

Canonical gradings of monads

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Example

The writer monad Wr for lists over a set C has:

object mapping $\text{Wr} : \text{Set} \rightarrow \text{Set}$ $\text{Wr } X = \text{List } C \times X$

unit functions $\eta_X : X \rightarrow \text{Wr } X$ $\eta_X x = ([], x)$

multiplication functions $\mu_X : \text{Wr} (\text{Wr } X) \rightarrow \text{Wr } X$ $\mu_X (s_1, (s_2, x)) = (s_1 \# s_2, x)$

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We can **grade** this by

- ▶ natural numbers $e \in \mathbb{N}$, to get a graded monad WrL :

$\text{WrL } e X = \text{List}_{\leq e} C \times X$ $\eta : X \rightarrow \text{WrL } 0 X$ $\mu : \text{WrL } e_1 (\text{WrL } e_2 X) \rightarrow \text{WrL} (e_1 + e_2) X$

- ▶ subsets $e \subseteq C$, to get a graded monad WrS :

$\text{WrS } e X = \text{List } e \times X$ $\eta : X \rightarrow \text{WrS } \emptyset X$ $\mu : \text{WrS } e_1 (\text{WrS } e_2 X) \rightarrow \text{WrS} (e_1 \cup e_2) X$

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multiplication functions $\mu_X : \text{Wr} (\text{Wr } X) \rightarrow \text{Wr } X$ $\mu_X (s_1, (s_2, x)) = (s_1 \# s_2, x)$

We can **grade** this by

- ▶ subsets $\Sigma \subseteq \text{List } C$, to get a graded monad WrC :

$\text{WrC } \Sigma X = \Sigma \times X$ $\eta : X \rightarrow \text{WrC } J X$ $\mu : \text{WrC } \Sigma_1 (\text{WrC } \Sigma_2 X) \rightarrow \text{WrC} (\Sigma_1 \boxplus \Sigma_2) X$

where

$$J = \{[]\} \quad \Sigma_1 \boxplus \Sigma_2 = \{s_1 \# s_2 \mid s_1 \in \Sigma_1, s_2 \in \Sigma_2\}$$

Example

WrC is the canonical grading of Wr:

- ▶ WrL is

$$\mathbb{N} \xrightarrow{F} \mathcal{P}(\text{List } C) \xrightarrow{\text{WrC}} [\text{Set}, \text{Set}]$$

where

$$Fe = \text{List}_{\leq e} C \subseteq \text{List } C$$

- ▶ WrS is

$$\mathcal{P}C \xrightarrow{F} \mathcal{P}(\text{List } C) \xrightarrow{\text{WrC}} [\text{Set}, \text{Set}]$$

where

$$Fe = \text{List } e \subseteq \text{List } C$$

This work

More generally: given

- ▶ a (skew) monoidal category \mathbf{D} (e.g. $[\mathbf{Set}, \mathbf{Set}]$)
- ▶ a class \mathcal{M} of \mathbf{D} -morphisms (e.g. componentwise injective natural transformations)
- ▶ a monoid T in \mathbf{D} (e.g. any monad on \mathbf{Set})

satisfying some reasonable conditions, we have

- ▶ a notion of \mathcal{M} -grading of T
- ▶ T has a canonical \mathcal{M} -grading
- ▶ every other \mathcal{M} -grading factors through the canonical one

Grading objects

Given

- ▶ a category \mathbf{D}
- ▶ a class of \mathbf{D} -morphisms \mathcal{M}
- ▶ an object $T \in \mathbf{D}$

an \mathcal{M} -grading (\mathbf{G}, G, g) of T is:

- ▶ a category \mathbf{G} of **grades**
- ▶ with a functor $G : \mathbf{G} \rightarrow \mathbf{D}$
- ▶ and a natural transformation

$$g_e : Ge \rightarrow T$$

whose components are in \mathcal{M}

Example: writer monad

$$\mathbf{D} = [\mathbf{Set}, \mathbf{Set}]$$

\mathcal{M} = componentwise injective nat. transformations

$$T = \text{List } C \times (-)$$

$$\mathbf{G} = (\mathbb{N}, \leq)$$

$$Ge = \text{List}_{\leq e} C \times (-)$$

$$g_{e,X} = \lambda(s, x). (s, x)$$

Canonical gradings of objects

The canonical \mathcal{M} -grading of $T \in \mathbf{D}$ has:

- ▶ category of grades \mathcal{M}/T : a grade is a pair (S, s) , where $s : S \rightarrow T$ is in \mathcal{M}
- ▶ functor $T_{\mathcal{M}} : (S, s) \mapsto S : \mathcal{M}/T \rightarrow \mathbf{D}$
- ▶ natural transformation $g_{(S,s)} = s : G(S, s) \rightarrow T$

Universal property:

for every other grading (G, G, g) of T , there is an essentially unique functor $F : G \rightarrow \mathcal{M}/T$ with isomorphisms $Ge \cong \hat{T}(Fe)$ for all e , commuting with the natural transformations

$$\begin{array}{ccc} \mathbf{G} & & \\ \downarrow F & \searrow G & \\ \mathcal{M}/T & \xrightarrow{T_{\mathcal{M}}} & \mathbf{D} \end{array}$$

\cong

In other words: \mathcal{M}/T is pseudoterminal in a 2-category of \mathcal{M} -gradings of T

Some examples

When $\mathbf{D} = [\mathbf{Set}, \mathbf{Set}]$, \mathcal{M} = componentwise injective nat. transformations,
canonical grades are subfunctors $S \hookrightarrow T$

- ▶ For $T = \text{Id}$:

$$\mathcal{M}/T \simeq \{\perp \leq \top\} \quad T_{\mathcal{M}}\perp = \emptyset \quad T_{\mathcal{M}}\top = \text{Id}$$

- ▶ For $T = \text{List}C \times (-)$:


$$\mathcal{M}/T \simeq (\mathcal{P}(\text{List}C), \subseteq) \quad T_{\mathcal{M}}(\Sigma \subseteq \text{List}C) = \Sigma \times (-)$$

Some examples

If $T = V \Rightarrow (-)$ (for a set V), subfunctors $S \hookrightarrow T$ are equivalently upwards-closed sets

$$\Sigma \subseteq \text{Equiv}_V$$

$R \in \Sigma \Rightarrow R' \in \Sigma$ whenever $R \subseteq R'$



of equivalence relations of V , via

$$\Sigma = \{R \in \text{Equiv}_V \mid [-]_R \in S(V/R)\}$$

and these give a canonical grading

$$T_{\mathcal{M}} \Sigma X = \{f : V \rightarrow X \mid \exists R \in \Sigma. \forall v, v'. v R v' \Rightarrow f v = f v'\}$$

Graded monads [Smirnov '08, Melliès '12, Katsumata '14]

- ▶ A monad on \mathbf{C} is a monoid in $([\mathbf{C}, \mathbf{C}], \text{Id}, \circ)$
- ▶ A (\mathbf{G}, I, \odot) -graded monad on \mathbf{C} is a lax monoidal functor

$$\mathbb{T} = (T, \eta, \mu) : (\mathbf{G}, I, \odot) \rightarrow ([\mathbf{C}, \mathbf{C}], \text{Id}, \circ)$$

Explicitly:

$$T : \mathbf{G} \rightarrow [\mathbf{C}, \mathbf{C}] \quad \eta_X : X \rightarrow TXI \quad \mu_{e_1, e_2, X} : Te_1(Te_2X) \rightarrow T(e_1 \odot e_2)X$$

Example:

$$\text{WrL } e X = \text{List}_{\leq e} C \times X \quad \eta : X \rightarrow \text{WrL } 0 X \quad \mu : \text{WrL } e_1 (\text{WrL } e_2 X) \rightarrow \text{WrL } (e_1 + e_2) X$$

Grading monoids

Given

- ▶ a monoidal category (\mathbf{D}, I, \otimes)
- ▶ a class of \mathbf{D} -morphisms \mathcal{M}
- ▶ a monoid T in \mathbf{D}

an \mathcal{M} -grading (\mathbf{G}, G, g) of T is:

- ▶ a monoidal category \mathbf{G} of **grades**
- ▶ with a lax monoidal functor $G : \mathbf{G} \rightarrow \mathbf{D}$
- ▶ and a monoidal nat. trans.

$$g_e : Ge \rightarrow T$$

whose components are in \mathcal{M}

Example: writer monad

$$(\mathbf{D}, I, \otimes) = ([\mathbf{Set}, \mathbf{Set}], \text{Id}, \circ)$$

$\mathcal{M} =$ componentwise injective nat. trans.

$$T = \text{List } C \times (-) \text{ (a writer monad)}$$

$$\mathbf{G} = (\mathbb{N}, \leq), \text{ with addition}$$

$$Ge = \text{List}_{\leq e} C \times (-) \text{ (a graded writer monad)}$$

$$g_{e,X} = \lambda(s, x). (s, x)$$

Canonical gradings of monoids

If

- ▶ T is a monoid in (\mathbf{D}, I, \otimes)

then \mathbf{D}/T forms a monoidal category:

$$\begin{array}{ccc} J = (I, \eta) & & (S, s) \boxtimes (S', s') = (S \otimes S', (s \otimes s') \circ \mu_T) \\ I \xrightarrow{\eta} T & & S \otimes S' \xrightarrow{s \otimes s'} T \otimes T \xrightarrow{\mu} T \end{array}$$

and

$$(S, s) \mapsto S : \mathbf{D}/T \rightarrow \mathbf{D}$$

forms a monoidal functor

Canonical gradings of monoids

If

- ▶ T is a monoid in $(\mathbf{D}, \mathbf{l}, \otimes)$
- ▶ \mathcal{M} forms a factorization system $(\mathcal{E}, \mathcal{M})$
- ▶ \mathcal{E} is closed under \otimes on both sides

then \mathcal{M}/T forms a monoidal category:

$$\begin{array}{c} \mathcal{M}/T \\ \uparrow L \quad \downarrow \\ \mathbf{D}/T \end{array}$$

$$J = L(I, \eta)$$

$$\begin{array}{ccc} \mathbf{l} & \xrightarrow{\eta} & T \\ & \searrow & \nearrow \\ & J & \end{array}$$

$$(S, s) \boxtimes (S', s') = L(S \otimes S', (s \otimes s') \circ \mu_T)$$

$$\begin{array}{ccccc} S \otimes S' & \xrightarrow{s \otimes s'} & T \otimes T & \xrightarrow{\mu} & T \\ & \searrow & \downarrow & \nearrow & \\ & & S \boxtimes S' & & \end{array}$$

and

$$(S, s) \mapsto S : \mathcal{M}/T \rightarrow \mathbf{D}$$

forms a lax monoidal functor

Canonical gradings of monoids

The canonical \mathcal{M} -grading of a monoid T in (\mathbf{D}, I, \otimes) has:

- ▶ monoidal category of grades \mathcal{M}/T
- ▶ lax monoidal functor $T_{\mathcal{M}} : (S, s) \mapsto S : \mathcal{M}/T \rightarrow \mathbf{D}$
- ▶ monoidal natural transformation $g_{(S,s)} = s : G(S, s) \rightarrow T$

Universal property:

for every other grading (G, G, g) of the monoid T ,
there is an essentially unique lax monoidal
 $F : G \rightarrow \mathcal{M}/T$ with isomorphisms $Ge \cong \hat{T}(Fe)$ for all
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$$\begin{array}{ccc} \mathbf{G} & & \\ \downarrow F & \searrow G & \\ \mathcal{M}/T & \xrightarrow{T_{\mathcal{M}}} & \mathbf{D} \end{array}$$

In other words: \mathcal{M}/T is pseudoterminal in a 2-category of \mathcal{M} -gradings of T

Example: writer

Take

- ▶ $\mathbf{D} = [\mathbf{Set}, \mathbf{Set}]$, with endofunctor composition
- ▶ $(\mathcal{E}, \mathcal{M}) = (\text{componentwise surjective}, \text{componentwise injective})$
- ▶ T is a writer monad

$$T = \text{List}C \times (-)$$

Then:

- ▶ Subfunctors $S \hookrightarrow \text{Wr}$ are equivalently subsets $\Sigma \subseteq \text{List}C$ via

$$\Sigma = \{s \in \text{List}C \mid (s, \star) \in S1\} \quad SX = \{(s, x) \in \text{List}C \times X \mid s \in \Sigma\}$$

So the canonical grading is $\mathcal{P}(\text{List}C)$ with

$$\begin{aligned} T_{\mathcal{M}} : (\mathcal{P}(\text{List}C), \subseteq) &\rightarrow [\mathbf{Set}, \mathbf{Set}] & T_{\mathcal{M}}\Sigma &= \Sigma \times (-) \\ J &= \{[]\} & \Sigma_1 \boxplus \Sigma_2 &= \{s_1 \# s_2 \mid s_1 \in \Sigma_1, s_2 \in \Sigma_2\} \end{aligned}$$

Grading by sets of shapes

For $\mathbf{D} = [\mathbf{Set}, \mathbf{Set}]$, there is also a factorization system

\mathcal{E} = natural transformations α such that α_1 is surjective

\mathcal{M} = cartesian natural transformations α such that α_1 is injective

satisfying

$$\mathcal{M}/T \simeq (\mathcal{P}(T1), \subseteq)$$

Summary

Given a suitable class \mathcal{M} of morphisms, every monoid T has a **canonical \mathcal{M} -grading**
 $T_{\mathcal{M}} : \mathcal{M}/T \rightarrow \mathbf{D}$

$$\begin{array}{ccc} \mathbf{G} & & \\ \downarrow F & \searrow G & \\ \mathcal{M}/T & \xrightarrow{T_{\mathcal{M}}} & \mathbf{D} \end{array}$$

In particular, we can canonically grade monads (and algebraic operations for them)