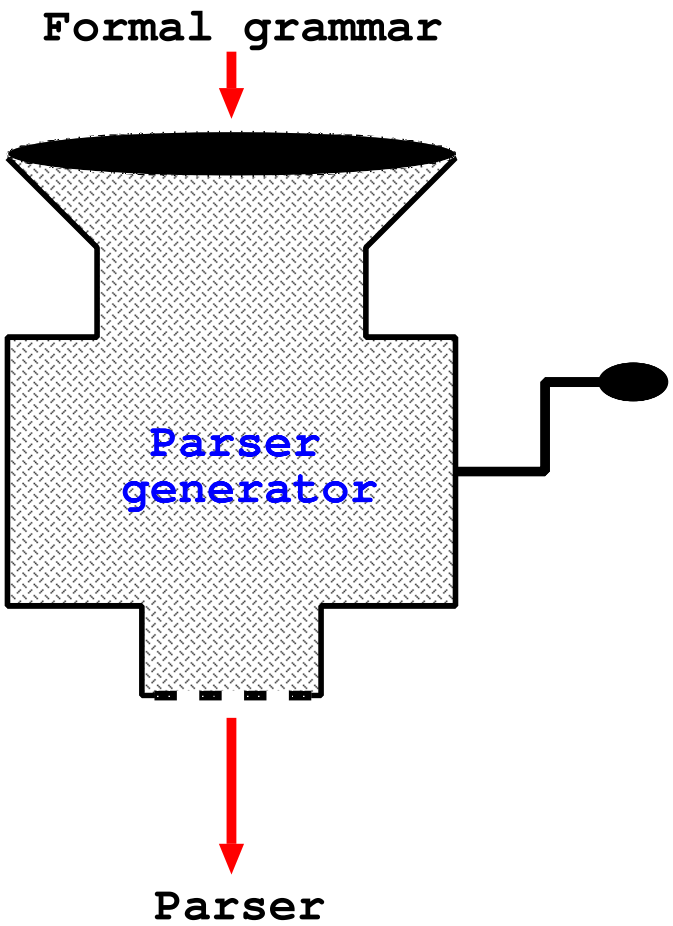
Executing the formal semantics of the Accellera Property Specification Language

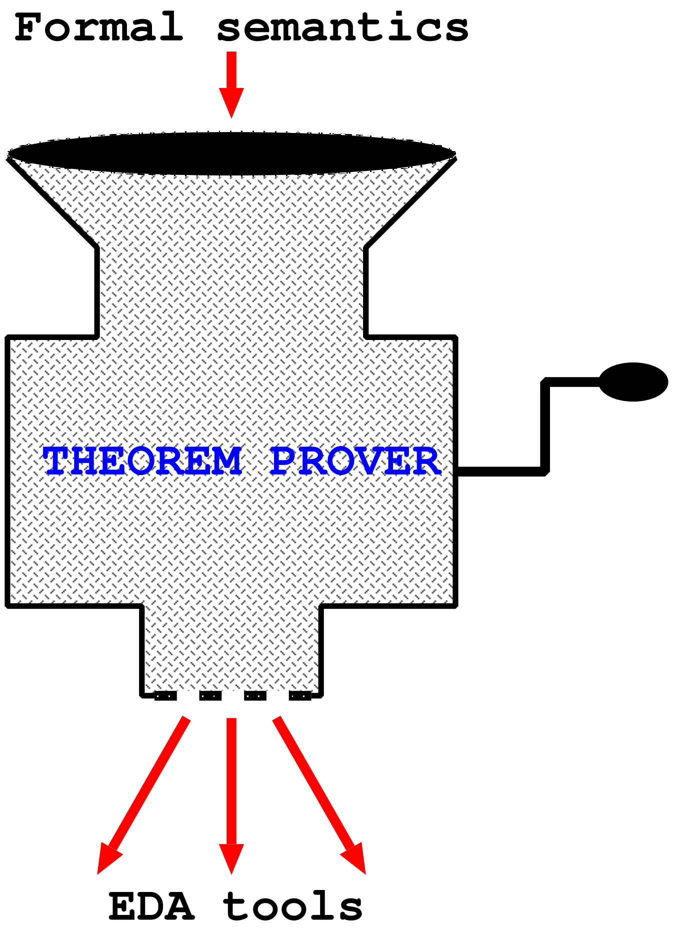
— joint work with Joe Hurd & Konrad Slind —

Standard practice: generate tools from formal syntax



- Input a grammar
- Apply theory of formal languages
- Generate a parser

Idea of this talk: generate tools from formal semantics



- Input ‘golden’ semantics from LRM
- Perform mechanised proof
- Generate tools

Goals and non-Goals

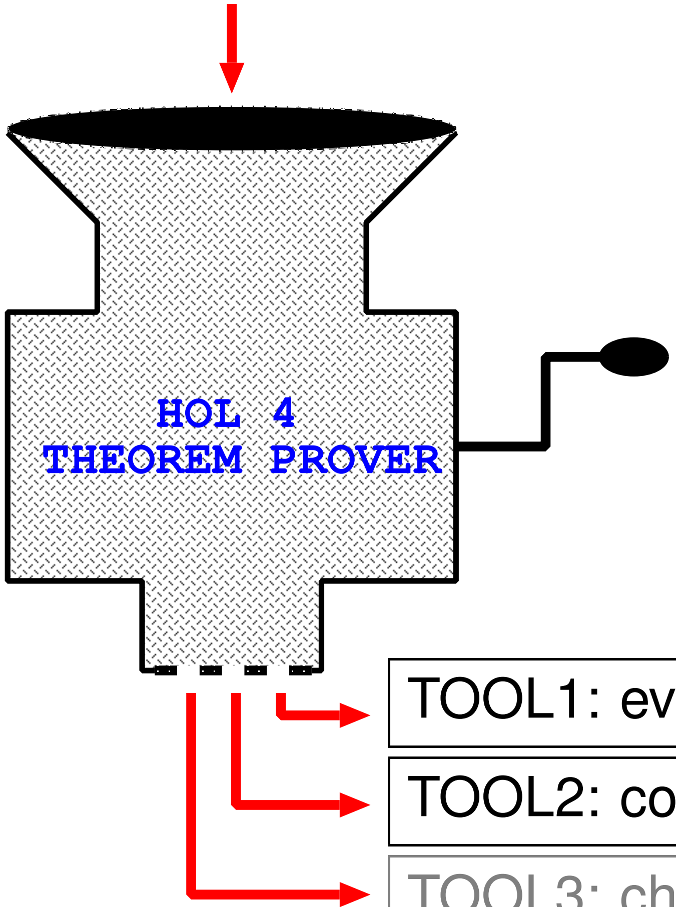
- Goal is to show formal semantics is not just documentation
  - can run the Language Reference Manual (LRM)
- Correctness primary, efficiency secondary
  - but need sufficient efficiency!
- Programming methodology, not new verification algorithms
  - EDA tools with theorem prover inside (*c.f.* PROSPER)

Accellera's PSL (formerly IBM's Sugar 2.0)

- PSL is a property specification language combining
  - boolean expressions (Verilog syntax)
  - patterns (Sequential Extended Regular Expressions SEREs)
  - LTL formulas (Foundation language FL)
  - CTL formulas (Optional Branching Extension OBE)
- Designed both for model checking and simulation testbenches
- Intended to be the industry standard

Generating PSL tools

Official semantics of PSL



TOOL1: evaluate FL properties on a specific path

TOOL2: compile properties to HDL checkers (idea from FoCs)

TOOL3: check OBE properties against a model (Amjad's PhD)

Tools use standard algorithms

- TOOL1: semantic calculator
  - match regexps using automata; evaluate formulas recursively
  - automata constructed and executed by proof inside HOL
- TOOL2: checker compiler
  - compile regexps to automata, then ‘pretty print’ to HDL (Verilog)
  - treatment of formulas incomplete and *ad hoc*
- TOOL3: symbolic model checker
  - classical McMillan-style μ-calculus checker
  - uses BDD representation judgements to link HOL terms to BDDs
  - see Gordon (TPHOLs2001), Amjad (TPHOLs2003)
- No new algorithms, but maybe a new kind of logic programming

Heroic proofs versus logic programming

- Theorem proving often associated with heroic proofs
  - *e.g.* Gödel's theorem (Shankar), relative consistency of AC (Paulson)
- We are not doing heroic proofs, but a kind of logic programming
  - computation by deduction
- HOL has a relatively fast call-by-value symbolic evaluator `EVAL`
  - by Bruno Barras using Coq technology (explicit substitutions)
  - doesn't compete with ACL2 or PVS ground evaluators (or C, C++)
  - runs ARM6 microarchitecture at a few seconds per instruction
  - key tool for our PSL evaluator

Executing the semantics (note: $(0 \mathinner{..} |w|)$ should be $[0 \mathinner{..} |w|)$)

- By rewriting and evaluation (PSL in red, HOL in blue):

  $\vdash\ s_0s_1s_2s_3s_4s_5s_6s_7s_8s_9 \models p \wedge X!\, f \;=\; s_0 \models p \;\wedge\; s_1s_2s_3s_4s_5s_6s_7s_8s_9 \models f$

  $\vdash\ \{a\}\{a,b\}\{b\} \models a \wedge X!\, b \;=\;$ T

- LRM semantics of the until-operator not directly executable

  $$w \models [f_1\ U\ f_2] \;=\; \exists\, k \in (0 \mathinner{..} |w|).\ w^k \models f_2 \wedge \forall\, j \in (0 \mathinner{..} k).\ w^j \models f_1$$

- Standard reformulation makes it directly executable

  $$\vdash\ w \models [f_1\ U\ f_2] \;=\; |w| > 0 \wedge (w \models f_2 \vee w \models f_1 \wedge w^1 \models [f_1\ U\ f_2])$$

- If f_1, f_2 are boolean expressions and the path is arbitrary of length 5:

  $$\begin{aligned} \vdash\ & s_0s_1s_2s_3s_4 \models [b_1\ U\ b_2] = \\ & s_0 \models b_2 \vee \\ & s_0 \models b_1 \wedge (s_1 \models b_2 \vee s_1 \models b_1 \wedge \\ & (s_2 \models b_2 \vee s_2 \models b_1 \wedge (s_3 \models b_2 \vee s_3 \models b_1 \wedge s_4 \models b_2))) \end{aligned}$$

Matching regular expressions

- Semantics of PSL SEREs is self-explanatory

$$
\begin{array}{lcll}
(w \models b & = & (|w| = 1) \wedge w_0 \models b) & \wedge \\
(w \models r_1;\ r_2 & = & \exists\, w_1 w_2.\ (w = w_1 w_2) \wedge w_1 \models r_1 \wedge w_2 \models r_2) & \wedge \\
(w \models r_1 : r_2 & = & \exists\, w_1 w_2 l.\ (w = w_1 [l] w_2) \wedge w_1 [l] \models r_1 \wedge [l] w_2 \models r_2) & \wedge \\
(w \models \{r_1\} \mid \{r_2\} & = & w \models r_1 \vee w \models r_2) & \wedge \\
(w \models \{r_1\}\&\&\{r_2\} & = & w \models r_1 \wedge w \models r_2) & \wedge \\
(w \models r[*] & = & \exists\, wlist.\ (w = \mathsf{Concat}\ wlist) \wedge \mathsf{Every}(\lambda\, w.\ w \models r)\, wlist) &
\end{array}
$$

- Make executable by proving

$$\vdash\ \ \forall\, w\ r.\ w \models r = \mathsf{amatch}(\mathsf{sere2regexp}(r))\, w$$

where:

- sere2regexp converts a SERE to a HOL regular expression
- amatch is an executable matcher for regular expressions

Suffix implication $\{r\}(f)$

- Semantics is:

  $w \models \{r\}(f) \;=\; \forall j \in (0 \,..\, |w|).\ w^{0,j} \models r \Rightarrow w^j \models f$

- Have defined an efficient executable function acheck so that, for example:

  $$\begin{array}{l} \mathsf{acheck}\ r\ f\ [x_0;\ x_1;\ x_2;\ x_3] = \\ \quad (\mathsf{amatch}\ r\ [x_0] \Rightarrow f[x_0;\ x_1;\ x_2;\ x_3]) \wedge \\ \quad (\mathsf{amatch}\ r\ [x_0;\ x_1] \Rightarrow f[x_1;\ x_2;\ x_3]) \wedge \\ \quad (\mathsf{amatch}\ r\ [x_0;\ x_1;\ x_2] \Rightarrow f[x_2;\ x_3]) \wedge \\ \quad (\mathsf{amatch}\ r\ [x_0;\ x_1;\ x_2;\ x_3] \Rightarrow f[x_3]) \end{array}$$

  N.B. execution only costs the same as the last amatch call

- Then proved

  $\vdash\ \forall w\ r\ f.\ w \models \{r\}(f) \;=\; \mathsf{acheck}(\mathsf{sere2regexp}(r))(\lambda x.\ x \models f)w$

- Rewrite with this, then execute

Other constructs: $\{r_1\} \mapsto \{r_2\}!$ and $\{r_1\} \mapsto \{r_2\}$

- Reduce $\{r_1\} \mapsto \{r_2\}!$ to suffix implication by proving

  $\vdash \ \forall w\ r_1\ r_2.\ w \models \{r_1\} \mapsto \{r_2\}! = w \models \{r_1\}(\neg\{r_2\}(\mathsf{F}))$

- Handle $\{r_1\} \mapsto \{r_2\}$ with regular expression $\mathsf{Prefix}(r)$ (in HOL not PSL)

  $\vdash \ \forall r\ w.\ w \models \mathsf{Prefix}(r) \ = \ \exists w'.\ w\ w' \models r$

- Execution of $w \models \mathsf{Prefix}(r)$ uses Dijkstra's algorithm

- Have proved:

  $$\begin{array}{l}\vdash \ \forall w\ r_1\ r_2.\\ \quad w \models \{r_1\} \mapsto \{r_2\} =\\ \quad \mathsf{acheck}(\mathsf{sere2regexp}\ r_1)\\ \qquad\qquad (\lambda x.\ x \models \neg\{r_2\}(\mathsf{F}) \vee \mathsf{amatch}\ (\mathsf{Prefix}\ (\mathsf{sere2regexp}\ r_2))\ x)\ w\end{array}$$

- Rewrite with this, then execute

Remaining formulas: aborts and clocking

- Semantics of abort formulas needs a reachability algorithm
  - have implemented a partial method
  - awaiting new abort semantics before attempting complete solution
- Clocked formulas f@c, f@c! can be translated to unclocked formulas
  - translation to unclocked formulas is by a recursive function
  - can be directly executed
  - have proved a theorem that says (roughly):

$$\vdash\ \forall r.\ \textit{unclocked\_semantics(translation } r\textit{)} = \textit{clocked\_semantics(}r\textit{)}$$

Example

- PSL Reference Manual Example 2, page 45

```
time  0  1  2  3  4  5  6  7  8  9
-----------------------------------
clk1  0  1  0  1  0  1  0  1  0  1
a     0  0  0  1  1  1  0  0  0  0
b     0  0  0  0  0  1  0  1  1  0
c     1  0  0  0  0  1  1  0  0  0
clk2  1  0  0  1  0  0  1  0  0  1
```

- Define w to be this path, so w is:
  `{c,clk2}{clk1}{}{clk1,a,clk2}{a}{clk1,a,b,c}{c,clk2}{clk1,b}{b}{clk1,clk2}`

- Can evaluate in SML, or via a command line wrapper

- Example: to evaluate `(c && next!(a until b))@clk1` at all times in w:

```
% pslcheck -all \
    -fl   '(c && next!(a until b))@clk1' \
    -path '{c,clk2}{clk1}{}{clk1,a,clk2}{a}{clk1,a,b,c}{c,clk2}{clk1,b}{b}{clk1,clk2}'
> > true at times 4,5,10
```

Uses of TOOL1 (calculating $w \models^{\top} f$ from semantics)

- Teaching and learning tool for exploring semantics
- Checking one has the right property before using it in verification
- Post simulation analysis (path is generated by simulator)
  - compare with “TransEDA VN-Property” property checker and analyzer
  - our tools much slower – but not necessary too slow!
  - guaranteed PSL compliant by construction: golden reference

TOOL2: Compile the semantics to checkers

- Idea pinched from IBM FoCs project
- A defined operator: $\forall\, r.\ \mathit{never}(r) \ = \ \{\mathrm{T}[*];\ r\} \mapsto \{\mathsf{F}\}$
- Example property: $\mathit{never}(\neg\mathtt{StoB\_REQ} \wedge \mathtt{BtoS\_ACK};\ \mathtt{StoB\_REQ})$
- Use semantics to generate a Verilog checker

```
module Checker (StoB_REQ, BtoS_ACK, BtoR_REQ, RtoB_ACK);

input StoB_REQ, BtoS_ACK, BtoR_REQ, RtoB_ACK;
reg   [1:0] state;

initial state = 0;

always @ (StoB_REQ or BtoS_ACK or BtoR_REQ or RtoB_ACK)
begin
 $display ("Checker: state = %0d", state);
 case (state)
  0: if (StoB_REQ) state = 1; else if (BtoS_ACK) state = 2; else state = 1;
  1: if (StoB_REQ) state = 1; else if (BtoS_ACK) state = 2; else state = 1;
  2: if (StoB_REQ) state = 3; else if (BtoS_ACK) state = 2; else state = 1;
  3: begin $display ("Checker: property violated!"); $finish; end
  default: begin $display ("Checker: unknown state"); $finish; end
 endcase
end

endmodule
```

Example of how the checker works and is justified

- The following theorem is first proved

  $\vdash \ |w| = \infty \ \Rightarrow \ w \models \mathit{never}(r) = \forall\, n.\ \neg\mathsf{amatch}\ (\mathsf{sere2regexp}\ \mathrm{T}[*];\ r)(w^{0,n})$

- Thus there's an error if $\mathsf{amatch}\ (\mathsf{sere2regexp}\ \mathrm{T}[*];\ r)(w^{0,n})$ is ever true
- Generate a DFA from $\mathsf{sere2regexp}\ \mathrm{T}[*];\ r$
- So far everything is by proof, so correct by construction
- Final step is to pretty print checker into HDL (Verilog)
  - this may introduce errors
  - no formal semantics of Verilog `:-(`
- Only have 'proof of concept' for checkers: more work to cover all formulas

Conclusions

- Two tools: semantic calculator and checker generator
- Correct by construction
- More work needed (especially for checkers)
- Illustrates new kind of logic programming using a theorem prover
  - prototyping standards compliant tools
  - theorem proving is slow .... but not necessarily too slow
  - maybe OK for some industrial strength *performance-non-critical* tools

THE END