

Limits and Colimits in a Category of Lenses

Emma Chollet, Bryce Clarke, Michael Johnson, Maurine Songa,
Vincent Wang, Gioele Zardini

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ETH zürich



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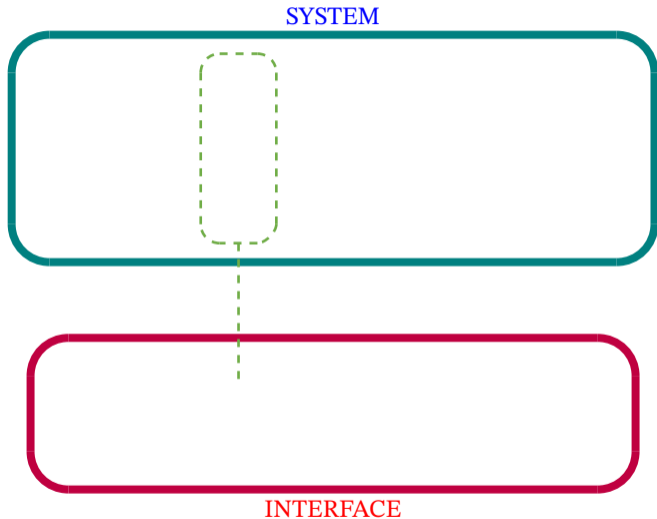


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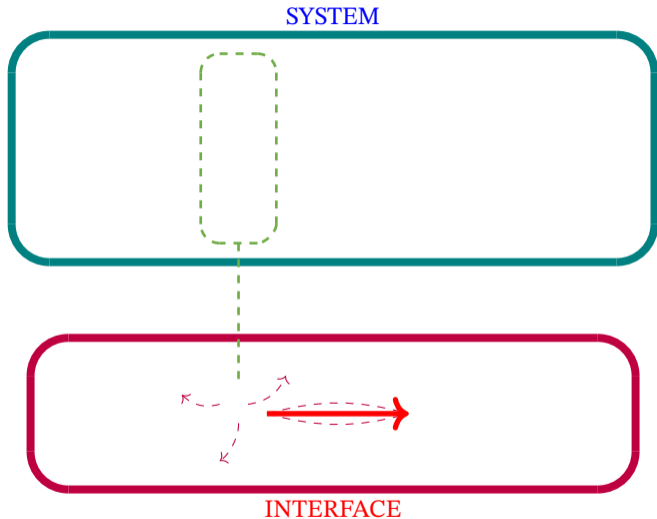


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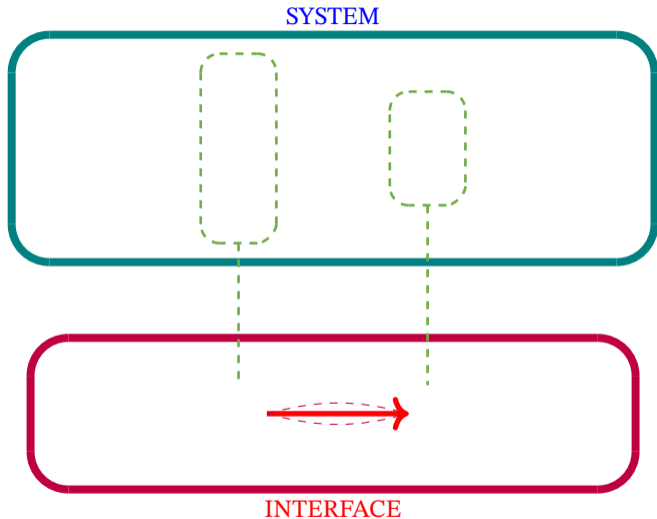
Lenses, informally



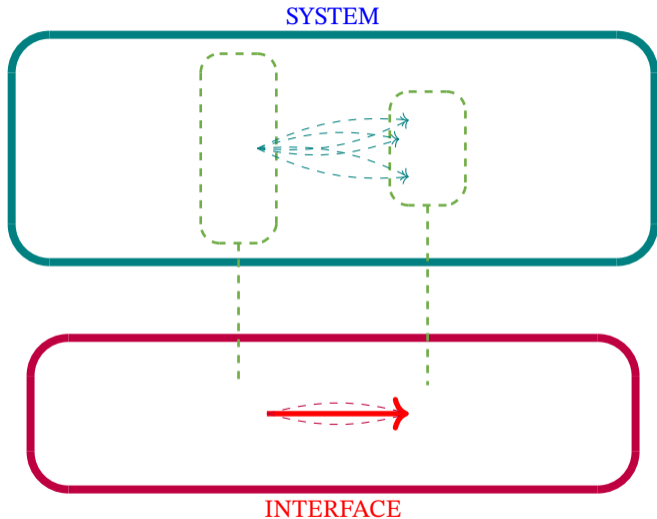
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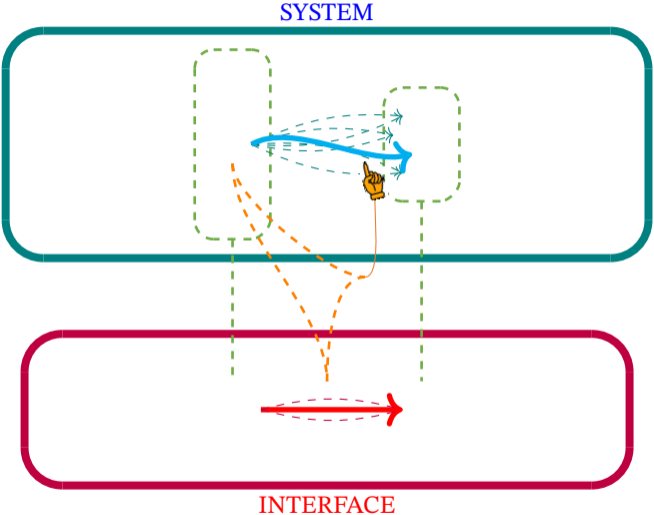
Lenses, informally



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“Lenses”

Lawful

Category-Based

Asymmetric

“Lenses”

Lawful

- (“Lawless”): Wild-West of Machine Learning, Game Theory, Economics...

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Category-Based

- Strictly generalises Set-Based Lenses
- Also known as “Delta Lenses”

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Asymmetric

- Asymmetric: One system knows everything the other does
- Symmetric: Either system may know something the other doesn't
- All Symmetric Lenses can be constructed from Asymmetric ones

“Lenses”

Lawful (Specification for)

- (“Lawless”): Wild-West of Machine Learning, Game Theory, Economics...

Category-Based (Generalisation)

- Strictly generalises Set-Based Lenses
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Asymmetric (Sufficiency)

- Asymmetric: One system knows everything the other does
- Symmetric: Either system may know something the other doesn't
- All Symmetric Lenses can be constructed from Asymmetric ones

! but Lenses

Lenses, formally (Nuts-and-Bolts)

Definition

Let \mathcal{C} and \mathcal{D} be categories. A \mathcal{C} - \mathcal{D} lens consists of a functor $L: \mathcal{C} \rightarrow \mathcal{D}$ and a lifting operation,

which satisfies the following axioms:

1. $L \circ \text{get} = \text{id}$;
2. $\text{put} \circ L = \text{id}$;
3. $\text{put} \circ \text{get} = \text{id}$.

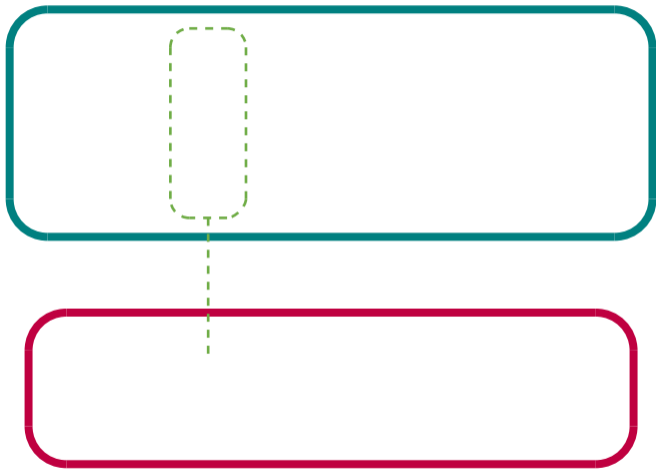
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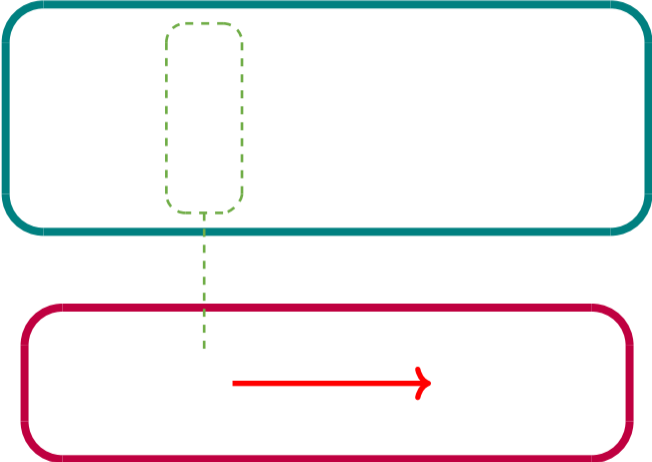
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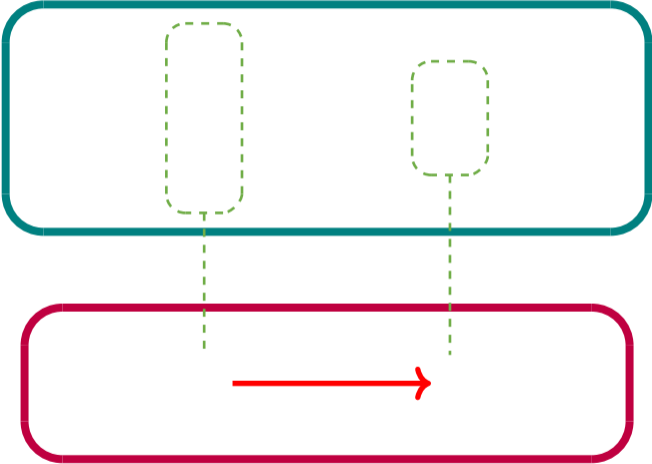
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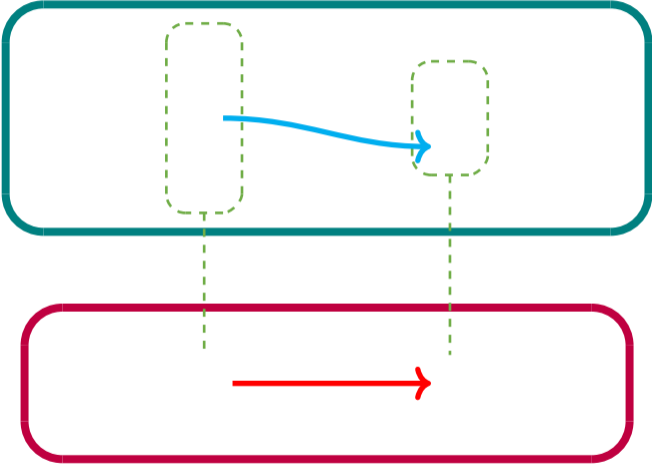
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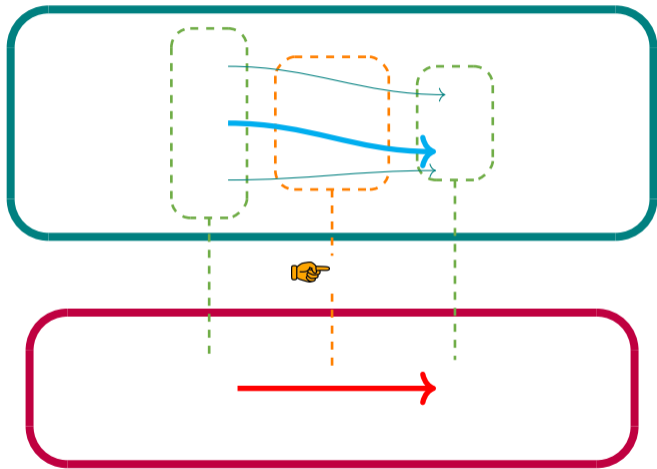
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- 1.
- 2.
- 3.

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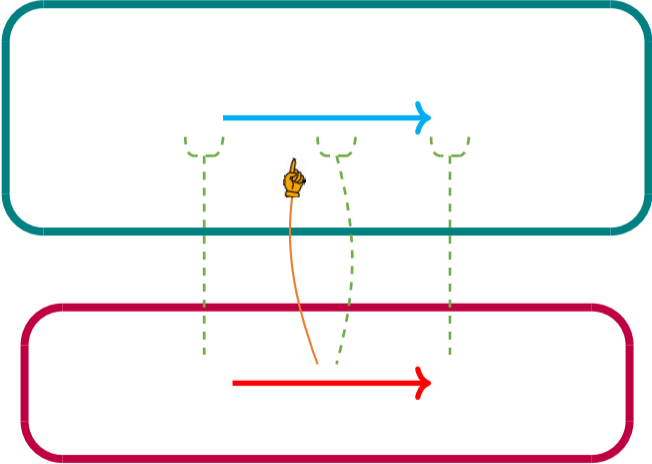
Definition

Let \mathcal{C} and \mathcal{D} be categories. A \mathcal{C} - \mathcal{D} lens consists of a functor $\text{Get} : \mathcal{C} \rightarrow \mathcal{D}$ and a lifting operation,

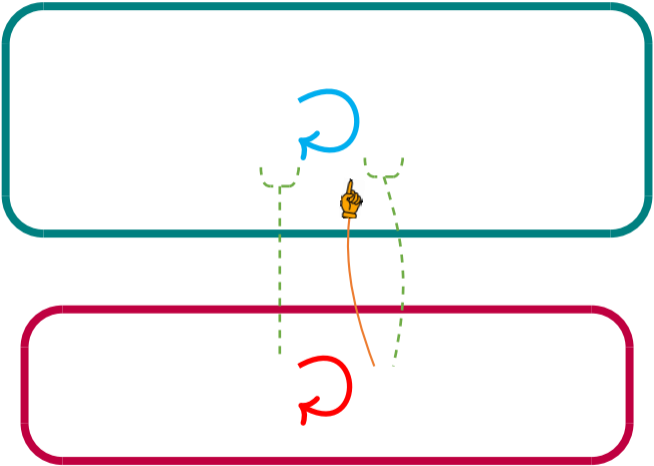
which satisfies the following axioms:

1. **Put followed by Get is trivial for morphisms**
2. **Get followed by Put preserves identities**
3. **The Put of composites is the composite of Puts**

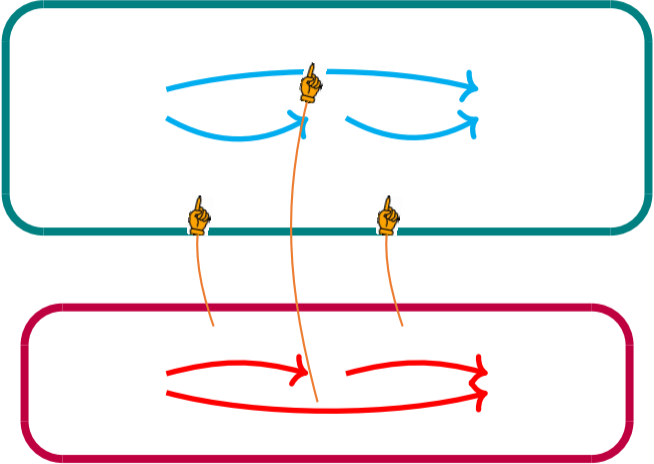
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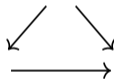
Lenses, formally (Nuts-and-Bolts)



Lenses, formally (slick)

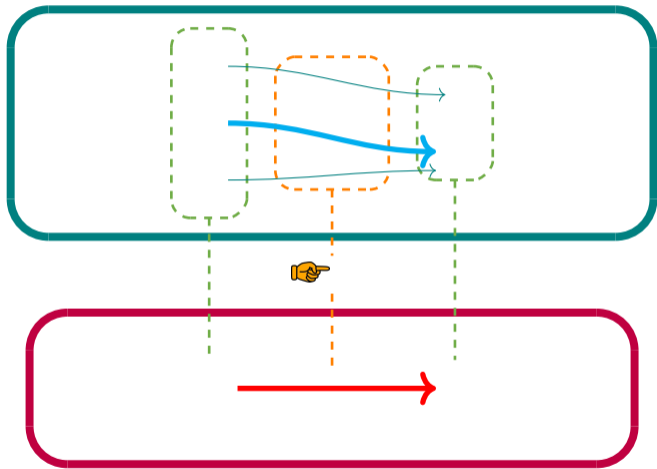
Proposition (Lenses as Functors and Cofunctors)

Every lens L may be represented as a commutative diagram of functors,

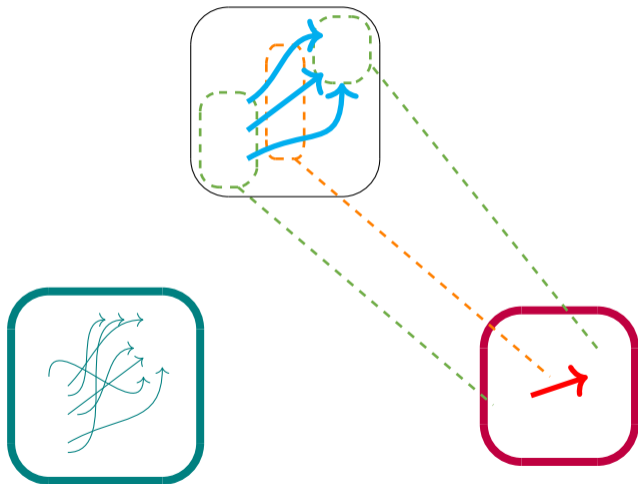


where L is a faithful, identity-on-objects functor and \mathcal{C} is a discrete opfibration.

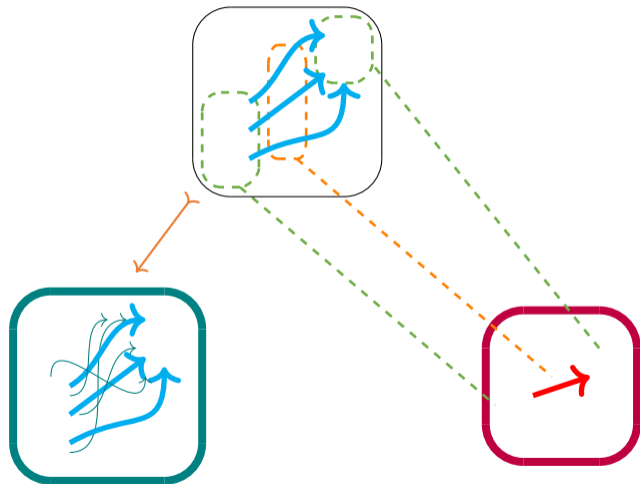
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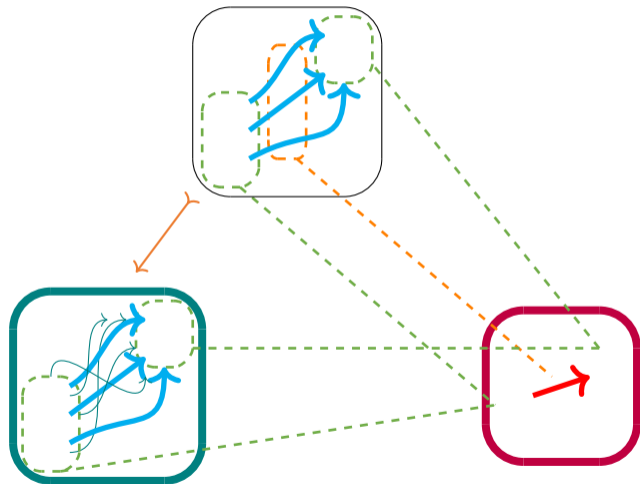
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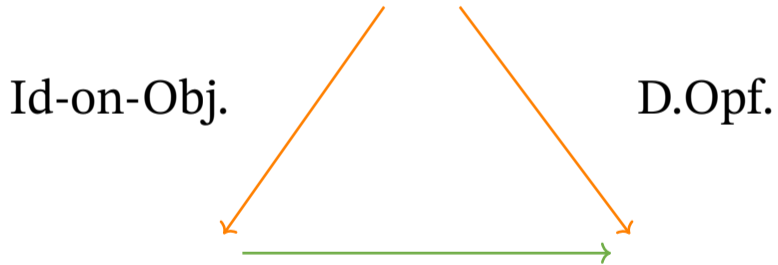
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Lenses, formally (slick)



The category

Definition

Let \mathcal{C} denote the category whose objects are categories and whose morphisms are lenses. Given a pair of lenses L and R , their composite is given by the functor $L \circ R$ together with the lifting operation:

Actually, is a pretty place

Actually, \mathcal{A} is a pretty place

has **initial** and **terminal** objects;

Actually, \mathbf{Set} is a pretty place

has **initial** and **terminal** objects;

has **small coproducts**;

Actually, \mathbf{Set} is a pretty place

has **initial** and **terminal** objects;

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has **equalisers**;

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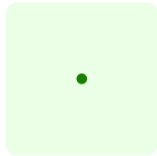
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(Next Talk): \mathcal{A} has **(certain) coequalisers**

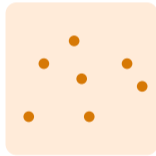
Engineering co-design as a guiding example

Design is characterised by three spaces:

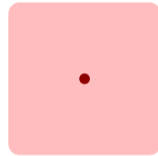
- **implementation space**: the options we can choose from;
- **functionality space**: what we need to provide/achieve;
- **requirements/costs space**: resources we need to have available;



functionality



implementations

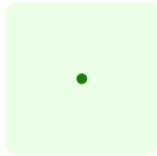


requirements

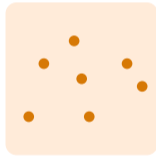
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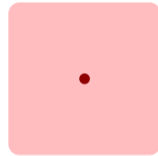
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functionality
speed



implementations
car models

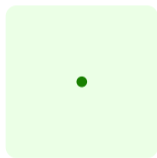


requirements
cost

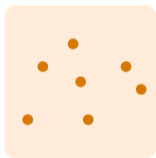
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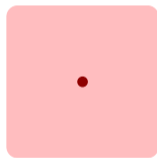
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functionality
speed
capacity max current



implementations
car models
battery models



requirements
cost
mass cost

Design problems, formally

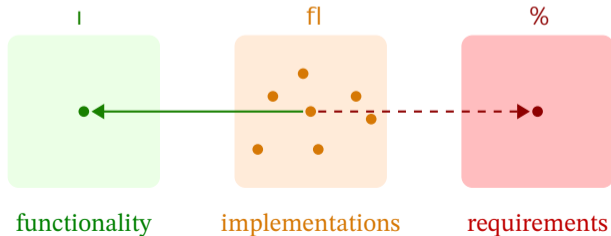
Definition

A Design Problem Instance (DPI) is a tuple $(I, \%, fl)$, where:

- I is a poset, called **functionality**;
- $\%$ is a poset, called **requirements**;
- fl is a set, called **implementations**;

the map $f_I : fl \rightarrow I$ maps an implementation to the functionality it provides;

the map $f_{\%} : fl \rightarrow \%$ maps an implementation to the resources it requires.



Practically, design problems can be understood as feasibility relations

For design purposes, we need to know **how** something is done: we need the **implementations**

For the algorithmic solution of co-design problems, we consider **feasibility relations** directly;

A fl is a **boolean profunctor**:

$I \quad \%$

fl

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A $\text{Feat} : \mathcal{R} \rightarrow \mathcal{R}$ is a **boolean profunctor**:

$$\begin{array}{ccc} \mathcal{R} & \text{Feat} & \mathcal{R} \\ \downarrow & & \downarrow \\ \mathcal{R} & & \mathcal{R} \end{array}$$

This is a **monotone** map (morphism in Feat):

- **Lower functionalities** do not require **more requirements**;
- **Higher requirements** do not provide **less functionalities**

Design problems form the category **DP**:

- Objects are posets, morphisms are design problems;
- Covered in detail in ACT4E (<https://applied-compositional-thinking.engineering>)

Realizing design problems as lenses

Consider the design problem related to buying a car based on its speed:

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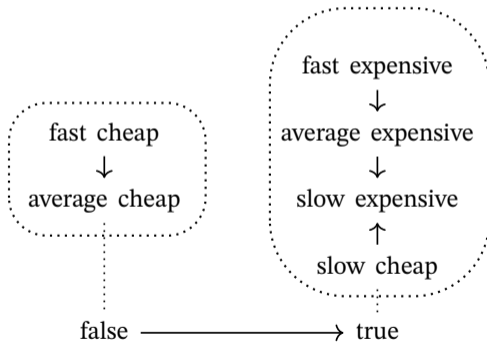
We consider $\{ \text{fast, average, slow} \}$ with posets

fast	
average	expensive
slow	cheap
	%

& $\{ \text{cheap, expensive} \}$ vehicles are the only $\{ \text{fast, average, slow} \}$ ones, the rest are $\{ \text{cheap, expensive} \}$.

Realising design problems as lenses

We can represent the functor $\mathbb{I} \quad \%$.

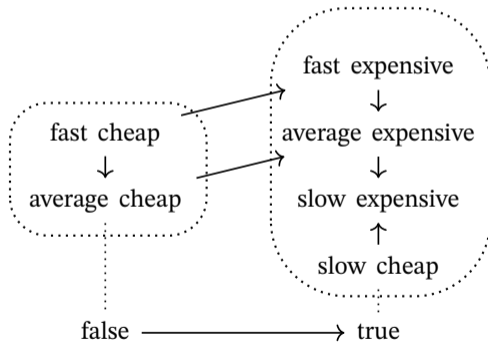


Realising design problems as lenses

We can represent the functor $| \quad \%$.

A lens over the functor provides a pair in $| \quad \%$ from each infeasible pair.

A lens models **feasibility** and informs **compromises** to make the unfeasible feasible.

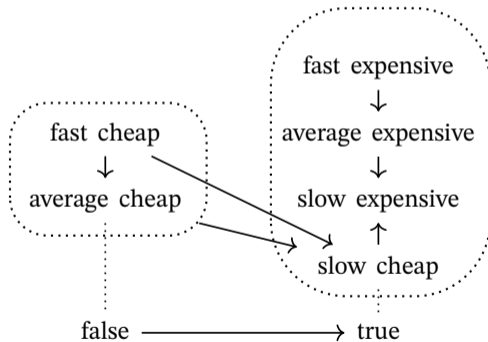


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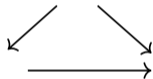


has small coproducts

Given lenses L_1, L_2 , take the coproduct in \mathcal{L} : $L_1 \amalg L_2$;

In \mathcal{L} , coproduct injection functors are **injective-on-objects discrete opfibrations**

Given lenses L_1, L_2 , we have a **unique** lens $L_1 \amalg L_2$ with:

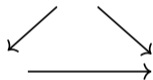


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Given lenses L_1, L_2 , take the coproduct in \mathcal{L} : $L_1 \sqcup L_2$;

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Example

From **speed** **cost** and **seats** **weight** you get

speed **cost** **seats** **weight**

has equalizers

Consider lenses L and R and

One can construct the equaliser E of the **underlying functors** in \mathcal{C} .

Then the equaliser is the largest subobject E such that $\text{pr}_L \circ \text{pr}_R$ is a discrete opfibration which forms a cone over the parallel pair in \mathcal{C} .

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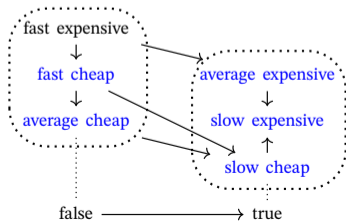
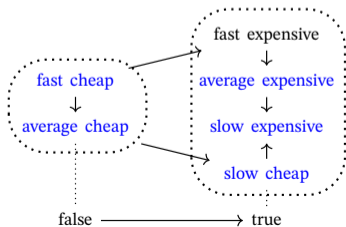
Example

Consider two design problems (two experts)

L $\%$, R $\%$

Their equalizer E $\%$:

- **embeds** into L $\%$, and selects pairs in E $\%$ for which experts **agree**
- In the worst case, **total disagreement**, i.e. $E = \emptyset$.



Conclusion and Outlook

We considered but Lenses to model problems of:

- synchronisation
- coordination
- interoperation

Conclusion and Outlook

We considered Lenses but Lenses to model problems of:

- synchronisation
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We studied the category Lenses to look for canonical constructions...

Conclusion and Outlook

We considered \mathcal{C} but Lenses to model problems of:

- synchronisation
- coordination
- interoperation

We studied the category \mathcal{L} to look for canonical constructions...

...and we found some.