

## MATHEMATICAL TRIPOS      Part III

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Wednesday 2 June, 2004    1.30 to 3.30

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## PAPER 31

## LARGE DEVIATIONS AND QUEUES

*Attempt **THREE** questions.*

*There are **four** questions in total.*

*The questions carry equal weight.*

*While rigorous answers are preferred, heuristic answers will still gain partial credit.*

<p><b>You may not start to read the questions printed on the subsequent pages until instructed to do so by the Invigilator.</b></p>
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**1** Let  $(X_n, n \in \mathbb{N})$  satisfy a large deviations principle in some space  $\mathcal{X}$  with good rate function  $I$ . Let  $f$  be a bounded continuous function  $\mathcal{X} \rightarrow \mathbb{R}$ .

(a) Let  $C_1, \dots, C_m$  be closed subsets of  $\mathcal{X}$  with  $\bigcup_i C_i = \mathcal{X}$ . Prove that

$$\limsup_{n \rightarrow \infty} \frac{1}{n} \log \mathbb{E} e^{nf(X_n)} \leq \max_{1 \leq i \leq m} \left\{ \sup_{x \in C_i} f(x) - \inf_{x \in C_i} I(x) \right\}.$$

(b) Let  $f(\mathcal{X})$  be contained in the interval  $[a, b]$ . Pick any  $\varepsilon > 0$  and define the closed intervals

$$D_i = [a + (i-1)\varepsilon, a + i\varepsilon], \quad i = 1, \dots, \lceil (b-a)/\varepsilon \rceil.$$

Let  $C_i = f^{-1}(D_i)$ . Using your answer to part (a), or otherwise, prove that

$$\limsup_{n \rightarrow \infty} \frac{1}{n} \log \mathbb{E} e^{nf(X_n)} \leq \sup_{x \in \mathcal{X}} [f(x) - I(x) + \varepsilon].$$

(c) Pick any  $\hat{x} \in \mathcal{X}$  and any  $\varepsilon > 0$ . Define the open interval

$$D = (f(\hat{x}) - \varepsilon, f(\hat{x}) + \varepsilon).$$

Let  $B = f^{-1}(D)$ . Using this set, or otherwise, prove that

$$\liminf_{n \rightarrow \infty} \frac{1}{n} \log \mathbb{E} e^{nf(X_n)} \geq f(\hat{x}) - I(\hat{x}) - \varepsilon.$$

(d) Deduce that

$$\lim_{n \rightarrow \infty} \frac{1}{n} \log \mathbb{E} e^{nf(X_n)} = \sup_{x \in \mathcal{X}} [f(x) - I(x)].$$

**2** A sequence of random variables  $(X_n, n \in \mathbb{N})$  taking values in a metric space  $\mathcal{X}$  is said to have *Hurstiness*  $H \in (0, 1)$  if the following three conditions are satisfied:

- $(X_n, n \in \mathbb{N})$  satisfies a large deviations principle with good rate function  $I$  at speed  $n^{2(1-H)}$ ;
- there is some  $\hat{x} \in \mathcal{X}$  such that  $0 < I(\hat{x}) < \infty$ ;
- there is some  $\mu \in \mathcal{X}$  such that  $I(x) = 0$  only if  $x = \mu$ .

Suppose  $(X_n, n \in \mathbb{N})$  has Hurstiness  $H$ . Let  $G > H$ ,  $G \in (0, 1)$ , and define the good rate function.

$$I'(x) = \begin{cases} 0 & \text{if } I(x) = 0 \\ \infty & \text{otherwise.} \end{cases}$$

(a) Prove that for any closed set  $C$

$$\limsup_{n \rightarrow \infty} \frac{1}{n^{2(1-G)}} \log \mathbb{P}(X_n \in C) \leq - \inf_{x \in C} I'(x).$$

(b) Using your answer to (a), or otherwise, show that if  $D$  is an open set containing  $\mu$  then

$$\mathbb{P}(X_n \notin D) \rightarrow 0.$$

Hence (or otherwise) show that for any open set  $E$

$$\liminf_{n \rightarrow \infty} \frac{1}{n^{2(1-G)}} \log \mathbb{P}(X_n \in E) \geq - \inf_{x \in E} I'(x).$$

(c) Suppose that  $(X_n, n \in \mathbb{N})$  has Hurstiness  $H$ , that  $(Y_n, n \in \mathbb{N})$  has Hurstiness  $G$ , that  $X_n$  is independent of  $Y_n$ , and that both take values in  $\mathbb{R}$ . Show that  $(X_n + Y_n, n \in \mathbb{N})$  has Hurstiness equal to the greater of  $H$  and  $G$ .

*Note. You should mention any general results you use, but you need not state them formally. Recall that  $(X_n, n \in \mathbb{N})$  is said to satisfy an LDP with rate function  $I$  and speed  $n^{2(1-H)}$  if for all measurable sets  $B \subset \mathcal{X}$*

$$\begin{aligned} - \inf_{x \in B^\circ} I(x) &\leq \liminf_{n \rightarrow \infty} \frac{1}{n^{2(1-H)}} \log \mathbb{P}(X_n \in B) \\ &\leq \limsup_{n \rightarrow \infty} \frac{1}{n^{2(1-H)}} \log \mathbb{P}(X_n \in B) \leq - \inf_{x \in \bar{B}} I(x). \end{aligned}$$

**3** Packets arrive at an Internet router as a Poisson process of rate  $\lambda$  packets per second. Each packet has a payload; payload sizes are independent of each other and of the arrival process, and have an exponential distribution with mean 1 kilobyte.

The router maintains two parallel queues, a ‘payload queue’ and a ‘header queue’. When a packet arrives, the payload is stored in the former, and a packet header is stored in the latter. Packets are served in the order they arrive. The payload queue is served at constant rate  $C$  kilobytes per second, and when the entire payload of a packet has been served then that packet’s header is removed from the header queue. Assume  $C > \lambda$ .

Both queues have finite space. The payload queue has space for 1000 kilobytes; the header queue has space for 1000 headers. As a queueing theorist, you are called in to advise on whether these are sensible choices.

(a) Let  $Q$  be the number of packet headers in the header queue. With reference to an  $M/M/1$  queue (or otherwise), estimate the probability that  $Q \geq q$ . (For modelling purposes, you can treat both queues as having infinite space.)

(b) The payload queue may be modelled by a discrete-time queue, with timeslots of length  $\delta$ , in which the number of packets arriving in each timeslot is a Poisson random variable with mean  $\delta\lambda$ , and the service offered in that timeslot is  $C\delta$ . Let  $R_\delta$  be the amount of work in this discrete-time queue. Estimate the probability that  $R_\delta \geq r$ . (Again, for modelling purposes, you can treat both queues as having infinite space.)

(c) Which queue is more likely to overflow? Give an intuitive explanation for your answer.

*Hint.* If  $N$  is a Poisson random variable with mean  $\lambda$  then  $\mathbb{E}t^N = e^{\lambda(t-1)}$ . If  $X$  is an exponential random variable with mean  $\lambda^{-1}$  then  $\mathbb{E}e^{\theta X} = \lambda/(\lambda - \theta)$ .

4 Consider a queue operating in continuous time, with constant service rate  $C$  and finite buffer  $B$ , with arrival process  $a \in \mathcal{C}_\mu$ . It is known that if  $\mu < C$  then the queue size at time 0 may be written as

$$\bar{q}(a) = \sup_{t \geq 0} \left\{ \left( \sup_{0 \leq s \leq t} x(-s, 0] \right) \wedge \left( B + \inf_{0 \leq s \leq t} x(-s, 0] \right) \right\}$$

where  $x(-s, 0] = a(-s, 0] - Cs$  and  $x \wedge y = \min(x, y)$ . When  $B = \infty$ , denote this function by  $q$ . It is also known that  $\bar{q}$  and  $q$  are continuous on  $(\mathcal{C}_\mu, \|\cdot\|)$ .

Suppose that  $(A^L, L \in \mathbb{N})$  satisfies a large deviations principle in  $(\mathcal{C}_\mu, \|\cdot\|)$  with good rate function

$$I(a) = \begin{cases} \int_{-\infty}^0 \Lambda^*(\dot{a}_s) ds & \text{if } a \text{ is absolutely continuous} \\ \infty & \text{otherwise} \end{cases}$$

for some strictly convex rate function  $\Lambda^*$  with  $\Lambda^*(\mu) = 0$ .

(a) Write down a large deviations principle for  $q(A^L)$ ; let it have rate function  $J$ . Also write down a large deviations principle for  $\bar{q}(A^L)$ ; let it have rate function  $\bar{J}$ .

(b) Show that  $\bar{q}(a) \leq q(a)$ . Hence (or otherwise) show that

$$\bar{J}(x) \geq \inf_{y \geq x} J(y).$$

(c) Show that  $J$  is increasing. Deduce that  $\bar{J}(x) \geq J(x)$ .

(d) Let  $x \leq B$ . Show that  $\bar{J}(x) \leq J(x)$ . *Hint. Let  $\hat{a}$  be the most likely path to attain  $q(a) = x$ . What is  $\bar{q}(\hat{a})$ ?*

(e) Deduce that  $\bar{q}(A^L)$  satisfies a large deviations principle with good rate function

$$\bar{J}(x) = \begin{cases} J(x) & \text{if } x \leq B \\ \infty & \text{otherwise.} \end{cases}$$

*Note. You may assume standard results about queues with infinite buffers.*