Make noise and whisper: a solution to relay attacks

Omar Choudary and Frank Stajano

Computer Laboratory University of Cambridge

Mafia relay attack



Image adapted from "Keep Your Enemies Close: Distance Bounding Against Smartcard Relay Attacks", Saar Drimer and Steven Murdoch

Mafia relay attack - implemented



Photo from "Keep Your Enemies Close: Distance Bounding Against Smartcard Relay Attacks", Saar Drimer and Steven Murdoch

Existing solutions: distance-bounding





round-trip time Alice - Bob

Our solution: noise-based



- Bob introduces errors on **some** of Alice's bits
- Alice can see **some** of Bob's errors
- Alice and Bob share a secret key

Our solution: how it works (case 1)



- Bob doesn't introduce any errors
- Once Alice and Bob have sent their bit, they can't modify it

Our solution: how it works (case 1)



• The T box computes the output and ensures that neither Alice or Bob can see the other's input

Our solution: how it works (case 1)



• Both Alice and Bob see the same output

Our solution: how it works (case 2)



- Bob introduces error
- Once Alice and Bob have sent their bit, they can't modify it

Our solution: how it works (case 2)



• The T box computes the output and ensures that neither Alice or Bob can see the other's input

Our solution: how it works (case 2)



- The output is different than Alice's input so she observes Bob's error
- Bob cannot determine Alice's input

Our solution: how it works (case 3)



- Bob introduces error
- Once Alice and Bob have sent their bit, they can't modify it

Our solution: how it works (case 3)



• The T box computes the output and ensures that neither Alice or Bob can see the other's input

Our solution: how it works (case 3)



- Bob cannot determine Alice's input, AND
- Alice cannot observe if Bob has introduced an error, since the output is the same as her input

Implementation - example 1



- Bidirectional channel, e.g. ISO 7816
- Alice and Bob can send 0 (GND) or 1 (pull-up resistor)

Implementation - example 1

• Protocol:

- \circ Alice and Bob share a secret key
- For i=1,N
 - A -> B: R_K[i]
 - Bob introduces errors in 50% of bits
- At the end of the run, Bob sends the N bits received to Alice
- Alice compares the number of correct bits with the number she expected
- if a pair of attackers were involved this should be detectable



- MA and MB can relay data on perfect channel
- But they are constrained by the T box



- Alice sends a 4-bit sequence to MB "1 1 0 0"
- MB inserts error in bits 2 and 3



- Alice sends a 4-bit sequence to MB "1 1 0 0"
- MB inserts error in bits 2 and 3



- Alice expects from Bob 3 out of 4 bits correct, i.e. sequence "1 0 0 0"
- MB inserts relays this to MA



- MA forwards the data to Bob
- Bob introduces errors at bits 1 and 3



- Bob receives "0 0 0 0" and sends this to Alice using encryption
- Alice spots the difference between this and the expected "1 0 0 0" => end of transaction

Limitations

- No protection against long-cable relay attack
- In the example implementations, the attackers can read/write bits faster than Alice/Bob if they have better hardware => capability race



Presenter:

Omar Choudary (osc22@cam.ac.uk) http://www.cl.cam.ac.uk/~osc22/

Implementation - example 2



- Bidirectional channel, e.g. ISO 7816
- Alice and Bob can send 0 (GND), 1 (VCC), or listen (state Z, pull-up resistor)

Implementation - example 2

- Protocol:
 - A, B: R_K (random based on secret)
 - \circ For i=1,N
 - Alice sends either R_K[i] or listens
 - Bob does the same
 - They should only listen simultaneously for 25% of the bits
 - if a pair of attackers were involved this should be detected by
 - causing a short-circuit (bad sequence)
 - listening too much