Green-Marl: A DSL for Easy and Efficient Graph Analysis

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Problem

- Difficult to execute graph analysis on large dataset
- Current graph processing frameworks force user to rewrite their program

The paper offers a Domain Specific Language (DSL)

- Easy to express graph algorithms
- Expose data-level parallelism
- Can compile to various backends

Betweenness Centrality algorithm in Green-Marl

```
Procedure Compute_BC(
     G: Graph, BC: Node_Prop<Float>(G)) {
 2
 3
      G.BC = 0; // initialize BC
      Foreach(s: G.Nodes) {
 4
 5
       // define temporary properties
6
        Node_Prop<Float>(G) Sigma;
 7
        Node_Prop<Float>(G) Delta;
8
        s.Sigma = 1; // Initialize Sigma for root
9
       // Traverse graph in BFS-order from s
10
        InBFS(v: G.Nodes From s)(v!=s) {
11
          // sum over BFS-parents
12
          v.Sigma = Sum(w: v.UpNbrs) {w.Sigma};
13
14
       // Traverse graph in reverse BFS-order
15
        InRBFS(v!=s) {
16
          // sum over BFS-children
17
          v.Delta = Sum (w:v.DownNbrs) {
18
             v.Sigma / w.Sigma * (1+ w.Delta)
19
          };
20
          v.BC += v.Delta @s; //accumulate BC
21
```

Scope of the language

- Graph is an ordered pair of nodes and edges G = (N, E)
- Each node/edge has some properties
- Given a graph and a set of properties Green Marl can:
 - Compute scalar value
 - Compute a new property
 - Select subgraph from original graph

Language Construct

- Data types 5 primitive types
 Int, bool, string, double and float
- Collections
 - Set, Order, Sequence

Semantics on collections

		sequential		parallel		1	
Group	Op-Name	S	O	Q	S	O	Q
Grow	Add	V			V		
	Push(Front/Back)		v	V		V	v
Shrink	Remove	V			V		
	Pop(Front/Back)		v	V		?	?
	Clear	v	v	V	v	v	V
Lookup	Has	V	v	V	V	V	V
	Front(Back)		V	V	8	V	V
	Size	v	v	V	v	V	V
Сору	=	V	v	V	X	Х	Х
Iteration	Items	v	v	V	V	v	V
Modification under iteration \rightarrow Shrink, Grow, or Copy: X							
Conflicts under parallel execution \rightarrow							
Grow-Shrink: X Lookup-Shrink: ? Lookup-Grow: ?							

Iterations and Traversals

- For
- Foreach
- Breadth-First Search
- Depth-First Search

Deferred Assignment

```
52 Foreach(s:G.Nodes) {
53     // no conflict. t.X gives 'old' value
54     s.X <= Sum(t:s.Nbrs) {t.X} @ s
55 }
56 // All the writes to X becomes visible simultaneously
57 // at the end of the s iteration.</pre>
```

Reductions

58	Int x,y;	In-place	Assignment	In-place	Assignment
59	$x = Sum(t:G.Nodes) \{t.A\};$	All	a &=	Sum	+=
60 61	y = 0; Foreach († : G. Nodes)	Any	=	Product	*=
62	y+= t.A;	Min	min=	Count	++
	_	Max	max=		

Overview of Green-Marl DSL-compiler Usage



Some optimizations

Set graph loop fusion

Loop fusion

```
103 Foreach(s: G.Nodes)(f(s))
104 s.A = X(s.B);
105 Foreach(t: G.Nodes)(g(t))
106 t.B = Y(t.A)
```

becomes

```
139 Node_Set S(G); // ...
140 Foreach(s: S.Items)
141 s.A = x(s.B);
142 Foreach(t: G.Nodes)(g(t))
143 t.B = y(t.A)
```

becomes

Name	LOC Original	LOC Green-Marl	Source
BC	350	24	[9] (C OpenMp)
Conductance	42	10	[9] (C OpenMp)
Vetex Cover	71	25	[9] (C OpenMp)
PageRank	58	15	[2] (C++, sequential)
SCC(Kosaraju)	80	15	[3] (Java, sequential)

Lines of code

Betweenness Centrality



Conductance



Vertex Cover



Page Rank



222 Foreach(s:G.Nodes)
223 For(t: s.Nbrs)
224 s.A = s.A + t.B;

becomes

```
225
     OMP (parallel for)
226
    for(index_t s = 0; s < G.numNodes(); s++) {</pre>
227
      // iterate over node's edges
228
      for(index_t t_=G.edge_idx[s]:t_<G.edge_idx[s+1];t_++) {</pre>
229
     // get node from the edge
230
     index_t t = G.node_idx[t];
231
     A[s] = A[s] + B[t];
232
   } }
```

Code Generation

Critique of Paper

- Like
 - Easy way to process graph
 - Increased programmer programmability
 - Architecture portability
- Dislike
 - Graph is immutable during the analysis
 - Syntax can take a bit of getting used to.
 - Only compares to SNAP