

Isomorphisms on Generic Mutually recursive polynomial types

Some Theory and an
application

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Richness of The subject

- ▶ Mathematical logic
- ▶ Type Theory
- ▶ Programming-language Theory
- ▶ Rewriting Theory
- ▶ Number Theory
- ▶ Category Theory
- ▶ Combinatorics
- ▶ Computational algebra
- ▶ Algorithms

The setting for this talk

Type Theory

$$\tau ::= x \mid \tau_1 + \tau_2 \mid \tau_1 \times \tau_2 \mid 0 \mid 1$$

$\frac{\text{px. } \tau}{\tau}$

generic recursive
Types

Lines of investigation

- ▶ Characterise the equational theory of type isomorphism
- ▶ Study the structure of isomorphisms

2

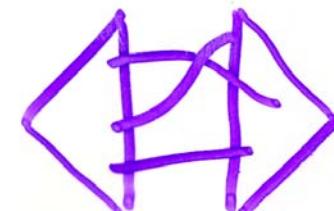
Isomorphisms between generic mutually recursive polynomial types

Example:

datatype $E = b \nmid E * E$

fun $p_0(x) = \underline{c} \underline{o} x \times \underline{s} \nmid$
 $b(x_1, x_2) \Rightarrow \underline{c} \underline{o} x \ x_1 \nmid$
 $b(x_{11}, b(x_{12}, x_2)) \Rightarrow$
 $b(x_{11}, b(x_{12}, x_2))$

$$p_0(b(b(b(x_{11}, x_{12}), x_2), x_3)) \\ = b(x_{11}, b(x_{12}, x_3))$$



3

Characterisation of type isomorphism

Thm: Two generic mutually recursive polynomial types are isomorphic iff they are provably equal in the equational theory of commutative semirings augmented with the equations $\rho x.z = z[\rho x.z/x]$.

Categorical version

Thm: The groupoid underlying

$$\mathbb{D}[x_1, \dots, x_n]/(x_1 \equiv p_1(x_1, \dots, x_n), \dots, x_n \equiv p_n(x_1, \dots, x_n))$$

is

$$\mathbb{R}[x_1, \dots, x_n]/(x_1 \equiv p_1(x_1, \dots, x_n), \dots, x_n \equiv p_n(x_1, \dots, x_n))$$

with Burnside semiring

$$\mathbb{N}[x_1, \dots, x_n]/(x_1 = p_1(x_1 - x_n), \dots, x_n = p_n(x_1 - x_n))$$

4

Deciding Type Isomorphism

Computational algebra

$$\mathbb{N}[x]/(x = p(x)) \hookrightarrow \mathbb{R} \times \mathbb{Z}[x]/(x - p(x))$$

This analysis somewhat extends to the multivariate case

5

The structure of isomorphisms

Example:

$$\text{datatype } E = b \text{ of } E * E$$

$$\begin{aligned} \text{fun } p_1(x) &= \underline{\text{co}}x \times \underline{\text{of}} \\ &\quad b(x_1, x_2) \\ &\Rightarrow \underline{\text{co}}x \times \underline{\text{of}} \\ &\quad b(x_{11}, x_{12}) \\ &\Rightarrow \underline{\text{co}}x \times \underline{\text{of}} \\ &\quad b(x_{111}, x_{112}) \\ &\Rightarrow \\ &b(b(x_{111}, b(x_{112}, x_{12})), x_2) \\ p_1(b(b(b(x_{111}, x_{112}), x_{12}), x_2), x_3) \\ &= b(b(x_{111}, b(x_{112}, x_{12})), x_2) \end{aligned}$$

$$p_1 = \begin{array}{c} \diagup \quad \diagdown \\ \square \quad \square \\ \diagdown \quad \diagup \end{array}$$

6

Thm: The automorphism group of the generic commutative idempotent is Richard Thompson's group V.

The groupoid underlying

$$G[x]/(x \cong x \circ x)$$

$$\cong S\mathbf{m}[x]/(x \equiv x \circ x)$$

and equivalent to

$$1 + V$$

► V is a very interesting (symmetry) group!

geometric view

7

Thm: The automorphism group of the generic idempotent is Richard Thompson's group F.

The groupoid

$$m[x]/(x \approx x \circ x)$$

is equivalent to

$$1 + F$$

► F is a very interesting (symmetry) group!

geometric view

Two proofs

Syntactic

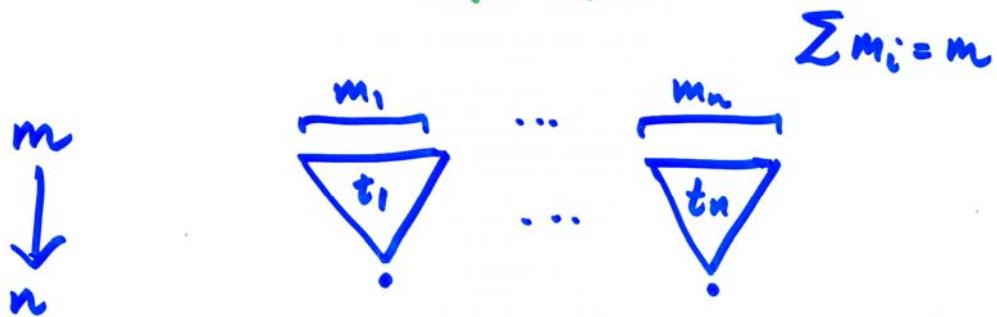
Type-theoretic analysis leads to isomorphisms being given by pattern-mappings of the form

$$b(b(x_{11}, x_{12}), x_2) \mapsto b(x_{11}, b(x_{12}, x_2))$$

Semantic

1. Free monoidal category on an operator $X \cdot X \rightarrow X$.
2. Free monoidal groupoid.

Free monoidal category
on a binary operator



10

Free (monoidal) groupoid
on a (monoidal) category

Morphisms

$$[\begin{smallmatrix} & \swarrow & \searrow & & \cdots & \swarrow & \searrow \\ A & & & & & z & \\ & \swarrow & \searrow & & \cdots & \swarrow & \searrow \\ & u & t & t' & u' & & u' \end{smallmatrix}] : A \rightarrow Z$$

where

$$\begin{smallmatrix} u & \nearrow & t & \nearrow & t' & \nearrow & u' \\ & \downarrow & & \downarrow & & \downarrow & \\ & u & & t & & t' & u' \end{smallmatrix} \sim \begin{smallmatrix} s & \nearrow & s' \\ u & \nearrow & s \\ & \downarrow & s' \\ & t & \nearrow & t' \\ & & \downarrow & \\ & & t' & \end{smallmatrix}$$

whenever

$$\begin{smallmatrix} s & \nearrow & s' \\ t & \nearrow & t' \\ & \downarrow & \\ & t' & \end{smallmatrix}$$

... This can be greatly simplified
by some rewriting ...

11

... for nice categories, like the one we are concerned with, the resulting system

$$u \begin{array}{c} \swarrow \\ t \\ \searrow \end{array} u' \rightsquigarrow u s \begin{array}{c} \swarrow \\ t' \\ \searrow \end{array} u's'$$

whenever $ts = t's'$

$$ts \begin{array}{c} \swarrow \\ t \\ \searrow \end{array} t's \rightsquigarrow t \begin{array}{c} \swarrow \\ t' \\ \searrow \end{array} t' \quad \text{sfrd}$$

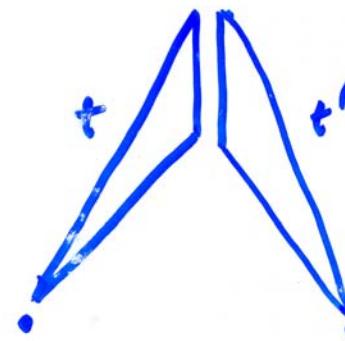
is locally confluent and terminating

12

So the automorphism group of the generic idempotent in

$$m[x]/(x \tilde{x} x \cdot x) = \text{fg}\left(m[x]/(x \cdot x \rightarrow x)\right)$$

has morphisms given by tree-cotree pairs



in normal form. (Composition is given by taking linear m.g.u.s.)

13

An application to algorithms

Sean
Cleary

Problem: Give an [efficient] algorithm that given any two binary trees with the same number of leaves yields a [minimum] number of rotations transforming one tree into the other.

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