

Q Suppose that X_n satisfies an LDP in some regular space \mathcal{X} , good rate func. I
 Y_n ————— \mathcal{Y} ————— J
 and that X_n and Y_n are independent.

Show that (X_n, Y_n) satisfies an LDP in $\mathcal{X} \times \mathcal{Y}$, good rate func.

$$K(x, y) = I(x) + J(y).$$

A Note first the topology on $\mathcal{X} \times \mathcal{Y}$.

If σ and τ are bases for \mathcal{X} and \mathcal{Y} , then $\{O \times P : O \in \sigma, P \in \tau\}$ is a basis for $\mathcal{X} \times \mathcal{Y}$.

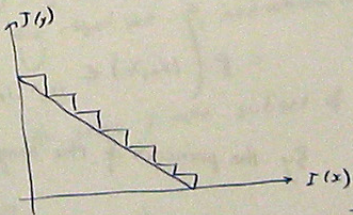
Thus open sets are of the form $\bigcup_n O_n \times P_n$ O_n open in \mathcal{X} , P_n open in \mathcal{Y}
 closed ————— $\bigcap_n ((C_n \times \mathcal{Y}) \cup (\mathcal{X} \times D_n))$ C_n closed in \mathcal{X} , D_n closed in \mathcal{Y} .

Goodness of K .

(clearly $K(x, y) \geq 0$.)

A typical level set is

$$\begin{aligned} \{(x, y) : K(x, y) \leq \alpha\} &= \{(x, y) : I(x) + J(y) \leq \alpha\} \\ &= \bigcap_{n \in \mathbb{N}} \bigcup_{m \in \mathbb{N}} \left\{ (x, y) : I(x) \leq \frac{m}{n} \alpha, J(y) \leq \frac{n+1-m}{n} \alpha \right\} \end{aligned}$$



$$\begin{aligned} &= \bigcap_{n \in \mathbb{N}} \bigcup_{m \in \mathbb{N}} \left\{ x : I(x) \leq \frac{m}{n} \alpha \right\} \times \left\{ y : J(y) \leq \frac{n+1-m}{n} \alpha \right\} \\ &= \bigcap_{n \in \mathbb{N}} \bigcup_{m \in \mathbb{N}} \text{closed} \times \text{closed} = \bigcap_{n \in \mathbb{N}} \text{closed} = \text{closed}. \end{aligned}$$

Furthermore, this level set is a subset of the compact set

$$\left\{ x : I(x) \leq \alpha \right\} \times \left\{ y : J(y) \leq \alpha \right\}.$$

So this level set is compact.

Lower Bound

Let B be open in $X \times Y$. Then $B = \bigcup_n O_n \times P_n$.

Pick $(x, y) \in B$. Then $(x, y) \in O \times P$ for some O open in X
 P open in Y
 $O \times P \subset B$.

$$\mathbb{P}((x_n, y_n) \in B) \geq \mathbb{P}((x_n, y_n) \in O \times P) = \mathbb{P}(x_n \in O) \mathbb{P}(y_n \in P).$$

$$\Rightarrow \liminf \frac{1}{n} \log \mathbb{P}((x_n, y_n) \in B)$$

$$\geq \liminf \frac{1}{n} \left[\log \mathbb{P}(x_n \in O) + \log \mathbb{P}(y_n \in P) \right]$$

$$\geq \liminf \frac{1}{n} \log \mathbb{P}(x_n \in O) + \liminf \frac{1}{n} \log \mathbb{P}(y_n \in P)$$

$$\geq - \inf_{x' \in O} I(x') - \inf_{y' \in P} J(y')$$

$$\geq - (I(x) + J(y)) = -K(x, y).$$

Upper Bd for Cylinder sets

Consider first a simple closed set of the form $B = \bigcap_{n \in \mathbb{N}} (C_n \times Y) \cup (X \times D_n)$.

For such a set,

$$\mathbb{P}((x_n, y_n) \in B) = \mathbb{P}((x_n, y_n) \in \bigcup_{(i_1, \dots, i_N) \in \{0,1\}^N} \bigcap_{n \in \mathbb{N}} \begin{cases} i_n=0: C_n \times Y \\ i_n=1: X \times D_n \end{cases})$$

$$= \mathbb{P}((x_n, y_n) \in \bigcup_{(i_1, \dots, i_N)} \left(\bigcap_{n: i_n=0} C_n \right) \times \left(\bigcap_{n: i_n=1} D_n \right))$$

By the principle of the largest term,

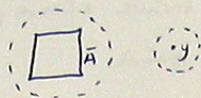
$$\limsup \frac{1}{n} \log \mathbb{P}((x_n, y_n) \in B) \leq \max_{i_1, \dots, i_N} \limsup \frac{1}{n} \log \mathbb{P}((x_n, y_n) \in \left(\bigcap_{n: i_n=0} C_n \right) \times \left(\bigcap_{n: i_n=1} D_n \right))$$

$$\leq - \min_{i_1, \dots, i_N} \left(\inf_{x \in \bigcap_{n: i_n=0} C_n} I(x) + \inf_{y \in \bigcap_{n: i_n=1} D_n} J(y) \right).$$

$$= - \inf_{(x, y) \in B} (I(x) + J(y)).$$

Upper bound for general closed sets

Recall: a space is regular if for every closed set \bar{A} and every point y there exist disjoint open neighbourhoods $\bar{A} \subset B_{\bar{A}}$ and $y \in B_y$.



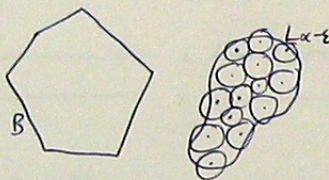
It can be shown that if X and Y are regular then $X \times Y$ is regular also. [simple exercise].

Let $B \subset X \times Y$ be closed.

Let $\alpha = \inf_{z \in B} K(z)$. If $\alpha = 0$, upper bd is trivial.

Otherwise, for any $\epsilon > 0$, consider the level set $L_{\alpha-\epsilon}$.

Clearly, $B \cap L_{\alpha-\epsilon} = \emptyset$. By goodness, $L_{\alpha-\epsilon}$ is compact.



For every $z \in L_{\alpha-\epsilon}$, by regularity, we can pick an open ball B_z which does not intersect the closed set B , and which is of the form $O \times P$.

By compactness of $L_{\alpha-\epsilon}$,

$$L_{\alpha-\epsilon} \subseteq \bigcup_{n \in \mathbb{N}} B_{z_n} \quad \text{for some finite subset of the } z \in L_{\alpha-\epsilon}.$$

Now, by construction of the B_z ,

$$B \subseteq \left(\bigcup_{n \in \mathbb{N}} B_{z_n} \right)^c = \bigcap_{n \in \mathbb{N}} C_{z_n} \quad \text{where } C_{z_n} = B_{z_n}^c \text{ is closed.}$$

$$\begin{aligned} \text{So } \limsup_n \frac{1}{n} \log P((x_n, y_n) \in B) &\leq \limsup_n \frac{1}{n} \log P((x_n, y_n) \in \bigcap_{n \in \mathbb{N}} C_{z_n}) \\ &\leq \inf_{z \in \bigcap_{n \in \mathbb{N}} C_{z_n}} K(z) \quad \text{by upper bd for cylinder sets} \end{aligned}$$

But for any such z , by choice of the B_z , $K(z) > \alpha - \epsilon$.

Hence $\limsup_n \frac{1}{n} \log P((x_n, y_n) \in B) \leq -(\alpha - \epsilon)$. Since ϵ arbitrary, the LD upper bound holds.

1 (i). • $\limsup \frac{1}{n} \log(a_n + b_n) = \limsup \frac{1}{n} \log a_n \vee \limsup \frac{1}{n} \log b_n$? wlog $a \geq b$.

$\forall \epsilon > 0 \exists n_0$ s.t. $n \geq n_0 \Rightarrow a_n = e^{n(A+\epsilon)}$ and $b_n = e^{n(B+\epsilon)}$.

$$\begin{aligned} \Rightarrow \limsup \frac{1}{n} \log(a_n + b_n) &\leq \limsup \frac{1}{n} \log [e^{n(A+\epsilon)} + e^{n(B+\epsilon)}] \\ &= \limsup \frac{1}{n} \log e^{n(A+\epsilon)} [1 + e^{-n(A-B)}] \\ &= A + \epsilon. \end{aligned}$$

But $\epsilon > 0$ arbitrary, so $\limsup \frac{1}{n} \log(a_n + b_n) = A$.

• $\liminf \frac{1}{n} \log(a_n + b_n) = \liminf \frac{1}{n} \log a_n \vee \liminf \frac{1}{n} \log b_n$? wlog $a \geq b$

$\rightarrow \liminf \frac{1}{n} \log a_n$ since $b_n \geq 0$

$= A$.

(ii). • $\lim \frac{1}{n} \log P(A_n \cup B_n) = \limsup \frac{1}{n} \log(a_n + b_n)$ $a_n = \log P(A_n)$ $b_n = \log P(B_n)$

$\left. \begin{aligned} &\leq (-a) \vee (-b) = -(a \wedge b). \\ &\liminf \frac{1}{n} \log a_n = -a \\ &\liminf \frac{1}{n} \log b_n = -b \end{aligned} \right\} \text{agree.}$

so $\lim = -(a \wedge b)$

• suppose $a > b$.

$$\begin{aligned} \frac{1}{n} \log P(A_n | A_n \cup B_n) &= \frac{1}{n} \log P(A_n) - \frac{1}{n} \log P(A_n \cup B_n) \rightarrow -a - (-(a \wedge b)) = -(a-b) < 0. \\ \& P(A_n | A_n \cup B_n) \rightarrow 0. \end{aligned}$$

suppose $a < b$.

By above, $P(B_n | A_n \cup B_n) \rightarrow 0$.

But $P(A_n | A_n \cup B_n) = 1 - P(B_n | A_n \cup B_n) \rightarrow 1$.



(i) $K \subset X$ compact, $I = \inf_{x \in K} f(x) < \infty$.

Pick $x_n \in K$, $I \leq f(x_n) < I + \frac{1}{n}$.

By compactness, (x_n) has convt subsequence $(x_{k(n)})$,

$$x_{k(n)} \in K, \quad x_{k(n)} \rightarrow x^* \in K, \quad f(x_{k(n)}) \rightarrow I.$$

compactness

What is $f(x^*)$?

$L_{I+\epsilon} = \{x: f(x) \leq I+\epsilon\}$ is closed, and $x_k \in L_{I+\epsilon}$ for suff. large k ,

so the limit $x^* \in L_{I+\epsilon}$. i.e. $f(x^*) \leq I+\epsilon$

This is true for all $\epsilon > 0$; so $f(x^*) = I$.

Since $x^* \in K$, $f(x^*) \geq I$.

Thus $f(x^*) = I$.

(ii) No.

But any convex function $f: \mathbb{R} \rightarrow \mathbb{R}$ is continuous.

PROOF. Suppose not. Let $x_n \rightarrow x$, $f(x_n) \not\rightarrow f(x)$.

So either $\limsup f(x_n) > f(x)$ or $\liminf f(x_n) < f(x)$.

Suppose $x_n \uparrow x$. (The case $x_n \downarrow x$ is similar, and these two cases are sufficient.)

(1) $x_n = (1-\lambda_n)x + \lambda_n x$ for some $\lambda_n \rightarrow 1$.

By convexity, $f[(1-\lambda_n)x + \lambda_n x] \leq (1-\lambda_n)f(x) + \lambda_n f(x) \rightarrow f(x)$.

Thus $\limsup f(x_n) = f(x)$. ✗

(2) By convexity, $f(y) \geq f(x) + m(y-x)$ [a supporting line, and $m < \infty$ since $f < \infty$].

$$f(x_n) \geq f(x) + m(x_n - x)$$

$\liminf f(x_n) \geq f(x)$. ✗

(iii) Consider a level set $\{x: g^*(x) \leq \alpha\} = \left\{x: \sup_{\theta \in \mathbb{R}} \theta x - g(\theta) \leq \alpha\right\}$
 $= \left\{x: \theta x \leq \alpha + g(\theta) \text{ for all } \theta \in \mathbb{R}\right\}$.

If $0 > \alpha + g(0)$, this set is empty, hence closed.

If $0 \leq \alpha + g(0)$:
 $= \left\{x: x \leq \frac{\alpha + g(\theta)}{\theta} \forall \theta > 0, x \geq \frac{\alpha + g(\theta)}{\theta} \forall \theta < 0\right\}$

$$= \left\{x: \sup_{\theta < 0} \frac{\alpha + g(\theta)}{\theta} \leq x \leq \inf_{\theta > 0} \frac{\alpha + g(\theta)}{\theta}\right\}$$

which is a closed interval, hence closed.

2 (iv) • Suppose f is lower-semicontinuous, i.e. level set closed.

Let $I = \liminf f(x)$, so that $f(x_n) \geq I + \epsilon$ for ∞ many n , any $\epsilon > 0$.

Since $x_n \rightarrow x$, $x \in L_{I+\epsilon}$ level set $L_{I+\epsilon}$ is closed, and these $x_{k(n)}$ all lie in $L_{I+\epsilon}$, so $\liminf f(x) \geq I + \epsilon$.

So $f(x) \geq I + \epsilon$.

This is true for all $\epsilon > 0$; thus $f(x) \geq I$.

• Suppose $x_n \rightarrow x \Rightarrow \liminf f(x_n) \geq f(x)$.

Is $L_I = \{x : f(x) \leq I\}$ closed? Pick $x_n \in L_I$, $x_n \rightarrow x$. Is $x \in L_I$?

Certainly $f(x) \leq \liminf f(x_n) \leq I$. So $x \in L_I$.

metric: closed \Leftrightarrow all limit points.

3 • $\limsup \frac{1}{n} \log P(Z_n \in B)$, B closed

$$= \limsup \frac{1}{n} \log [P(Z_n \in B, B_n = 0) + P(Z_n \in B, B_n = 1)]$$

$$= \limsup \frac{1}{n} \log [P(X_n \in B, B_n = 0) + P(Y_n \in B, B_n = 1)]$$

$$= \limsup \frac{1}{n} \log [(1-p)P(X_n \in B) + pP(Y_n \in B)]$$

$$\leq \left(\limsup \frac{1}{n} \log (1-p)P(X_n \in B) \right) \vee \left(\limsup \frac{1}{n} \log pP(Y_n \in B) \right)$$

$$= \limsup \frac{1}{n} \log P(X_n \in B) \quad \vee \quad \limsup \frac{1}{n} \log P(Y_n \in B) \quad \text{if } 0 < p < 1.$$

$$= \left(- \int_{x \in B} I(x) \right) \vee \left(- \int_{y \in B} J(y) \right) = - \int_{x \in B} I(x) \wedge J(y).$$

• $\liminf \frac{1}{n} \log P(Z_n \in B)$, B open

$$\geq \liminf \frac{1}{n} \log P(X_n \in B, B_n = 0) \geq - \int_{x \in B} I(x).$$

Also $\geq - \int_{x \in B} J(x).$

Hence $\geq - \int_{x \in B} I(x) \wedge J(x).$

Or: if the spaces are separable, rate func are good:

$$B_n \text{ satisfies LDP } \frac{1}{n} \log P(B_n \in B) = 0 \\ B \subset \{0,1\}.$$

So (x_n, y_n, B_n) satisfies LDP, grf $K(x, y, b) = I(x) + J(y) + 0$
 $\mathbb{P}^x \mathbb{P}^y \mathbb{P}^b \in \{0,1\}$.

So z_n acts function of (x_n, y_n, B_n) , satisfies LDP.

$$\text{grt } L(z) = \inf_{x, y, b} \{ I(x) + J(y) \} \\ \begin{matrix} 1=0 & x=z \\ \text{or } b=1 & y=z \end{matrix}$$

$$= \min \begin{cases} b=0: & \inf_y I(z) + J(y) \\ b=1: & \inf_x I(x) + J(z) \end{cases}$$

$$= \min \begin{cases} b=0: & I(z) \\ b=1: & J(z) \end{cases} \quad \text{since } \inf_{x \in X} I(x) = 0, \text{ by LD upper tail for } X.$$

$$= I(z) \wedge J(z).$$

• $\limsup \frac{1}{n} \log P(X_n \in B)$ for B closed in Σ

$\Rightarrow B^c$ open in $\Sigma \Rightarrow B^c \cap \Sigma$ open in X .

$\Rightarrow B^c \cap \Sigma = A^c$, for some A closed in X

$= A^c \cap \Sigma$ since $B^c \cap \Sigma \subset \Sigma$, ~~and closed~~

$\Rightarrow B \cup \Sigma^c = A \cup \Sigma^c = (A \cap \Sigma) \cup \Sigma^c$

$\Rightarrow B = A \cap \Sigma$, A closed in X .

$\leq \limsup \frac{1}{n} \log P(X_n \in A \cap \Sigma)$

$= -\inf_{x \in A \cap \Sigma} I(x) = -\inf_{x \in B} I(x)$.

• $\liminf \frac{1}{n} \log P(X_n \in B)$ for B open in Σ

$\Rightarrow B = A \cap \Sigma$, A open in X

$= \liminf \frac{1}{n} \log P(X_n \in A \cap \Sigma)$

$= \liminf \frac{1}{n} \log P(X_n \in A)$ since $P(\Sigma) = 1$

$\geq -\inf_{x \in A} I(x) \geq -\inf_{x \in A \cap \Sigma} I(x) = -\inf_{x \in B} I(x)$.

$\inf_{x \in A} I(x) \leq \inf_{x \in A \cap \Sigma} I(x)$

\uparrow
larger smaller

5. (i). $\limsup \frac{1}{n^{2(1-G)}} \log P(X_n \in C) ?$ If $\int_{x \in C} I'(x) = 0$, trivial. So assume $\int_{x \in C} I'(x) = \infty$.

$$= \limsup n^{2(G-H)} \cdot \frac{1}{n^{2(1-H)}} \log P(X_n \in C)$$

$$\limsup \frac{1}{n^{2(1-H)}} \log P(X_n \in C) \leq - \int_{x \in C} I(x)$$

What is this int? since C is good, the int is attained.

If $\int_{x \in C} I(x) = 0$, $\int_{x \in C} I'(x) = 0$ by non-negativity (iii) — discard this case.

Else $\int_{x \in C} I(x) > 0 \Rightarrow \int_{x \in C} I'(x) = \infty$. Let $\alpha = \int_{x \in C} I(x) > 0$.

$$\leq -\alpha$$

So $\frac{1}{n^{2(1-H)}} \log P(X_n \in C) \leq -\alpha + \varepsilon$ for n suff. large. (where $\varepsilon < \alpha$).

$$\text{So } n^{2(G-H)} \cdot \frac{1}{n^{2(1-H)}} \log P(X_n \in C) \leq -n^{2(G-H)} (\alpha - \varepsilon) \rightarrow -\infty$$

$$\text{Thus } \limsup \frac{1}{n^{2(1-G)}} \log P(X_n \in C) \leq - \int_{x \in C} I'(x)$$

(ii). • $\limsup \frac{1}{n} \log P(X_n \notin D) \leq - \int_{x \in D^c} I(x) < 0$ (by goodness, closure of $D \rightarrow$ int attained).

$$\text{So } P(X_n \notin D) \rightarrow 0$$

• $\liminf \frac{1}{n^{2(1-G)}} \log P(X_n \in E) ?$

If E contains μ , $\int_{x \in E} I'(x) = 0$;

and $P(X_n \in E) \rightarrow 1$ so $\log P(X_n \in E) \rightarrow 0$ so $\liminf \frac{1}{n^{2(1-G)}} \log P(X_n \in E) = 0$.

If E does not contain μ , $\int_{x \in E} I'(x) = \infty$, so bd is trivial.

(iii) so we have LDP at speed $n^{2(1-G)}$, rate function I' .

(clearly a prate function (non-negative, compact level sets).)

wlog $H < G$.

(iv). By LDP for product spaces, (X_n, Y_n) satisfies LDP at speed $n^{2(1-G)}$, gff $J'(x) + J'(y)$.
 so $X_n + Y_n$ satisfies LDP at same speed, gff $K(z) = \int_{x,y: x+y=z} I'(x) + J'(y) = J(\frac{z}{2}, \mu)$ non-trivial no!
 which is non-trivial. \square

[7] (i).

$B \sim \text{Exp}(\lambda)$.

so $\frac{B}{L}$ satisfies LDP in \mathbb{R}_+ , g.f $J(x) = \lambda x$.

(ii) By Cramér:

$\frac{A_1 + \dots + A_n}{L}$ satisfies LDP in \mathbb{R} , g.f $J(x) = \sup_{\theta} \theta x - \log \mathbb{E} e^{\theta A_1}$

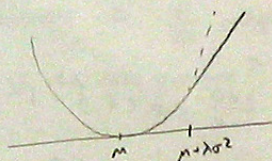
$$= \sup_{\theta} \theta x - (\theta \mu + \frac{1}{2} \theta^2 \sigma^2)$$

$$= \frac{1}{2\sigma^2} (x - \mu)^2$$

(iii) By contraction; Product Spaces:

$\frac{B}{L} + \frac{A_1 + \dots + A_n}{L}$ satisfies LDP, g.f

$$k(x) = \inf_{\substack{b, \theta, a: \\ b+a=x}} \lambda b + \frac{1}{2\sigma^2} (a-\mu)^2 = \begin{cases} x \geq \mu + \lambda \sigma^2: \lambda(x-\mu) - \frac{\lambda^2 \sigma^2}{2} \\ \quad \text{at } b = x - \mu - \lambda \sigma^2, a = \mu + \lambda \sigma^2 \\ x < \mu + \lambda \sigma^2: \frac{(x-\mu)^2}{2\sigma^2} \\ \quad \text{at } b=0, a=x. \end{cases}$$



Exp only kicks in at a certain pt.

NOTE: G-E won't work.

$$\frac{1}{L} \log \mathbb{E} e^{\theta(B + A_1 + \dots + A_n)} = \theta \mu + \frac{1}{2} \theta^2 \sigma^2 + \frac{1}{L} \cdot \frac{\lambda}{\lambda - \theta}$$

$$\rightarrow \begin{cases} \theta \mu + \frac{1}{2} \theta^2 \sigma^2 & \text{if } \theta < \lambda \\ \infty & \theta \geq \lambda \end{cases} \quad \text{Not steep!}$$

(iv) By Gärtner-Ellis.

$$\frac{1}{L} \log \mathbb{E} e^{\theta(C + A_1 + \dots + A_n)} \rightarrow \theta \mu + \frac{1}{2} \theta^2 \sigma^2$$

So $\frac{C}{L} + \frac{A_1 + \dots + A_n}{L}$ satisfies same LDP as $\frac{A_1 + \dots + A_n}{L}$.

OR by product & contraction. Try $C \sim N(0,1)$ first — result for $C \sim N(\mu, \sigma^2)$ follows by contraction.

$$\frac{C}{L} \text{ satisfies LDP at speed } L^2 \quad \mathbb{P}\left(\frac{C}{L} \approx x\right) \approx \mathbb{P}(N(0,1) \approx Lx) \approx \frac{1}{\sqrt{2\pi}} e^{-\frac{1}{2} L^2 x^2}$$

or use Cramér: $\frac{C}{L} \approx \frac{1}{L} (N(0,1) + \dots + N(0,1))$

Via Hoeffding's, satisfies LDP at speed L ,

$$\text{g.f } J'(x) = \begin{cases} 0 & x=0 \\ \infty & \text{else.} \end{cases}$$

8 (i) ~~we~~ We need only model the header queue.

This has Poisson arrivals of rate λ ,
and service times are Exp with mean $1/c$.

So queue size evolves like $M_\lambda / M_c / 1$.

Stationary queue size dist is

$$P(Q \geq q) = \left(\frac{\lambda}{c}\right)^q.$$

(ii) Let A = amount of work arriving in a typical timeslot.

→ made up from $N \sim \text{Poisson}(\lambda\delta)$ packets, of size $x_1, \dots, x_N \sim \text{Exp}(1)$

$$E e^{\theta A} = E e^{\theta(x_1 + \dots + x_N)}$$

$$= E \left[E e^{\theta(x_1 + \dots + x_N)} | N \right]$$

$$= E \left[(E e^{\theta x_1})^N \right] = E \left[\left(\frac{1}{1-\theta}\right)^N \right] = e^{\lambda\delta \left(\frac{1}{1-\theta} - 1\right)} = e^{\lambda\delta\theta / (1-\theta)}$$

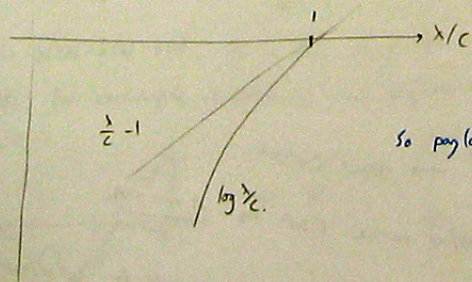
By theorem,

$$\frac{1}{r} \log P(R_\delta > r) \approx -\sup \left\{ \theta > 0 : \frac{\lambda\delta\theta}{1-\theta} < \theta < \delta \right\} = 1 - \frac{\lambda}{c}$$

NOTE: doesn't depend on δ ! Good!

(iii) Header-queue: $P(Q \geq q) = \left(\frac{\lambda}{c}\right)^q = e^{q \log \lambda/c}$

Payload-queue: $P(R_\delta > r) \approx e^{r \cdot (1 - \lambda/c)}$



So payload-queue more likely to overflow.

This is reasonable — there are more ways for the payload queue to overflow
e.g. average # packets which are all r large.

9

Note first:

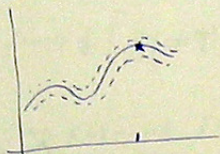
* a continuous function on a compact set attains its supremum.

[let $f(x_k) \uparrow$ supremum.

Then (x_k) has convt subsequence, $(x_j) \rightarrow x$

Clearly $f(x_j) \rightarrow$ supremum $= f(x)$.

* The function $x \mapsto \sup_{\Delta \leq t \leq 1} x(t)$ is cts w.r.t. sup-norm on $\mathcal{C}[0,1]$.

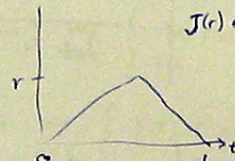


So $b \mapsto R(b) = \sup_{t \in [0,1]} x(t) - \inf_{t \in [0,1]} x(t)$, where $x(t) = b(t) - b(1)$,

is continuous. So $\frac{R}{\sqrt{N}}$ satisfies LDP.

What is the rate function? $J(r) = \inf_{b: R(b)=r} I(b)$, $I(b) = \begin{cases} \frac{1}{2} \int_0^1 \dot{b}_t^2 dt & \text{if } b \text{ is abs. cts} \\ \infty & \text{else.} \end{cases}$

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consider $\hat{b}(t) =$  $J(r) = I(\hat{b}) = \frac{1}{2} \left[\frac{1}{2} \left(\frac{r}{1/2}\right)^2 + \frac{1}{2} \left(r/\frac{1}{2}\right)^2 \right] = 2r^2$.

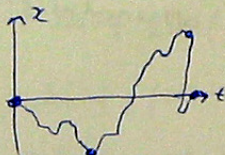
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Suppose $R(b) = r$. We will argue $I(b) \geq 2r^2 \dots$

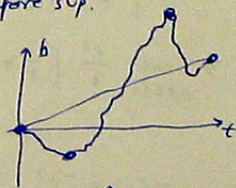
Either $I(b) = \infty$ (in which case we are done) or $I(b) < \infty$, in which case...

Look at $x(t) = b(t) - b(1)$. It is cts, so it attains its inf and sup.

Suppose for example it attains inf before sup.



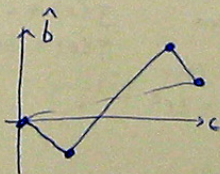
so b is



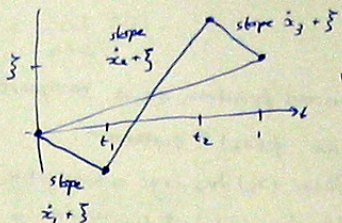
Now consider the straightened path $\hat{b} =$

By convexity of I ,

$$I(b) \geq I(\hat{b}).$$



The rate function of \hat{b} is



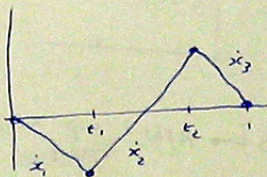
where $t_1 \dot{x}_1 + (t_2 - t_1) \dot{x}_2 + (1 - t_2) \dot{x}_3 = 0$

$$I(\hat{b}) = t_1 (\dot{x}_1 + \zeta)^2 + (t_2 - t_1) (\dot{x}_2 + \zeta)^2 + (1 - t_2) (\dot{x}_3 + \zeta)^2$$

$$= \zeta^2 + 2\zeta (t_1 \dot{x}_1 + (t_2 - t_1) \dot{x}_2 + (1 - t_2) \dot{x}_3) + \text{const.}$$

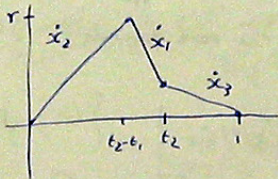
\Rightarrow const; min achieved at $\zeta = 0$.

Thus $I(\hat{b}) \geq I(\hat{\hat{b}})$, $\hat{\hat{b}}$ is



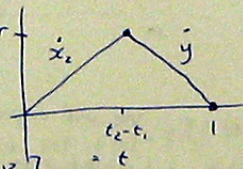
By just shuffling the pieces,

$I(\hat{\hat{b}}) \geq I(\hat{\hat{\hat{b}}})$, $\hat{\hat{\hat{b}}}$ is



By convexity,

$I(\hat{\hat{\hat{b}}}) \geq I(\hat{\hat{\hat{\hat{b}}}})$, $\hat{\hat{\hat{\hat{b}}}}$ is



where $t = t_2 - t_1$
 $\dot{x}_2 = r/t$
 $\dot{y} = -r/(1-t)$

$$I(\hat{\hat{\hat{\hat{b}}}}) = \frac{1}{2} \left[t \left(\frac{r}{t}\right)^2 + (1-t) \left(-\frac{r}{1-t}\right)^2 \right]$$

$$= \frac{r^2}{2} \left[\frac{1}{t} + \frac{1}{1-t} \right]$$

$$\Rightarrow \inf_{0 \leq t \leq 1} \frac{r^2}{2} \left[\frac{1}{t} + \frac{1}{1-t} \right] = 2r^2, \text{ attained at } t = \frac{1}{2}$$

We've shown: $J(r) \leq 2r^2$
 $J(r) \geq 2r^2$

Thus $J(r) = 2r^2$.

$$\square \quad (i) \quad \limsup \frac{1}{n} \log P(X_n \in C_i) \leq - \inf_{x \in C_i} I(x).$$

$$\begin{aligned} & \limsup \frac{1}{n} \log E e^{n f(x_n)} \\ & \leq \limsup \frac{1}{n} \log \left[\sum_i e^{n \sup_{x \in C_i} f(x)} \cdot P(X_n \in C_i) \right] \\ & \leq \max_i \limsup \frac{1}{n} \log e^{n \sup_{x \in C_i} f(x)} P(X_n \in C_i) \quad \text{by PLT} \\ & = \max_i \limsup \frac{1}{n} \left[\sup_{x \in C_i} f(x) + \frac{1}{n} \log P(X_n \in C_i) \right] \\ & \leq \max_i \sup_{x \in C_i} f(x) - \inf_{x \in C_i} I(x). \end{aligned}$$

$$(ii) \quad \limsup \frac{1}{n} \log E e^{n f(x_n)} \leq \max_i \sup_{x \in C_i} f(x) - \inf_{x \in C_i} I(x)$$

~~$\leq \max_i \sup_{x \in C_i} f(x) - \inf_{x \in C_i} I(x)$~~ where \hat{x} attains inf on C_i .

$$\leq \max_i \sup_{x \in C_i} \left\{ \sup_{z \in C_i} f(z) - I(x) \right\}$$

so $f(\hat{x}) \approx \sup_{x \in C_i} f(x) - \varepsilon$.

$$\leq \max_i \sup_{x \in C_i} \left\{ f(x) + \varepsilon - I(x) \right\}$$

$$= \sup_{x \in \mathcal{X}} f(x) - I(x) + \varepsilon.$$

But $\varepsilon > 0$ arbitrary, so $\limsup = \sup_{x \in \mathcal{X}} f(x) - I(x)$.

$$(iii) \quad \liminf \frac{1}{n} \log E e^{n f(x_n)} \geq \liminf \frac{1}{n} \log E (e^{n(f(\hat{x})-\varepsilon)} \mathbb{1}_{\{X_n \in B\}})$$

$$\geq \liminf \frac{1}{n} \log (e^{n(f(\hat{x})-\varepsilon)} P(X_n \in B))$$

$$= \liminf \left[f(\hat{x}) - \varepsilon + \frac{1}{n} \log P(X_n \in B) \right]$$

$$\geq f(\hat{x}) - \varepsilon - \inf_{x \in B} I(x)$$

$\geq f(\hat{x}) - \varepsilon - I(\hat{x})$. Since ε arb, $\geq f(\hat{x}) - I(\hat{x})$.

Since \hat{x} arb, $\geq - \inf_{x \in \mathcal{X}} f(x) - I(x)$.

$$(iv) \quad \liminf = \limsup.$$

Q15

INTUITION.

$$\mathbb{P}\left(\frac{X^{(N)}}{N} \approx x\right) \approx \mathbb{P}\left(\frac{X^{(N)}}{\alpha N} \approx \frac{x}{\alpha}\right) \approx e^{-\alpha N \cdot I(x/\alpha)}$$

so $\frac{X^{(N)}}{N}$ satisfies LDP, rate func. $J(x) = \alpha I(x/\alpha)$.

PROOF.

LED Upper bound for semi-infinite intervals $[x, \infty)$.

$$\begin{aligned} \limsup_{N \rightarrow \infty} \frac{1}{N} \log \mathbb{P}\left(\frac{X^{(N)}}{N} \geq x\right) &= \limsup_{N \rightarrow \infty} \frac{L(N)}{N} \cdot \frac{1}{L(N)} \log \mathbb{P}\left(\frac{X^{(N)}}{L(N)} \geq \frac{N}{L(N)} x\right) \\ &\leq \limsup_{N \rightarrow \infty} \frac{L(N)}{N} \cdot \frac{1}{L(N)} \log \mathbb{P}\left(\frac{X^{(N)}}{L(N)} \geq \frac{x}{\alpha} - \varepsilon\right) \quad \text{for any } \varepsilon > 0 \\ &\leq -\alpha \cdot \inf_{y \geq \frac{x}{\alpha} - \varepsilon} I(y). \end{aligned}$$

This is true for any $\varepsilon > 0$, so

$$\limsup_{N \rightarrow \infty} \frac{1}{N} \log \mathbb{P}\left(\frac{X^{(N)}}{N} \geq x\right) \leq -\alpha \lim_{\varepsilon > 0} \inf_{y \geq \frac{x}{\alpha} - \varepsilon} I(y) \leq -\alpha \inf_{y \geq \frac{x}{\alpha}} I(y).$$

For the last step, see Q2 (iv).

Specifically, suppose f is lower-semicontinuous and

$$\lim_{\varepsilon > 0} \inf_{y \geq z - \varepsilon} f(y) < \inf_{y \geq z} f(y).$$

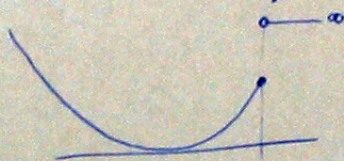
Then, for ~~arbitrarily small~~ $\varepsilon > 0$, $\exists y_\varepsilon$ with $z - \varepsilon \leq y_\varepsilon \leq z$

and $f(y_\varepsilon) < \inf_{y \geq z} f(y) - \delta$, for some $\delta > 0$.

By Q2 (iv), $\liminf_{\varepsilon > 0} f(y_\varepsilon) > f(\lim_{\varepsilon > 0} y_\varepsilon) = f(z)$. \times

Thus $\lim_{\varepsilon > 0} \inf_{y \geq z - \varepsilon} f(y) \geq \inf_{y \geq z} f(y)$.

To make this less abstract, convince yourself that I must look something like



2. LD Upper bound for semi-infinite intervals $(-\infty, x]$.
Proof similar to above.

3. LD Upper bound for general closed sets.
Proof exactly as for Gramér's theorem.

4. LD lower bound.

Pick any $\hat{x} \in B$, for some given open set B . Let $(\hat{x}-\delta, \hat{x}+\delta) \subset B$.

$$\liminf \frac{1}{N} \log \mathbb{P} \left(\frac{X^{L \times N}}{N} \in B \right)$$

$$\geq \liminf \frac{L \times N}{N} \cdot \frac{1}{L \times N} \log \mathbb{P} \left(\frac{X^{L \times N}}{L \times N} \in \frac{N}{L \times N} (\hat{x}-\delta, \hat{x}+\delta) \right)$$

$$\geq \liminf \frac{L \times N}{N} \cdot \frac{1}{L \times N} \log \mathbb{P} \left(\frac{X^{L \times N}}{L \times N} \in \left(\frac{\hat{x}}{\alpha} - \delta', \frac{\hat{x}}{\alpha} + \delta' \right) \right) \text{ for suff. small } \delta'$$

$$\geq -\alpha \cdot \inf_{\frac{\hat{x}}{\alpha} - \delta' \leq y \leq \frac{\hat{x}}{\alpha} + \delta'} I(y) \geq -\alpha I\left(\frac{\hat{x}}{\alpha}\right).$$

Take the inf to give the result.
 $\hat{x} \in B$

NOTE. The upper bound would be much simpler if we had assumed that I was everywhere finite (and hence, by convexity, continuous). That would make this question exam-standard.
As it is, it is too technical to be exam-standard.