

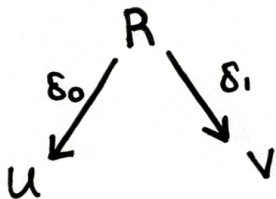
# Situated Transition Systems

(ACT 2021)

Chad Nester

Tallinn University of Technology

Span(RGraph)  $\leftrightarrow$  transition systems  $\bar{w}$  boundary



- Edges of  $U, V$  are events.
- $R$  is the transition system.
- $\delta_0, \delta_1$  map transitions to boundary events.

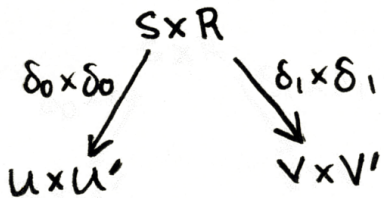
Composition  $\leftrightarrow$  Consistent events

Pullback ensures

- consistency along shared boundary



tensor product  $\leftrightarrow$  Independent systems



- trivial edges allow asynchronous execution

$$M = \begin{array}{c} \text{up} \\ \curvearrowright \\ \text{down} \end{array}$$

$$\text{Gear} : M \rightarrow M =$$



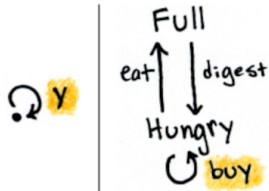
$$\text{Gear} \circ \text{Gear} : M \rightarrow M =$$



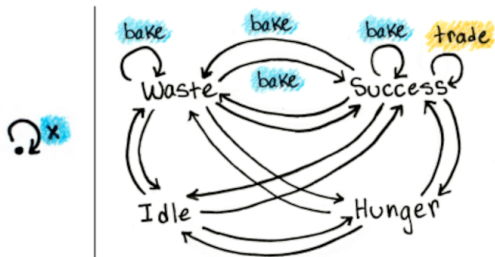
Baker :  $U \rightarrow V$



Eater :  $V \rightarrow I$



Eater  $\circ$  Baker :  $U \rightarrow I$



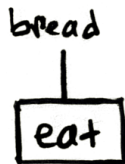
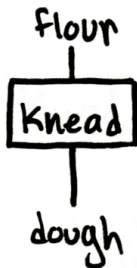
- Waste = (Open, Full)
- Success = (Open, Hungry)
- Idle = (Closed, Full)
- Hunger = (Closed, Hungry)

# Resource Theories $\leftrightarrow$ Symmetric Monoidal Categories

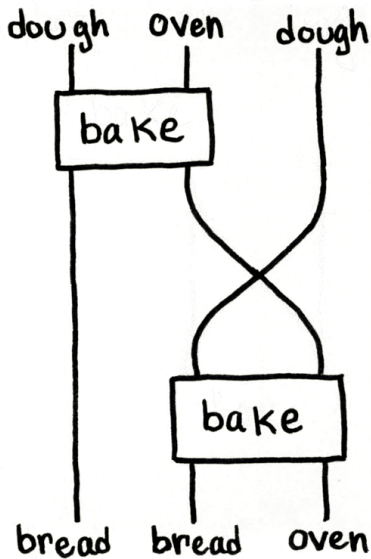
atomic objects:

$\{ \text{bread, dough, flour, oven} \}$

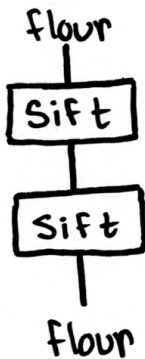
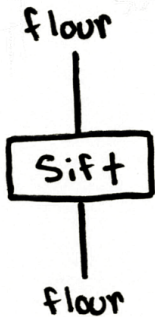
generating morphisms:



# Resource Transformations $\leftrightarrow$ String Diagrams



Same Effect  $\leftrightarrow$  Equal as Morphisms

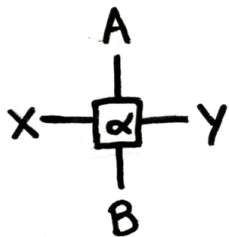


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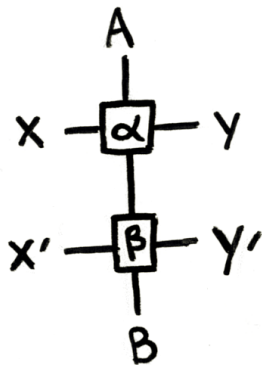
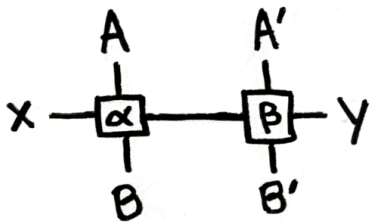




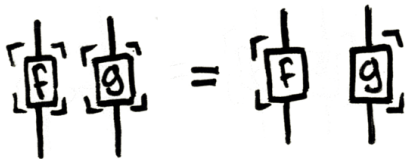
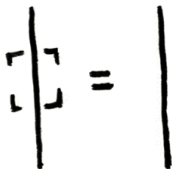
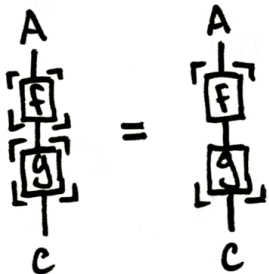
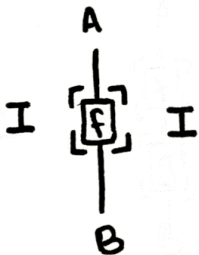
# Single object double categories



- Horizontal & vertical edge categories are monoids

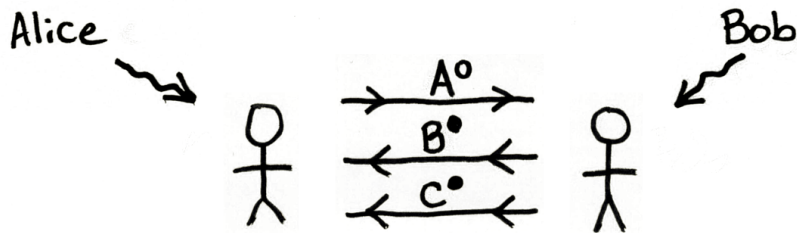


Build such a double category from  $\mathcal{A} \leftarrow \text{our resource theory}$

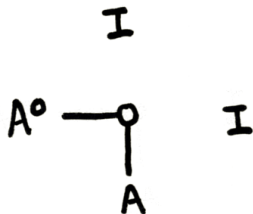
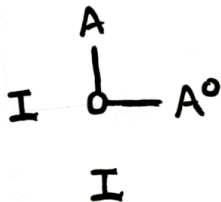


Vertical Edge Monoid  $\leftrightarrow$  Exchanges

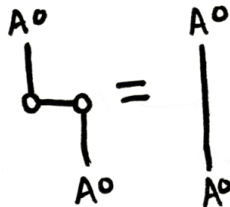
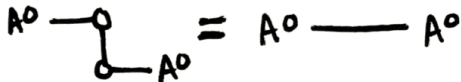
We understand  $A^\circ \otimes B^\circ \otimes C^\circ$  as in:



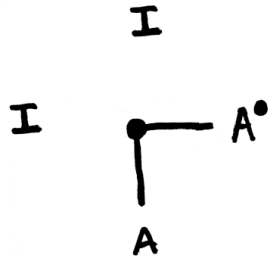
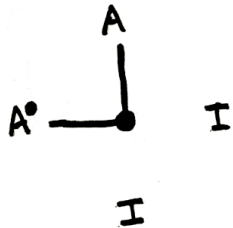
Add conjoints for each  $A \in \mathcal{A}_0$ . i.e., cells



Satisfying:



Add Companions for each  $A \in \mathcal{A}_0$ . i.e., cells

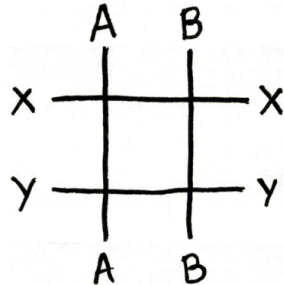
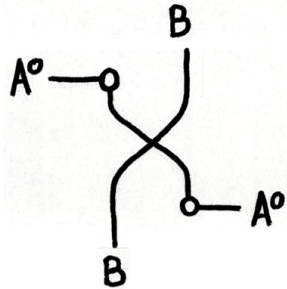
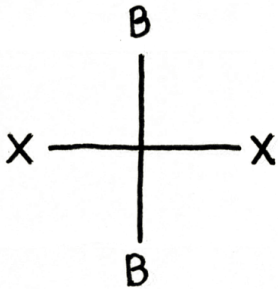


Satisfying:

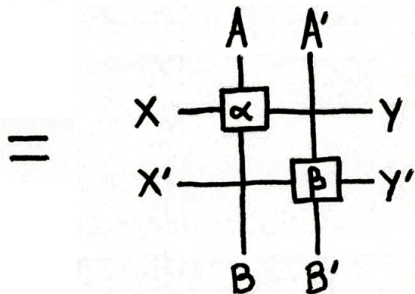
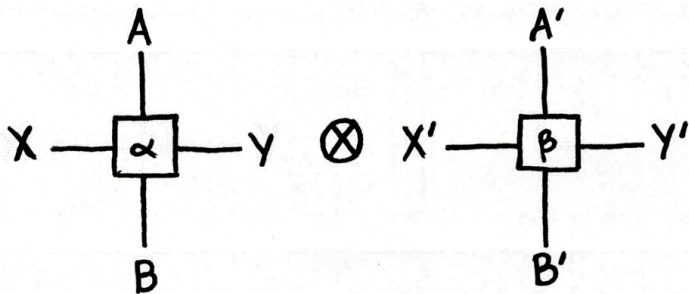
$$\begin{array}{c}
 \bullet \\
 | \\
 A^\circ - \bullet - A^\circ
 \end{array}
 = A^\circ - A^\circ$$

$$\begin{array}{c}
 A \\
 | \\
 \bullet - \bullet \\
 | \quad | \\
 A \quad A
 \end{array}
 = \begin{array}{c}
 A \\
 | \\
 A
 \end{array}$$

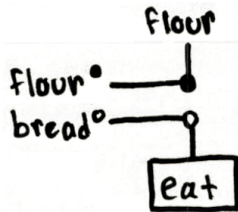
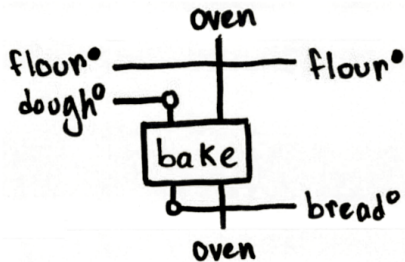
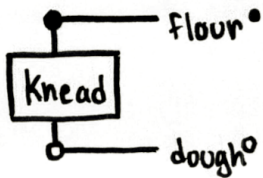
Crossing Cells arise from Corner Structure :



Crossing Cells  $\longrightarrow$  Monoidal Double Category

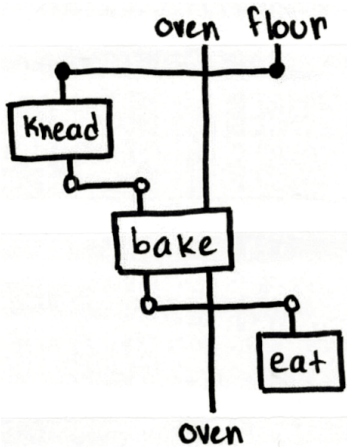


# Now Cells $\leftrightarrow$ Concurrent Resource Transformations

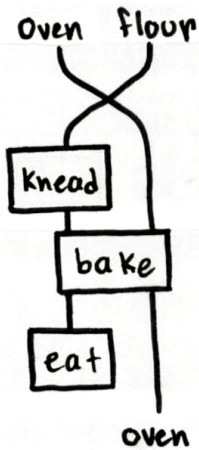




# Horizontal Composition $\leftrightarrow$ Interaction

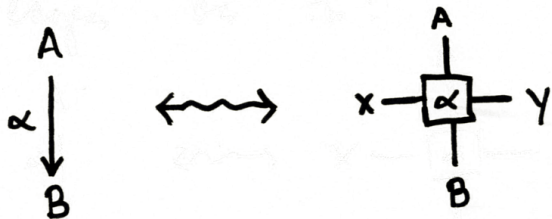


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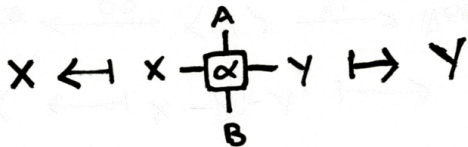
Write  $\mathbb{A}^{00}$  for the (vertical edge) monoid of exchanges

write  $\langle A \rangle$  for the graph with vertices  $\mathbb{A}_0$  and edges as in:



Then there is a span of RGraphs

$$\mathbb{A}^{00} \xleftarrow{\delta_0} \langle A \rangle \xrightarrow{\delta_1} \mathbb{A}^{00}$$



An  $A$ -situated boundary  $(U, \varphi_U)$  consists of a reflexive graph  $U$  and  $\varphi_U: U \rightarrow A^{00}$  in  $\mathbf{RGraph}$ .

An  $A$ -situated transition system  $(R, \varphi_R)$  consists of a  $\text{Span}(\mathbf{RGraph})$   $U \leftarrow R \rightarrow V$  and  $\varphi_R: R \rightarrow \langle A \rangle$  such that

$$\begin{array}{ccccc}
 U & \longleftarrow & R & \longrightarrow & V \\
 \varphi_U \downarrow & & \downarrow \varphi_R & & \downarrow \varphi_V \\
 A^{00} & \xleftarrow{\delta_0} & \langle A \rangle & \xrightarrow{\delta_1} & A^{00}
 \end{array}$$

in  $\mathbf{RGraph}$ .

$I\mathbb{A}$ -situated transition systems form a (Planar) Monoidal category,  $S(I\mathbb{A})$ .

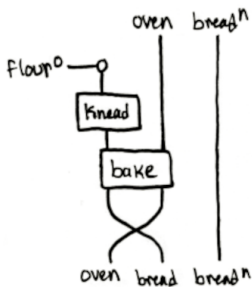
If  $I\mathbb{A}$  is compact closed, so is  $S(I\mathbb{A})$ .

$S(\mathbb{Z})$  is the category Accounts of systems with partita-doppia, introduced by Katz, Sabadini, and Walters in 1998.

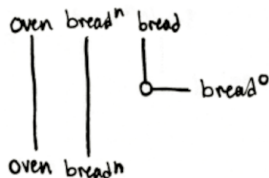
Baker : flour<sup>o</sup> → bread<sup>o</sup>



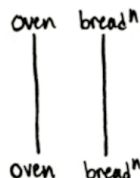
bake<sub>n</sub>



sell<sub>n</sub>



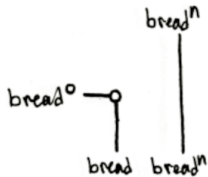
close<sub>n</sub> = open<sub>n</sub>



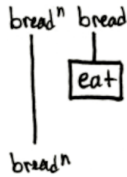
Eater :  $\text{bread}^0 \rightarrow I$



$\text{buy}_n$

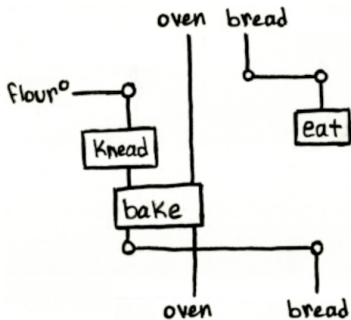
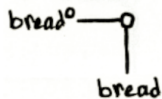
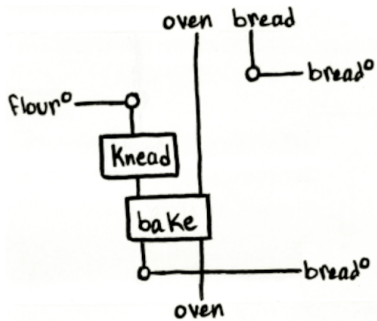


$\text{eat}_n$

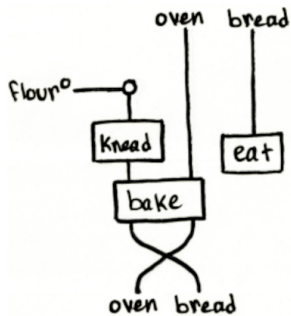


$\text{digest}_n$





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Thanks  
for  
Listening!

P. Katis, N. Sabadini, and R.F.C. Walters. On partita doppia. 1998.

Chad Nester. The Structure of Concurrent Process Histories. COORDINATION 2021.

Chad Nester. Situated Transition Systems. ACT 2021 (to appear).