

# The Algebra of DAGs

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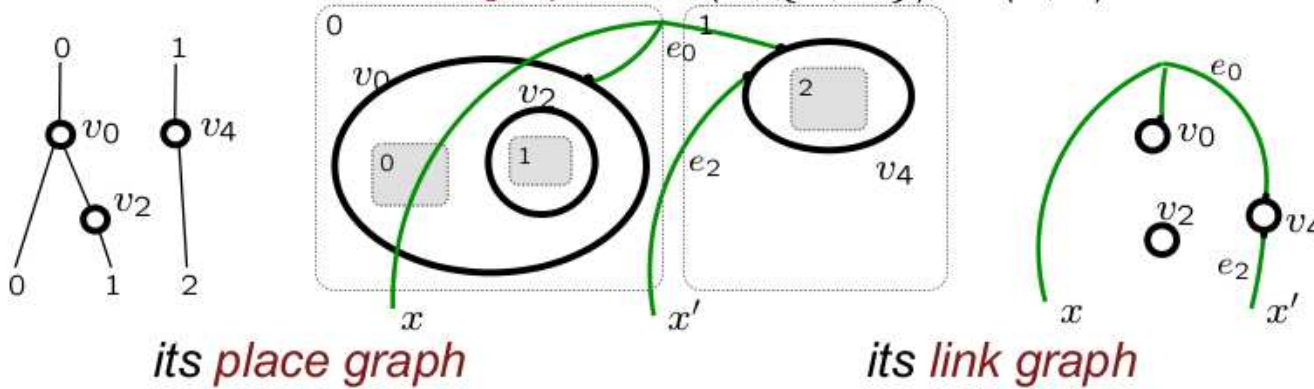
Samson@60

28.V.2013

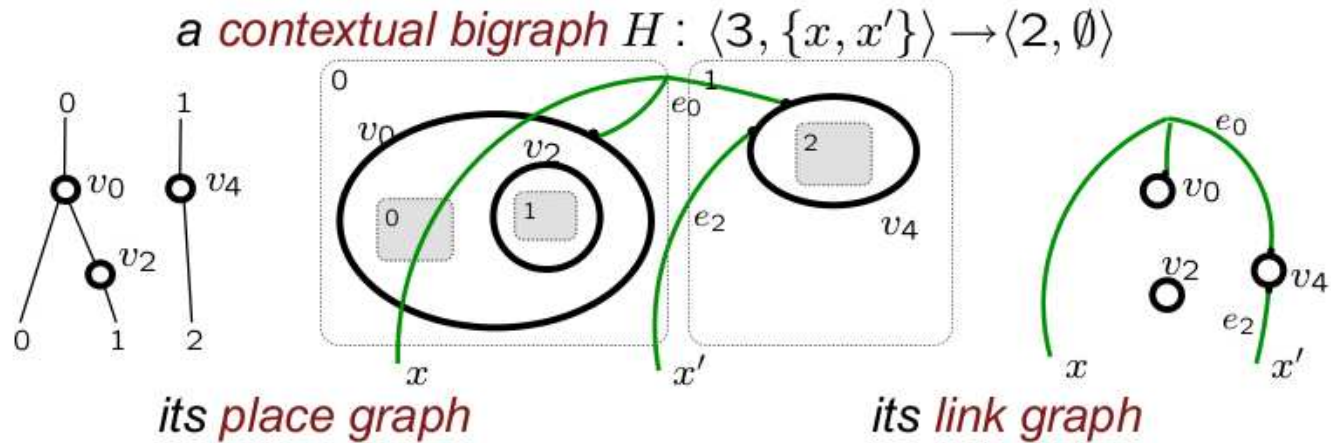
Joint work with Marco Devesas Campos

# A Question of Robin Milner

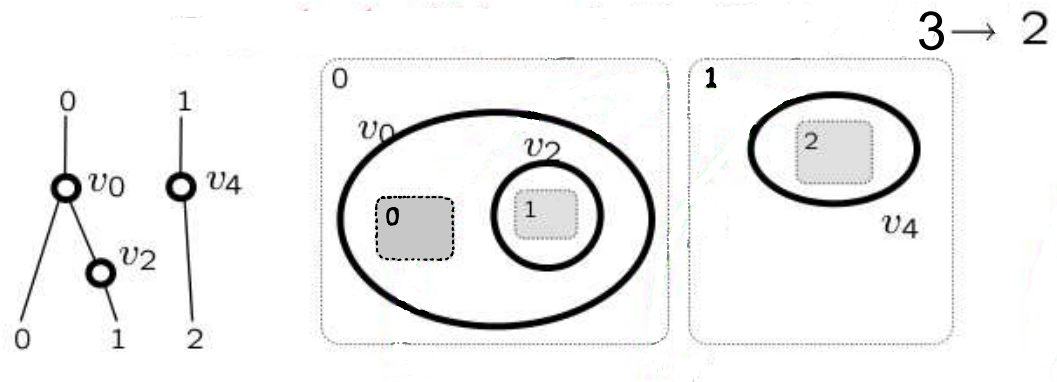
a contextual bigraph  $H: \langle 3, \{x, x'\} \rangle \rightarrow \langle 2, \emptyset \rangle$



# A Question of Robin Milner



On the generalization from tree structure ...



... to dag structure.

# Axioms for DAG structure

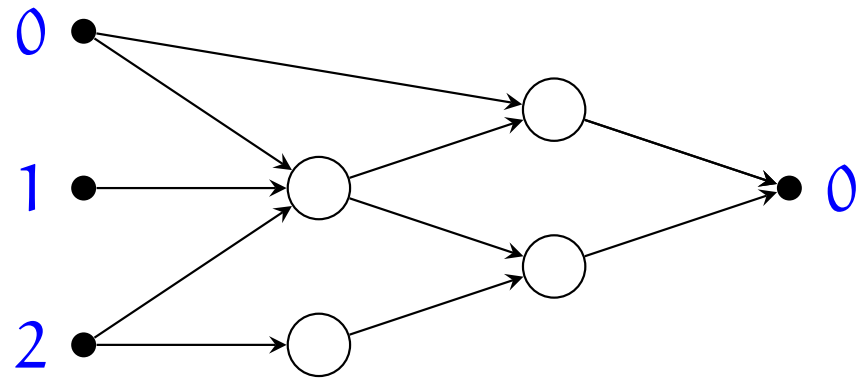
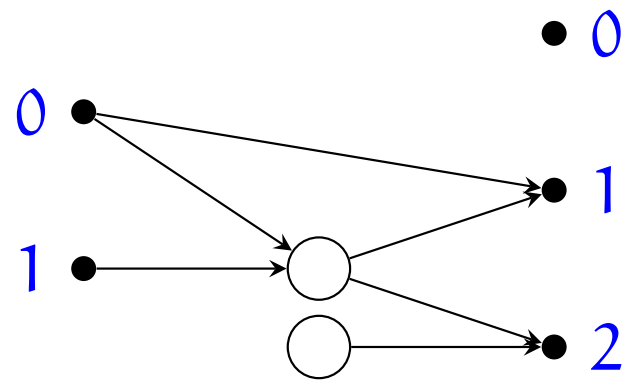
[Gibbons]

## Problem:

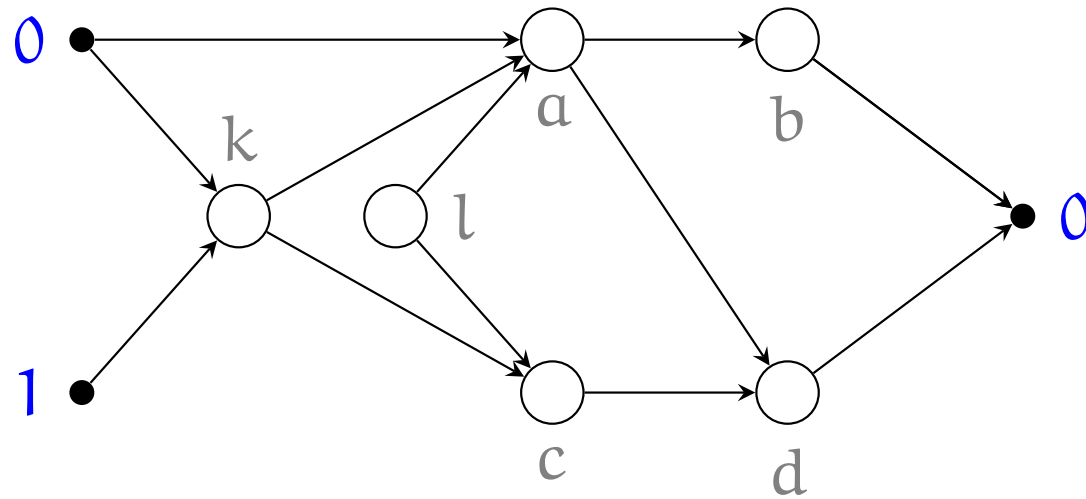
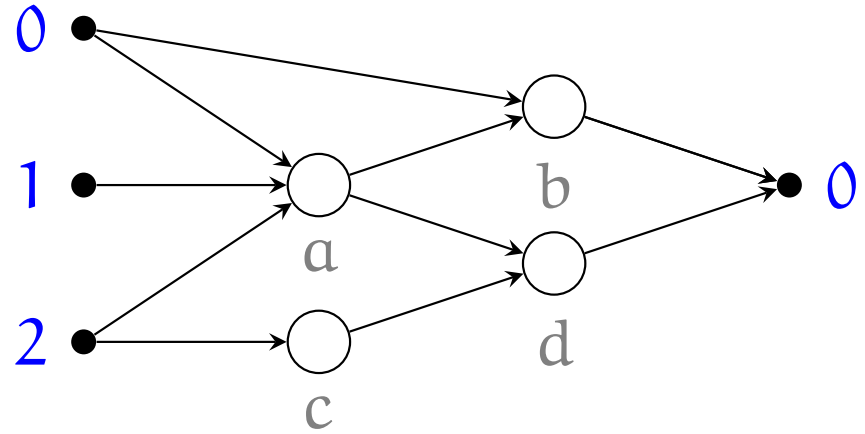
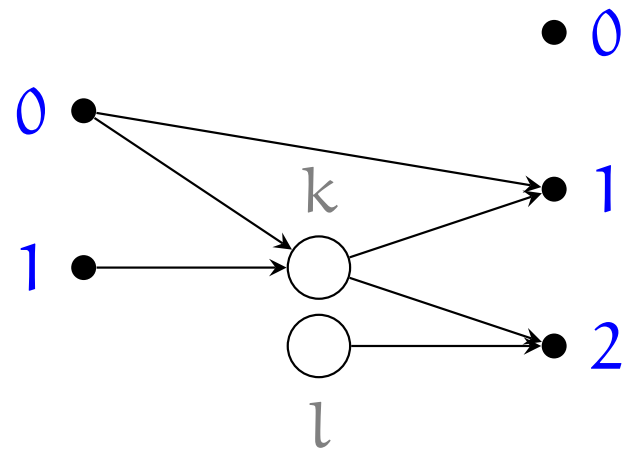
Give an algebraic characterisation of the symmetric monoidal category **Dag** with

**objects:** finite ordinals, and

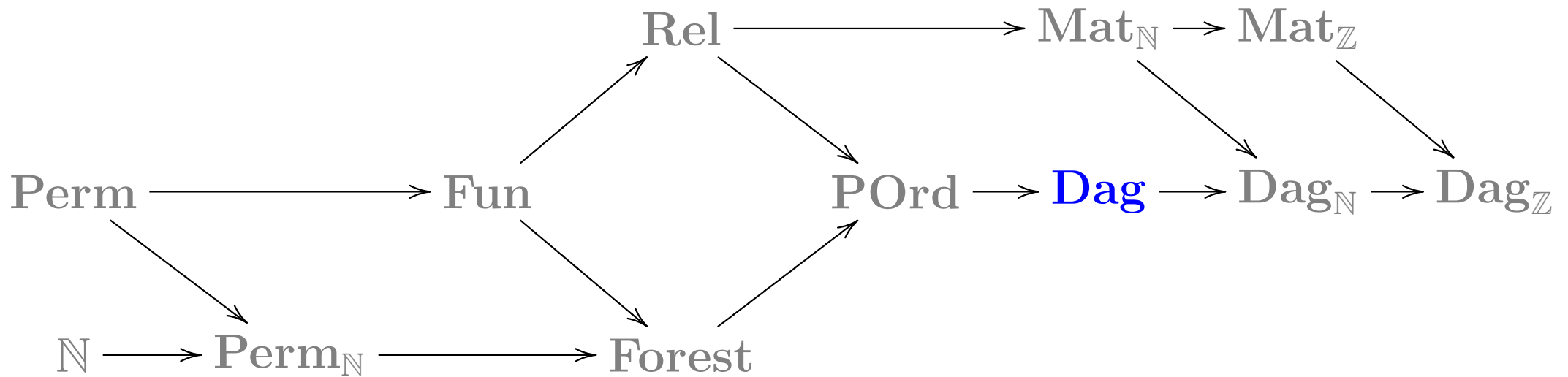
**morphisms:** finite interfaced dags.



# Composition:



# The Landscape of Algebraic Structures



# The Mathematical Setting

[Lawvere, MacLane]

## Symmetric Monoidal Equational Presentations



# The Mathematical Setting

[Lawvere, MacLane]

## Symmetric Monoidal Equational Presentations

### Examples:

#### 1. Commutative monoids

Operators

$$\eta : 0 \rightarrow 1, \quad \nabla : 2 \rightarrow 1$$

Equations

$$\nabla(x_0, \eta) \equiv x_0, \quad x_0 \equiv \nabla(\eta, x_0)$$

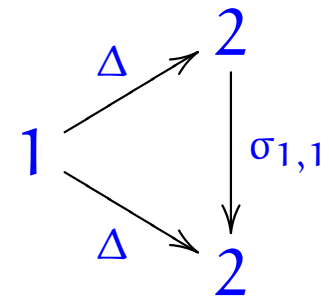
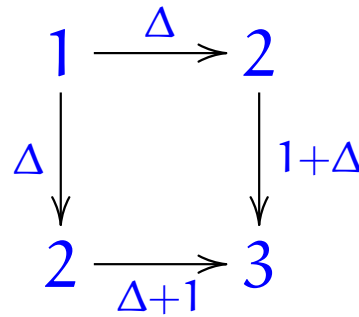
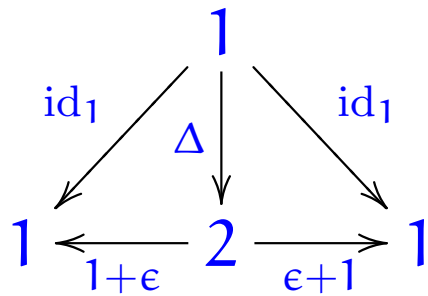
$$\nabla(x_0, \nabla(x_1, x_2)) \equiv \nabla(\nabla(x_0, x_1), x_2), \quad \nabla(x_0, x_1) \equiv \nabla(x_1, x_0)$$

## 2. Commutative comonoids

Operators

$$\epsilon : 1 \rightarrow 0, \quad \Delta : 1 \rightarrow 2$$

Equations



## PROduct and Permutation categories

**Definition:** A PROP is a symmetric strict monoidal category with underlying monoid structure on objects given by finite ordinals under addition.

### Examples:

1. **Dag**
2. The free PROP  $\mathbf{P}[\mathcal{E}]$  on a symmetric monoidal equational presentation  $\mathcal{E}$ .

$\mathbf{P}[\mathcal{E}]$  may be constructed syntactically, with morphisms given by equivalence classes of expressions generated by

$$\text{id}_n : n \rightarrow n \qquad \frac{f : \ell \rightarrow m, \quad g : m \rightarrow n}{f;g : \ell \rightarrow n}$$

$$\frac{f_1 : m_1 \rightarrow n_1, \quad f_2 : m_2 \rightarrow n_2}{f_1 + f_2 : m_1 + m_2 \rightarrow n_1 + n_2} \qquad \sigma_{m,n} : m + n \rightarrow n + m$$

$$\frac{o : n \rightarrow m \text{ an operator}}{o : n \rightarrow m}$$

under the congruence determined by the laws of symmetric strict monoidal categories together with the identities of the equational presentation  $\mathcal{E}$ .

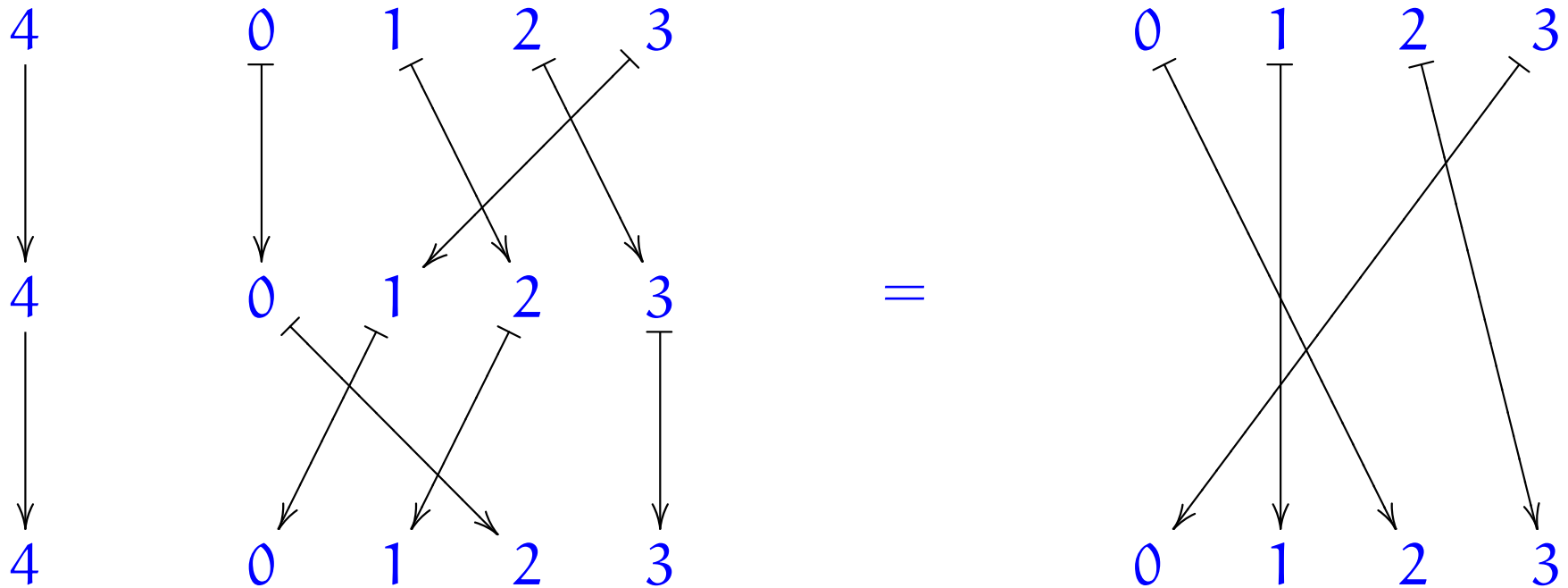
# Algebraic Characterization of DAG Structure

**Theorem:** For  $\mathcal{D}$  the symmetric monoidal equational presentation of a node together with that of degenerate commutative bialgebras,

$$\mathbf{P}[\mathcal{D}] \cong \mathbf{Dag} .$$

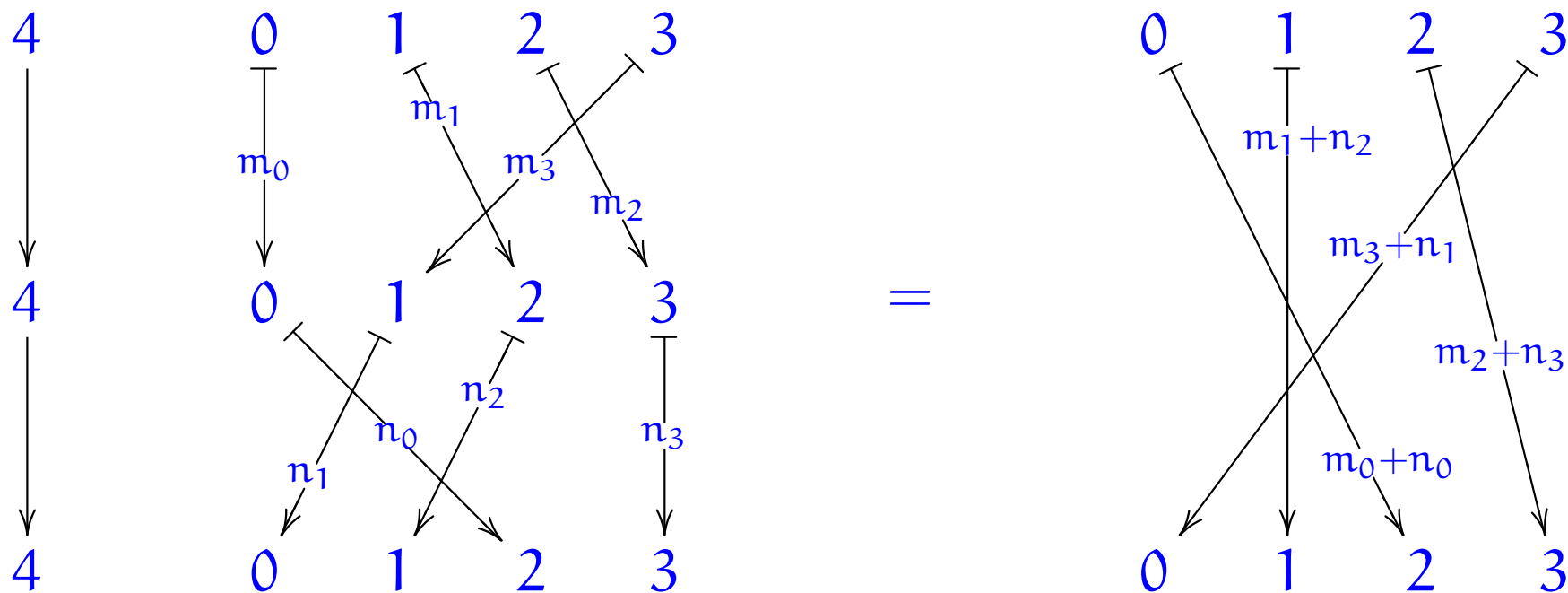
# Free PROPs

1. The free PROP  $\mathbf{P}[\emptyset]$  on the empty equational presentation is the free symmetric strict monoidal category on an object, viz. the category  $\mathbf{Perm}$  of *finite ordinals and permutations*.



2. The free PROP  $\mathbf{P}[\bullet]$  on the equational presentation of a node  
 $\bullet : 1 \rightarrow 1$  is the free symmetric strict monoidal category on the additive monoid of natural numbers

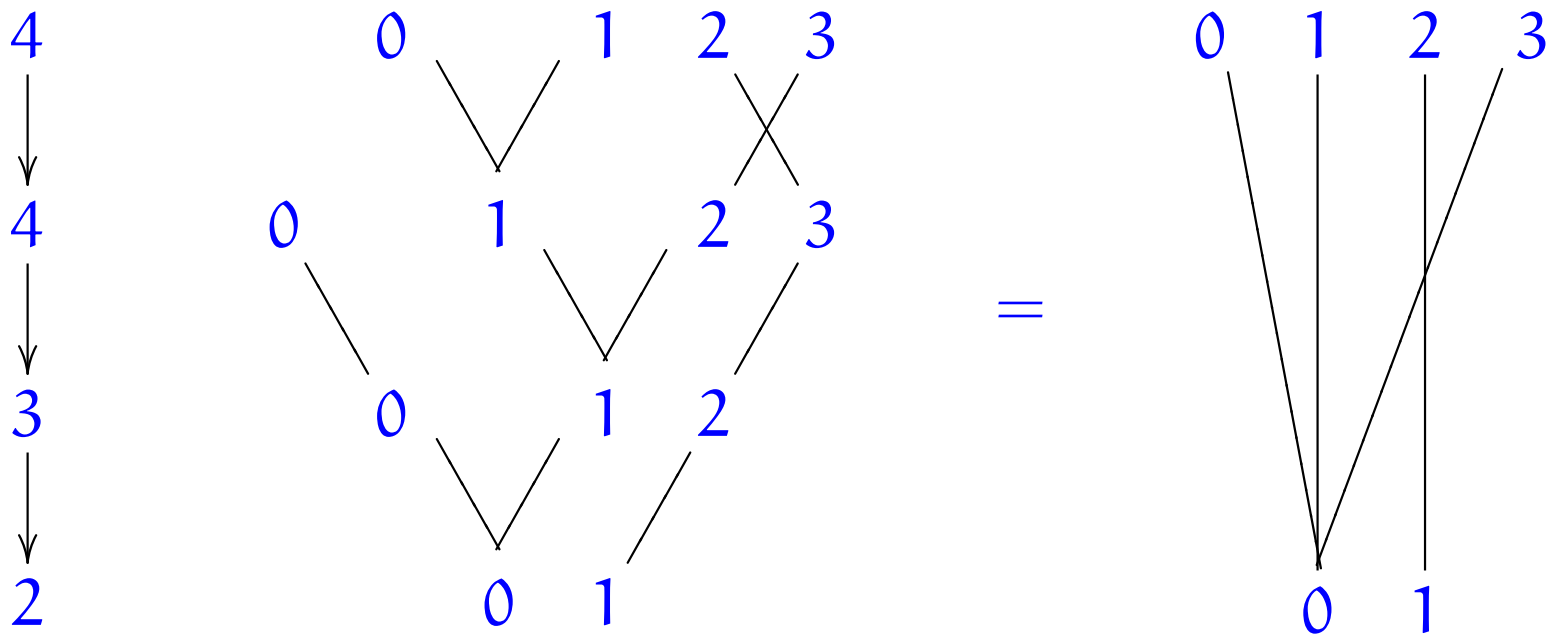
2. The free PROP  $\mathbf{P}[\bullet]$  on the equational presentation of a node  
 $\bullet : 1 \rightarrow 1$  is the free symmetric strict monoidal category on the  
 additive monoid of natural numbers, viz. the category  $\mathbf{Perm}_{\mathbb{N}}$   
 of *finite ordinals and  $\mathbb{N}$ -labelled permutations*.





3. The free PROP  $\mathbf{P}[\mathit{ComMon}]$  on the equational presentation of commutative monoids is the free cocartesian category on an object, i.e. the category  $\mathbf{Fun}$  of *finite ordinals and functions*.

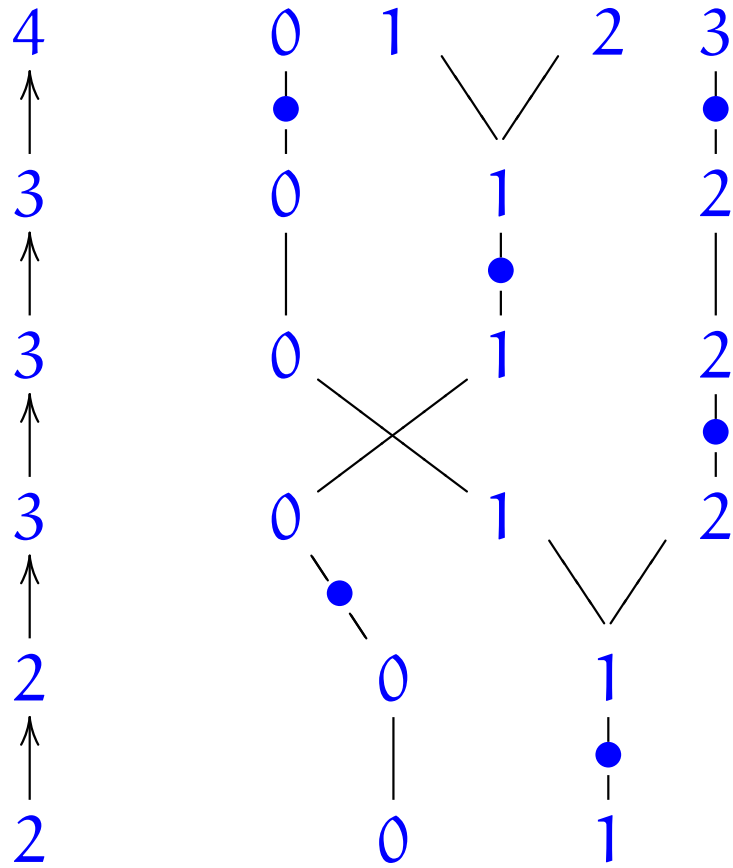
3. The free PROP  $\mathbf{P}[\mathit{ComMon}]$  on the equational presentation of commutative monoids is the free cocartesian category on an object, i.e. the category  $\mathbf{Fun}$  of *finite ordinals and functions*.



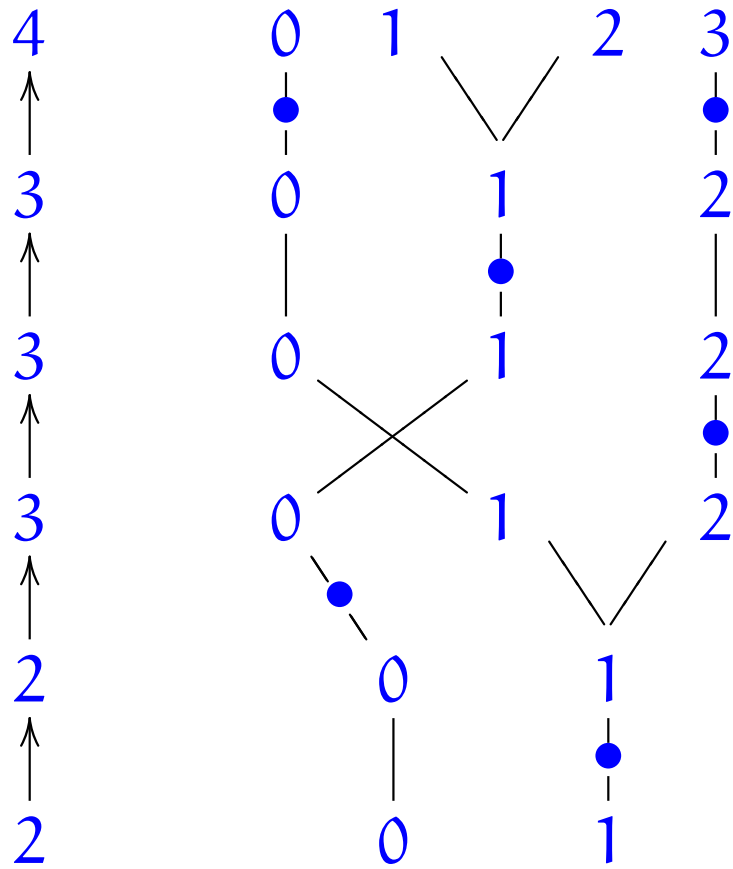
4. The free PROP  $\mathbf{P}[\bullet + \text{ComCoMon}]$  on the equational presentation of a node together with that of commutative comonoids is the subcategory **Forest** of **Dag** consisting of forests.

[Moerdijk, Milner]

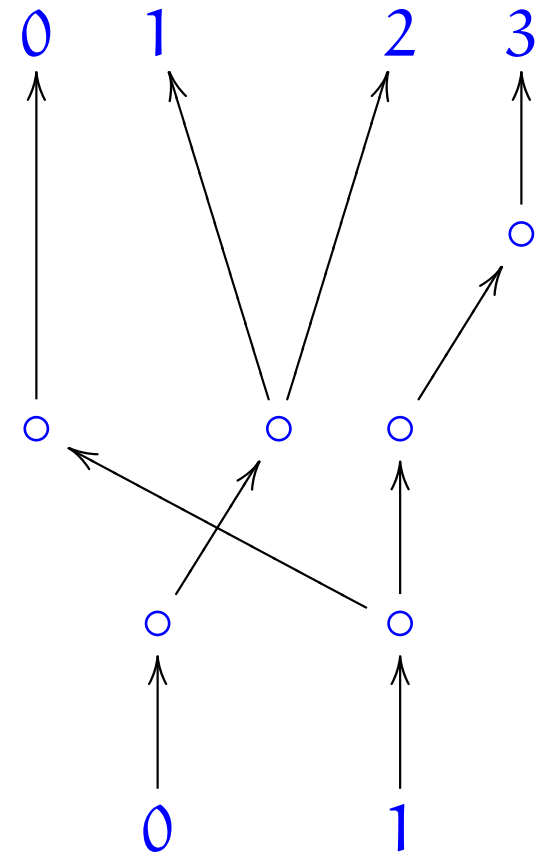
a morphism



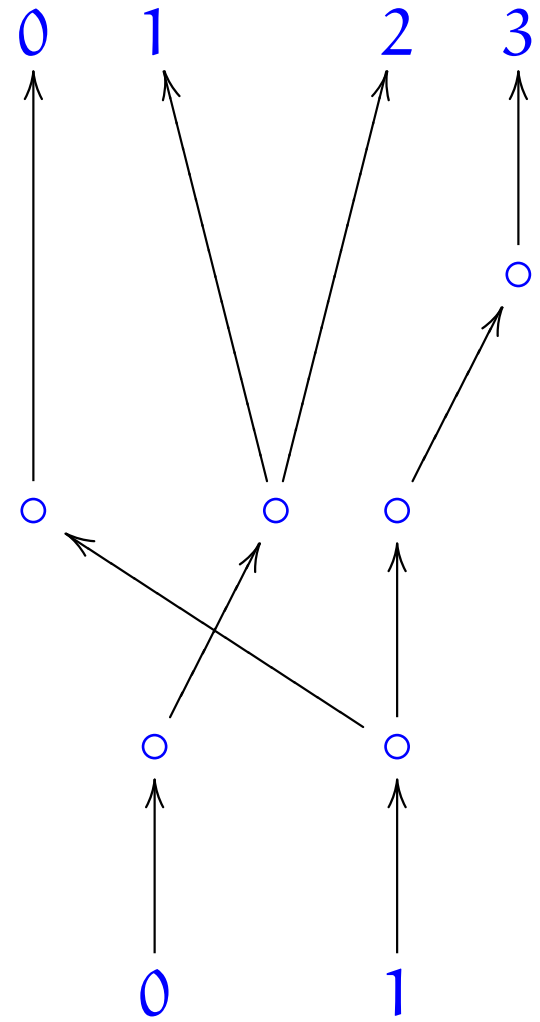
a morphism



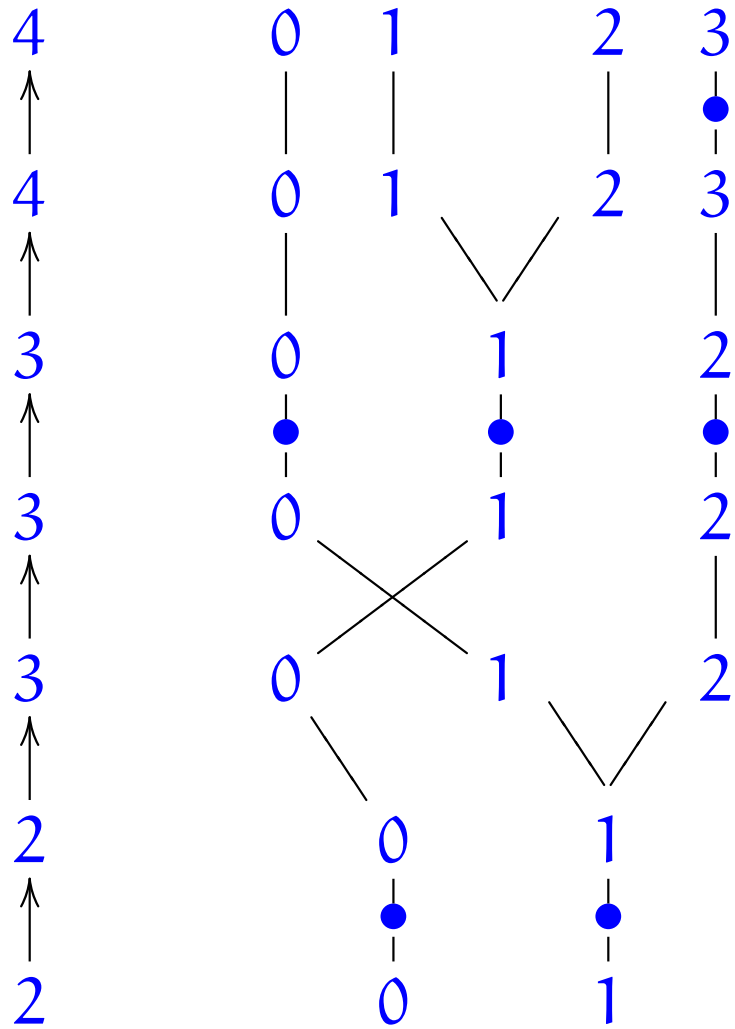
its forest representation



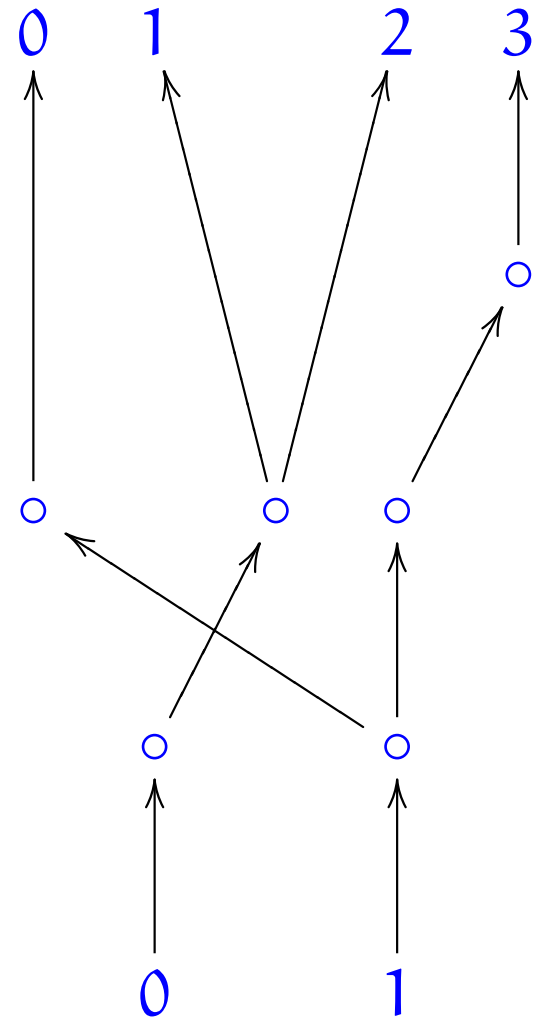
a forest



a layered normal form



a forest



5. The free PROP  $\mathbf{P}[\text{ComBiAlg}]$  on the equational presentation of commutative bialgebras is the free category with biproducts on an object, viz. the category  $\text{Mat}_{\mathbb{N}}$  of *finite ordinals and  $\mathbb{N}$ -valued matrices*.

[MacLane, Pirashvili, Lack]



5. The free PROP  $\mathbf{P}[ComBiAlg]$  on the equational presentation of commutative bialgebras is the free category with biproducts on an object, viz. the category  $\mathbf{Mat}_{\mathbb{N}}$  of *finite ordinals and  $\mathbb{N}$ -valued matrices*.

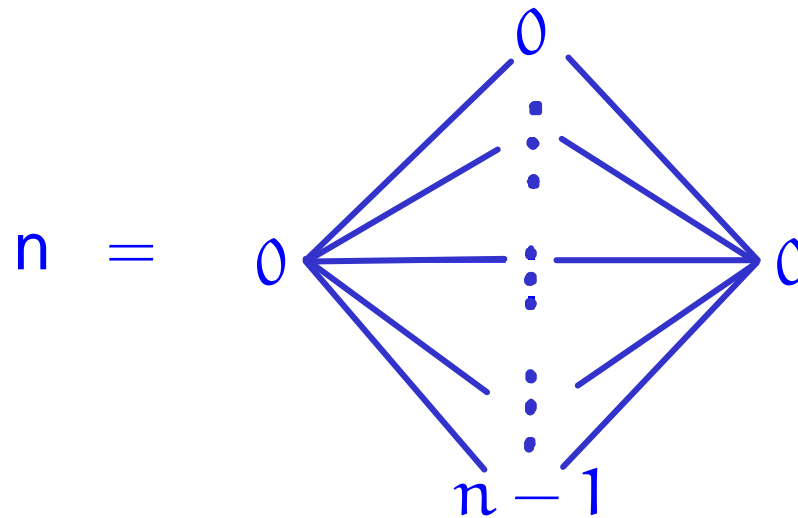
[MacLane, Pirashvili, Lack]

The equational presentation of commutative bialgebras is that of commutative monoids and commutative comonoids where the comonoid structure is a monoid homomorphism and the monoid structure is a comonoid homomorphism.

- (a) The commutative bialgebra structure turns the symmetric monoidal structure into biproduct structure.

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- (b) Every morphism  $m \rightarrow n$  has a unique representation as an  $m \times n$  matrix with entries in the endo-hom on  $1$ .

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- (b) Every morphism  $m \rightarrow n$  has a unique representation as an  $m \times n$  matrix with entries in the endo-hom on  $1$ .
- (c) The endo-hom on  $1$  is the multiplicative monoid of natural numbers.



6. The free PROP  $\mathbf{P}[\mathcal{D}eg\mathcal{C}om\mathcal{B}i\mathcal{A}lg]$  on the equational presentation of degenerate commutative bialgebras is the category  $\mathbf{Rel}$  of *finite ordinals and relations*.

The degeneracy axiom:

$$2 = 0 \begin{array}{c} 0 \\ \diamond \\ 1 \end{array} 0 = 0 - 0 = 1$$

7. The free PROP  $\mathbf{P}[\bullet + \mathcal{D}egComBiAlg]$  on the equational presentation of a node together with that of degenerate commutative bialgebras is **Dag**.

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**Proof:**

(a) Define *universal topological interpretations*  $[[D]]_\tau$  of dags  $D$  according to topological sortings  $\tau$  of  $D$ .

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- (a) Define *universal topological interpretations*  $[[D]]_{\tau}$  of dags  $D$  according to topological sortings  $\tau$  of  $D$ .
- (b) Prove the *invariance* of topological interpretations, viz. that  $[[D]]_{\tau} = [[D]]_{\tau'}$  for all topological sortings  $\tau$  and  $\tau'$  of  $D$ .



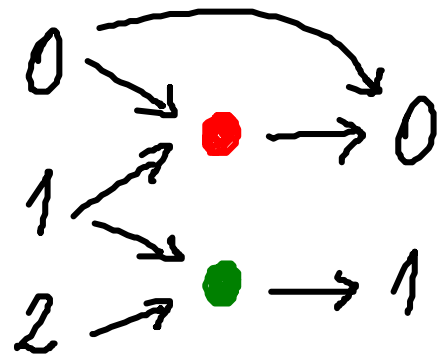
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**Proof:**

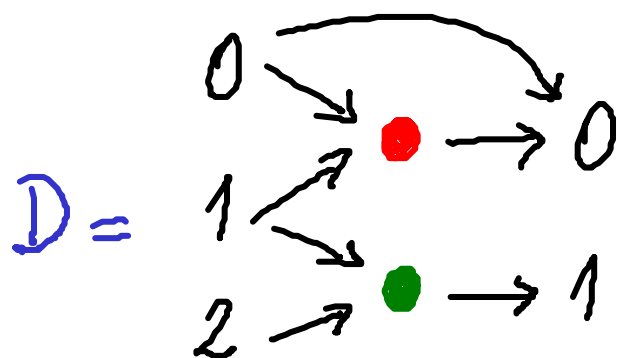
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- (b) Prove the *invariance* of topological interpretations, viz. that  $[[D]]_{\tau} = [[D]]_{\tau'}$  for all topological sortings  $\tau$  and  $\tau'$  of  $D$ .
- (c) Establish the *compositionality* of the interpretation function to obtain an *initial-algebra semantics*.

Example:

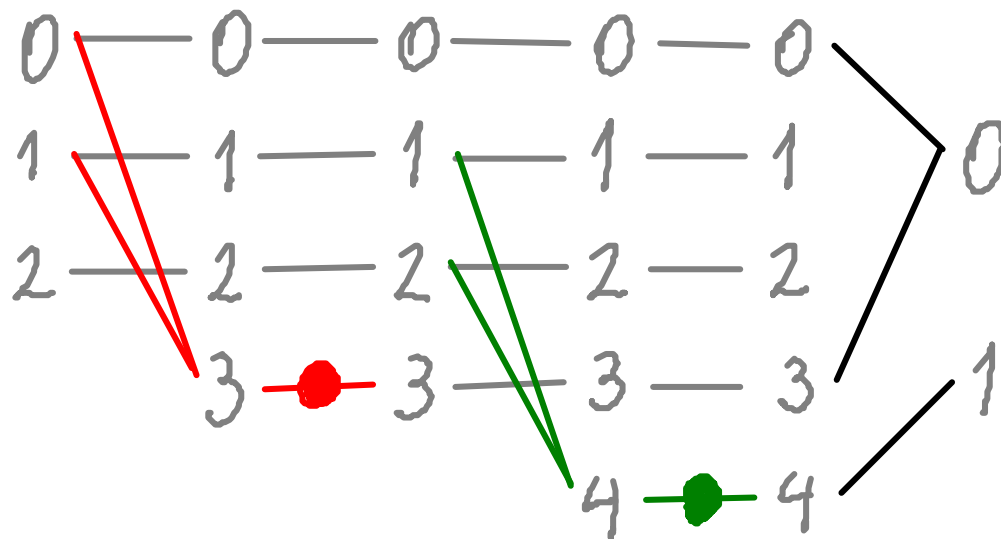
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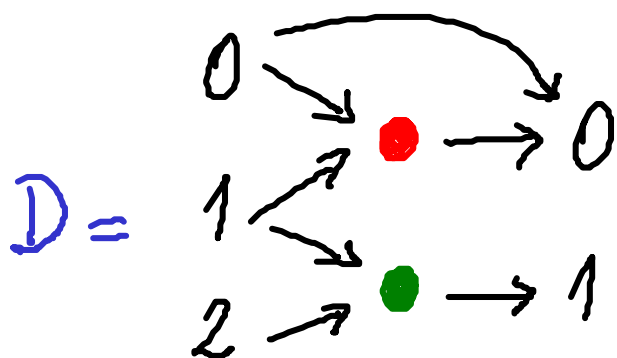
Example:



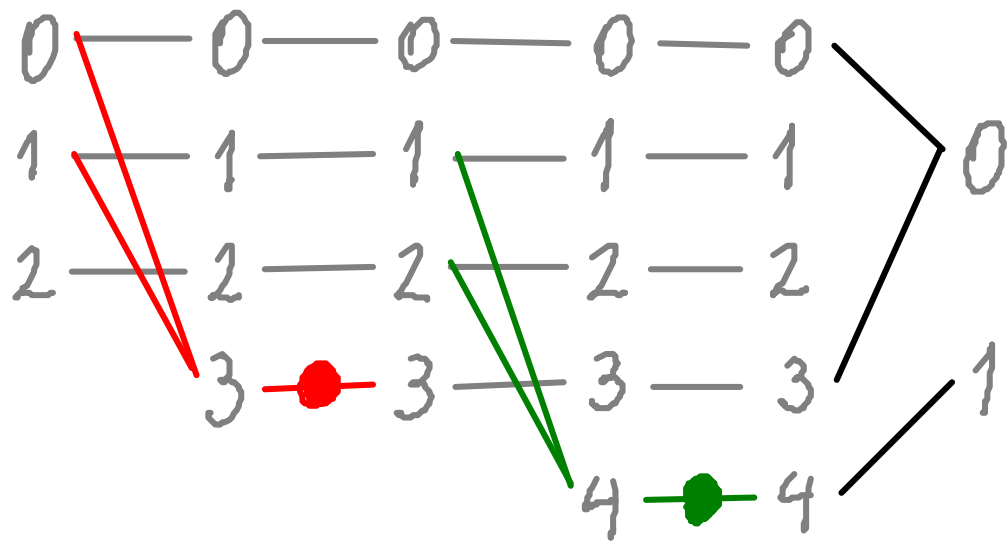
$[D]_{(\bullet < \bullet)}$



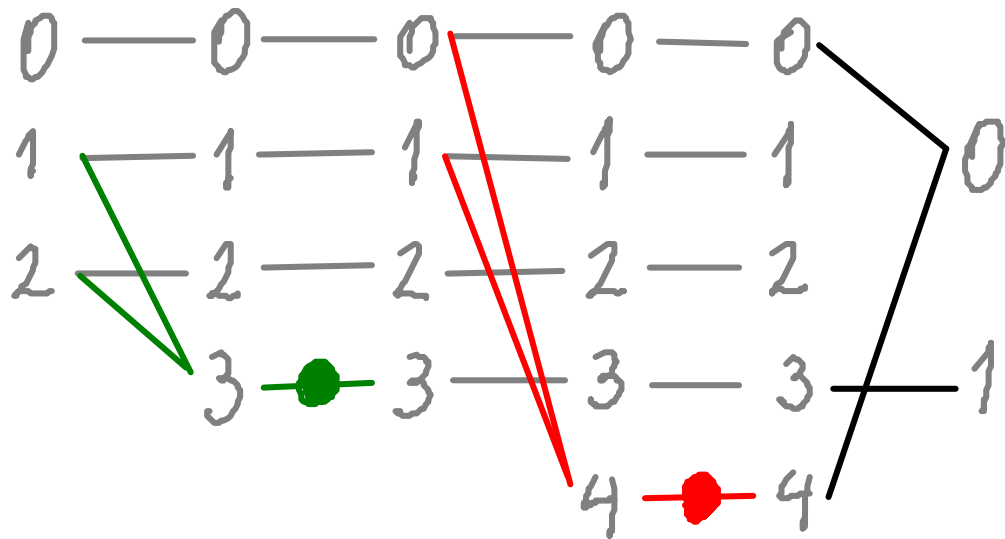
Example:



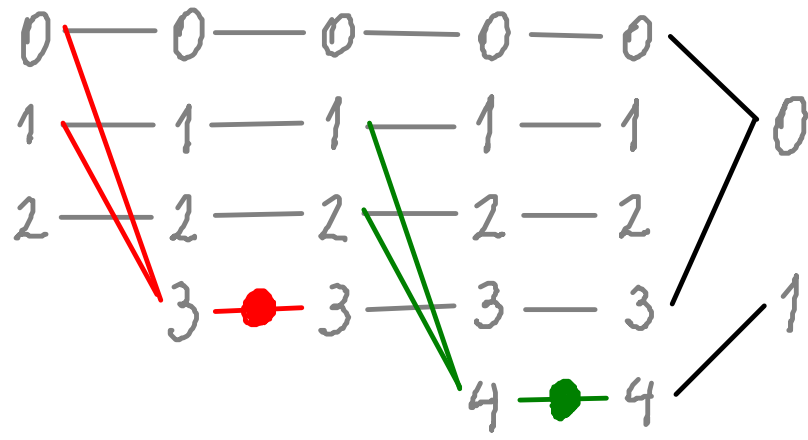
$[D]_{(\bullet < \bullet)}$



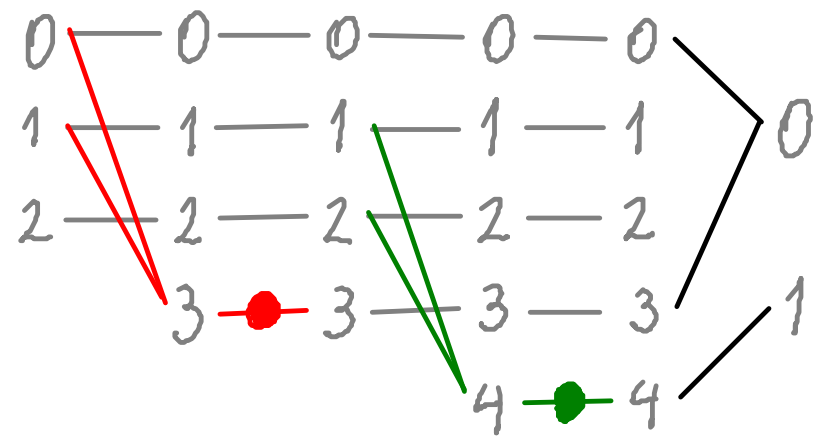
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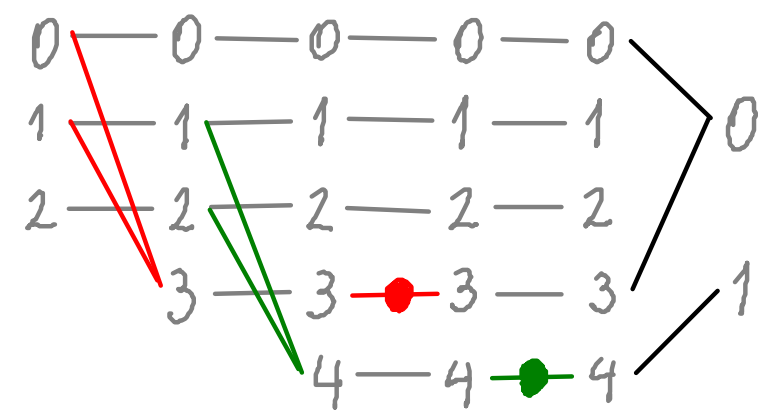
$[D] (\bullet < \bullet) =$

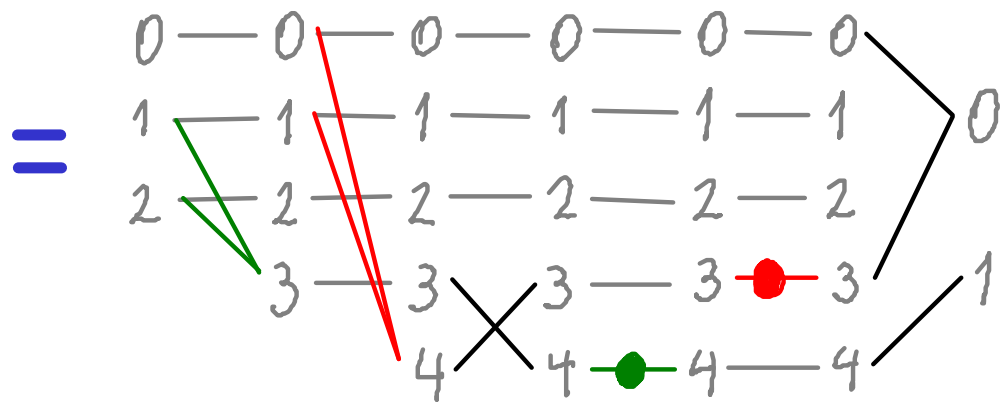
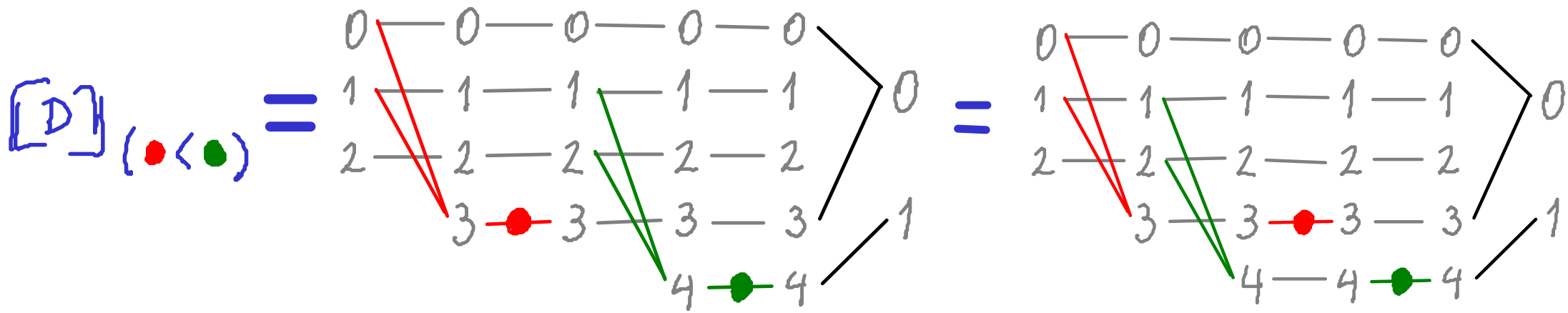


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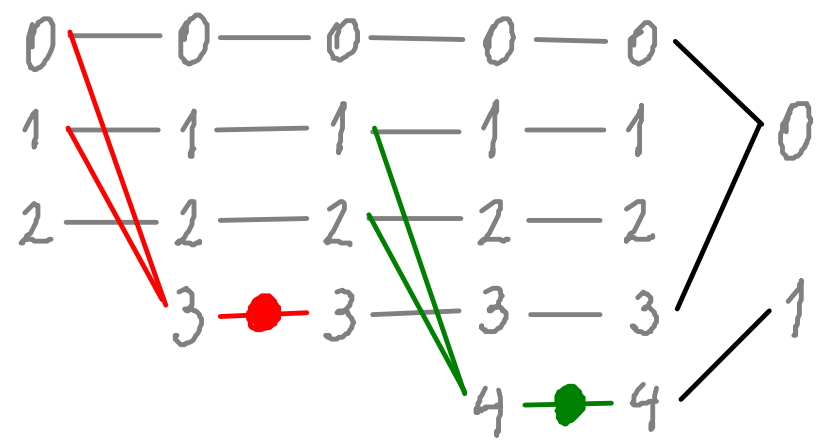
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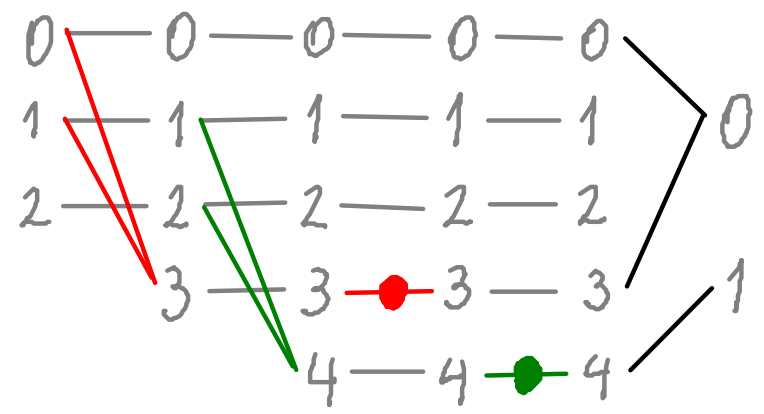


$[D] (\bullet < \bullet)$

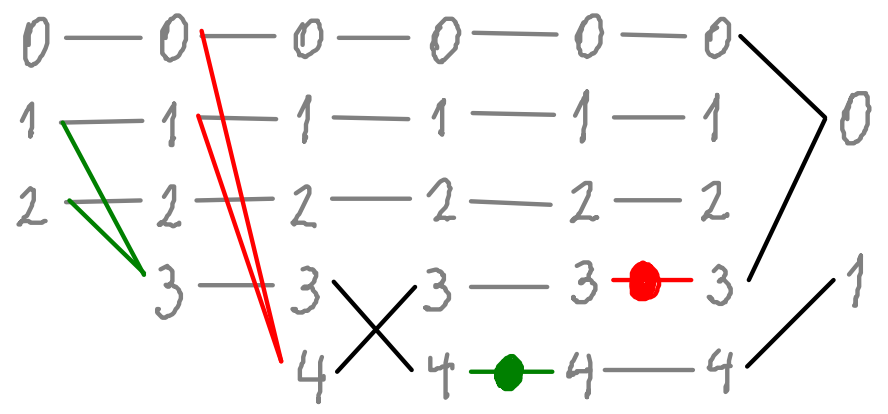
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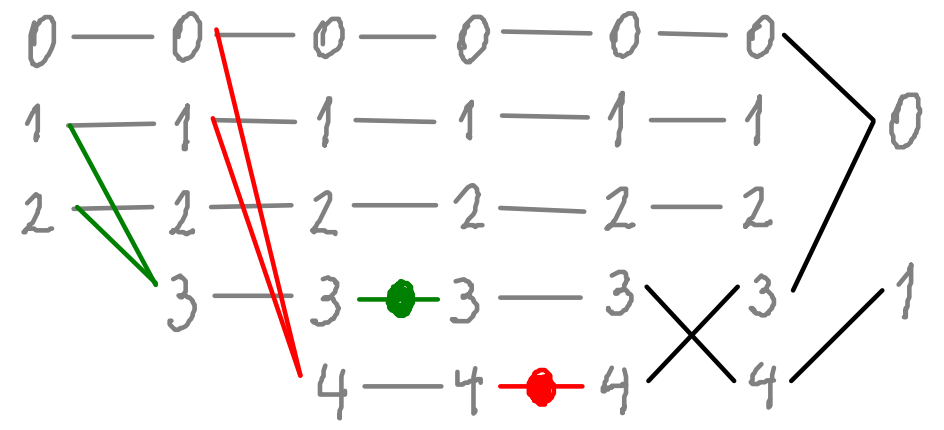
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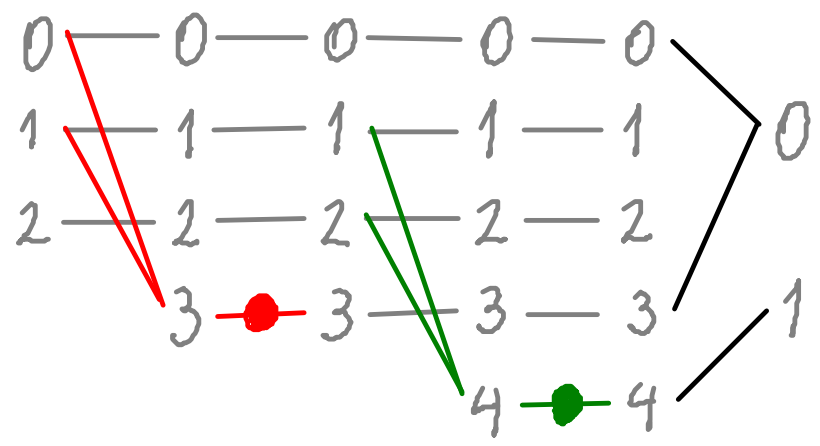
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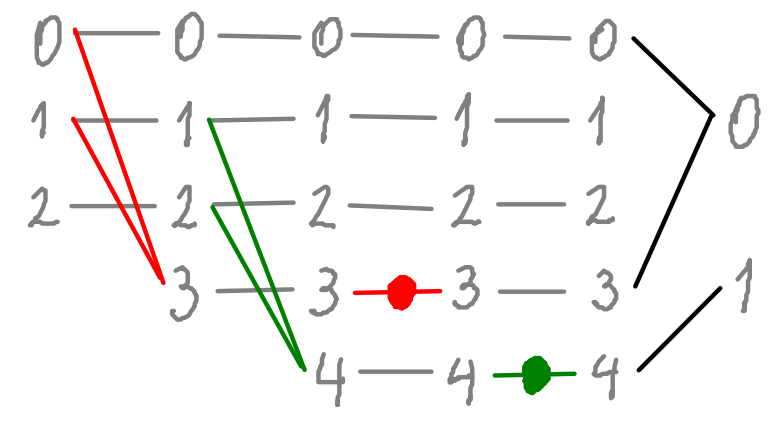


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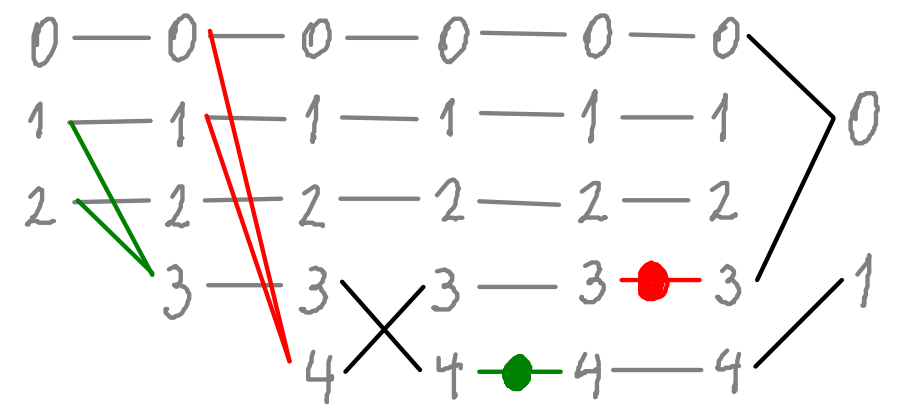
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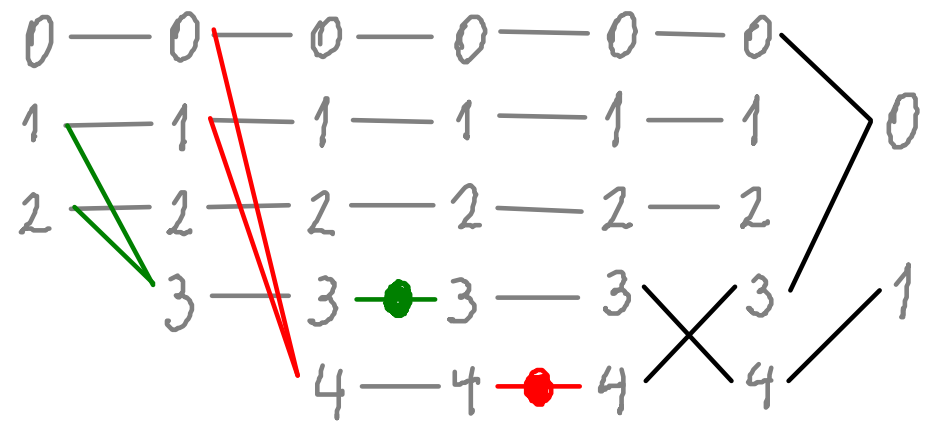
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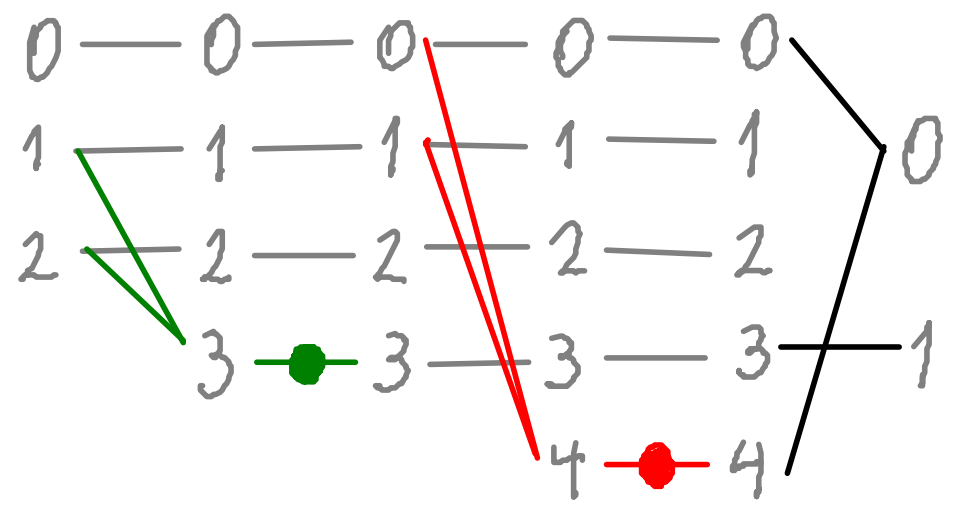
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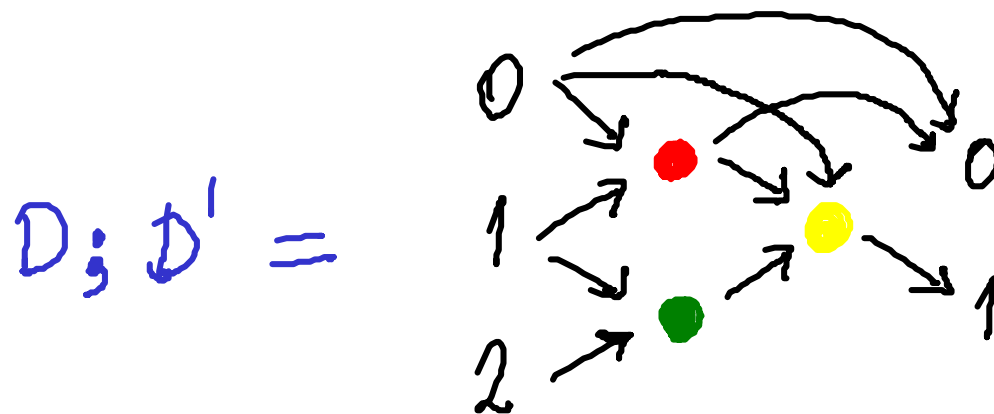
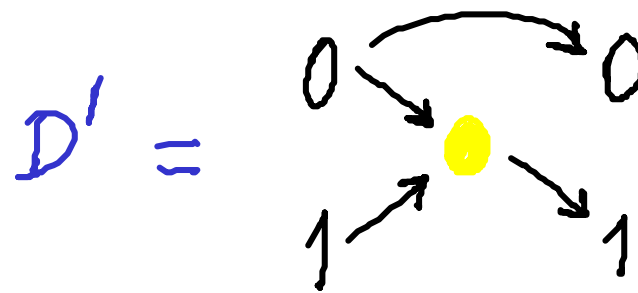
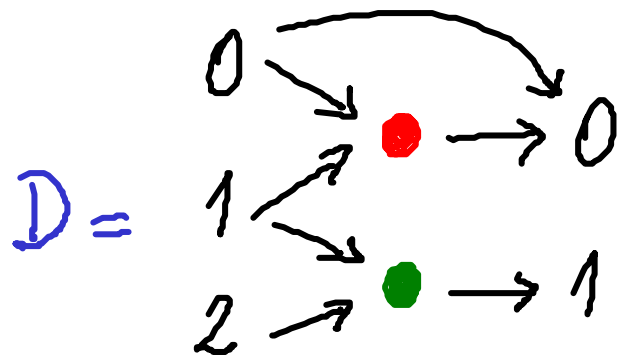


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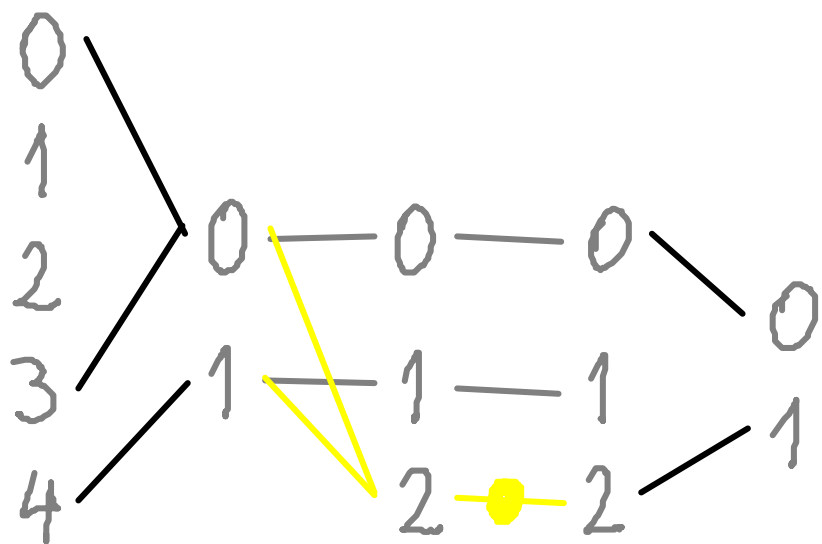


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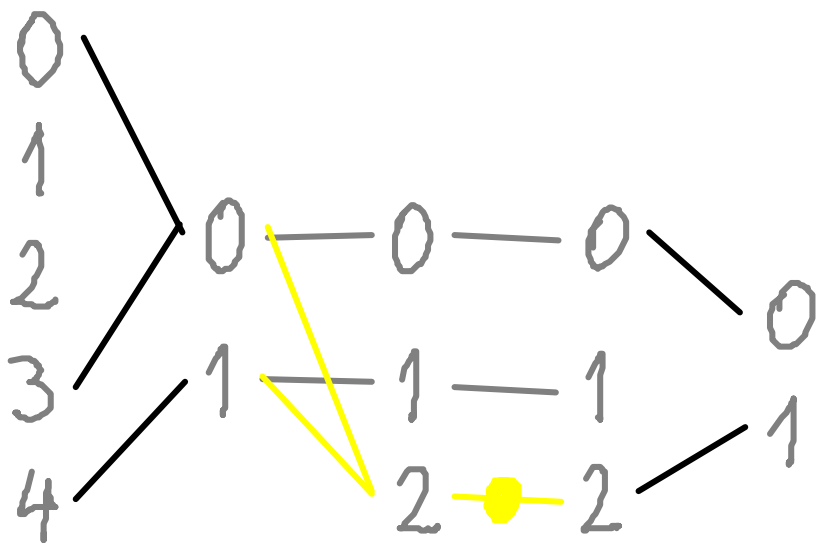
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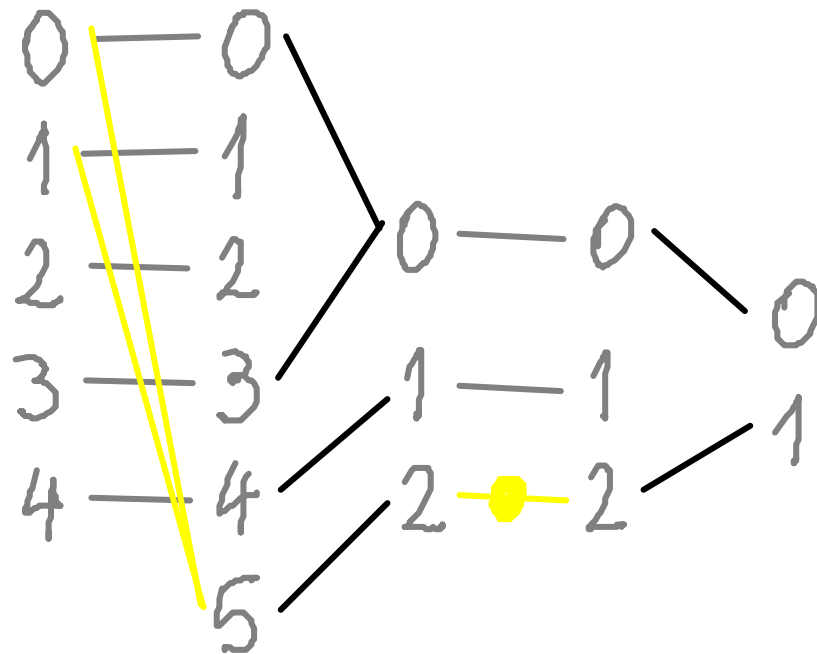
$$[D; D'](\bullet < \bullet < \bullet) = [D](\bullet < \bullet) ; [D'](\bullet)$$



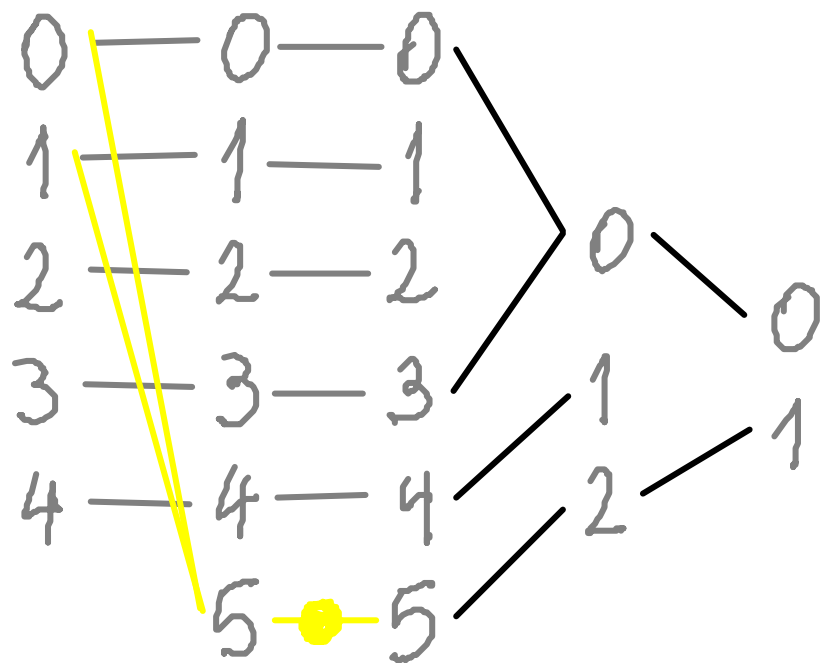


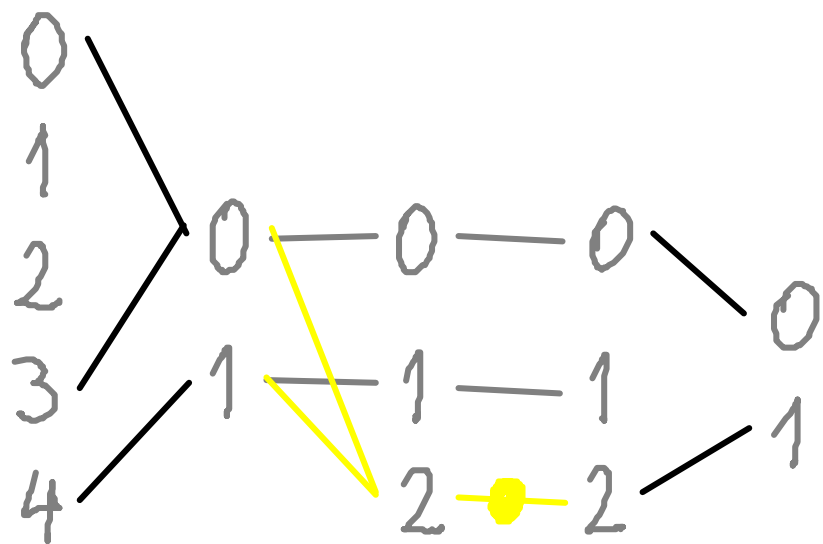


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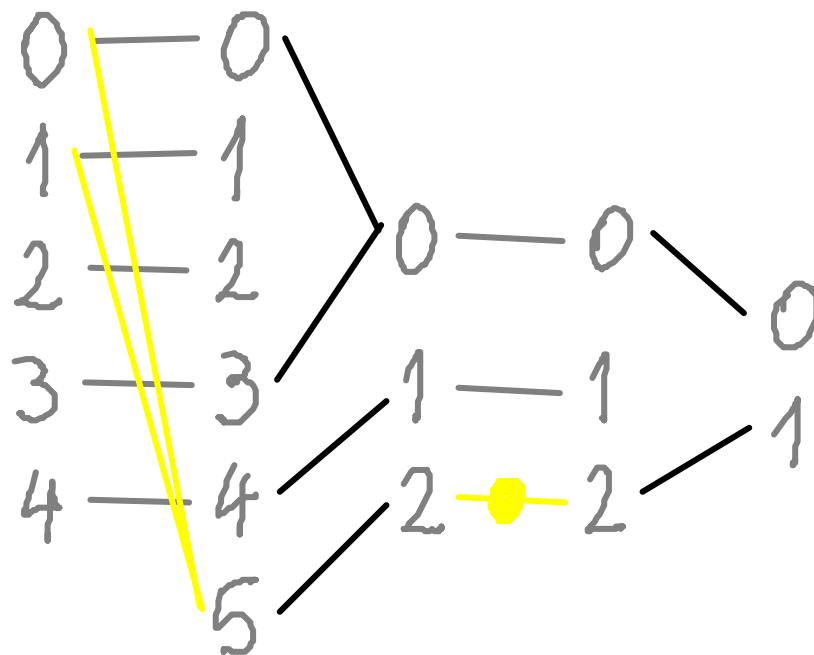


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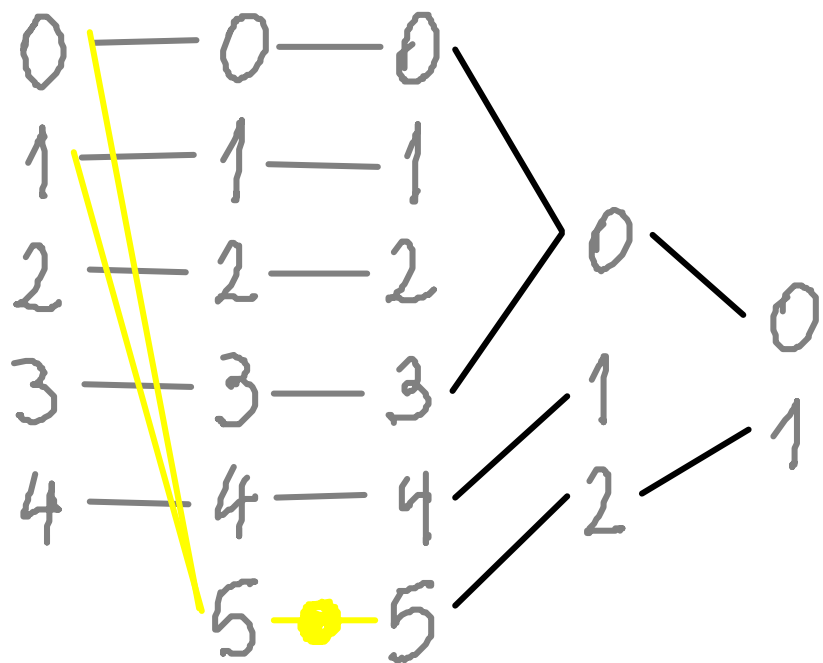




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