I_{\mathcal{Z}}) \leq (\neg I_S) \circ F$$

proposition

Alice and Bob have a winning strategy using quantum entanglement iff there is a quantum family of graph colorings of G by S .