Model-comparison Games with Algebraic Rules

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> Newton Institute 1 March 2012

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First-order logic	Ehrenfeucht-Fraïssé game

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Fixed-point logic	Pebble game

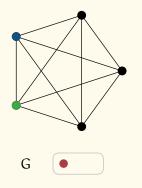
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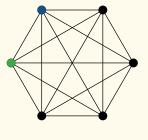
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??	Invertible-map game

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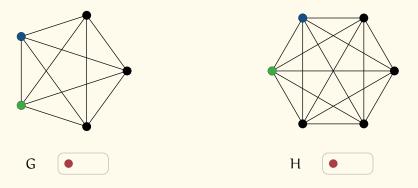






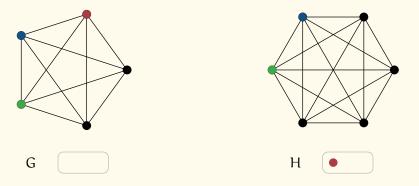
Spoiler Duplicator

 L^k - first order logic with variables x_1, \ldots, x_k .



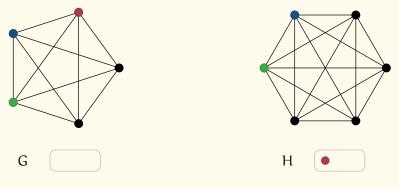
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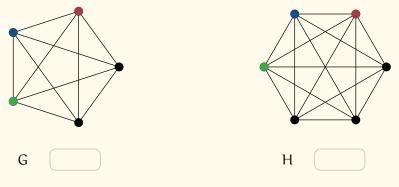
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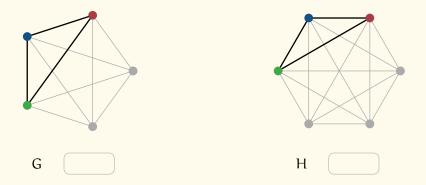
Spoiler Duplicator places the red pebble in H on some vertex

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Pebble mapping: partial isomorphism?

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• G_k has property \mathcal{P} but H_k does not; and

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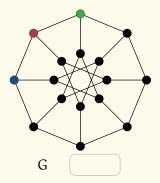
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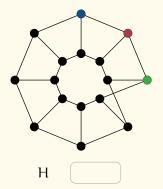
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- G_k has property \mathcal{P} but H_k does not; and
- ▶ Duplicator wins the k-pebble game on G_k and H_k.

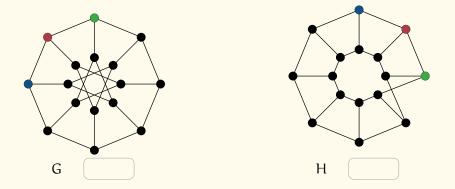
Pebble games for finite-variable counting logics

 C^k - extension of L^k with counting quantifiers: $\exists^{\geq i} x \, . \, \phi(x)$.

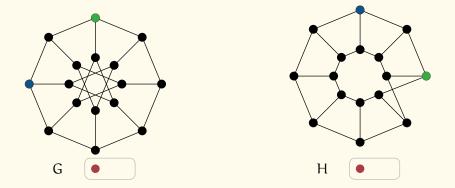




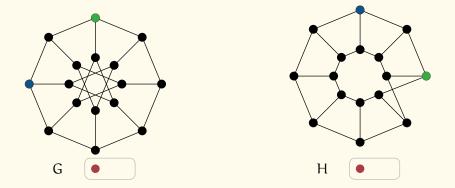
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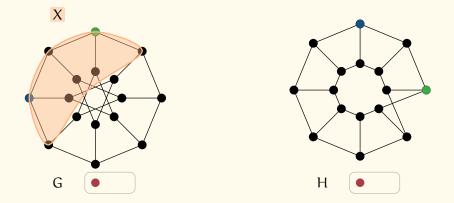
Spoiler chooses pebbles to remove from the two graphs Duplicator



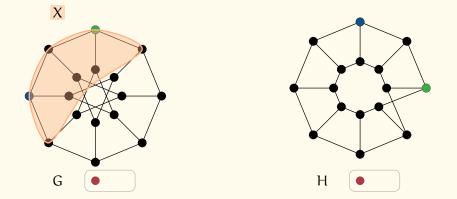
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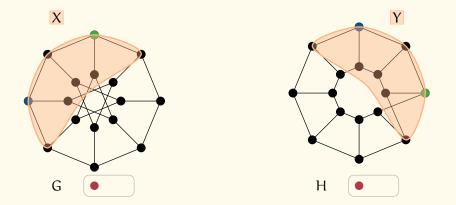
Spoiler chooses a set X of vertices in one of the structures Duplicator



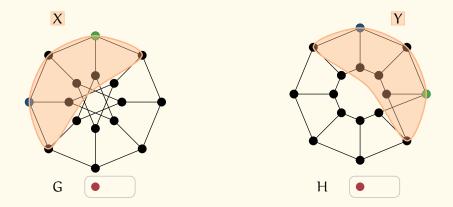
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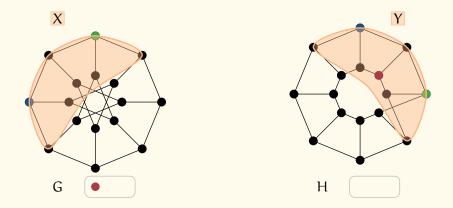
Spoiler Duplicator responds with a subset Y of H of the same cardinality



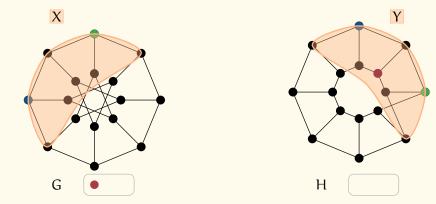
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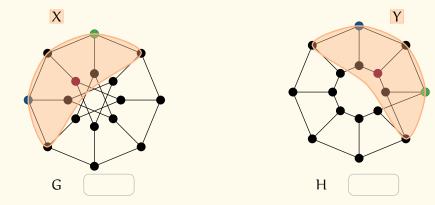
Spoiler places the red pebble in H on an element in Y Duplicator



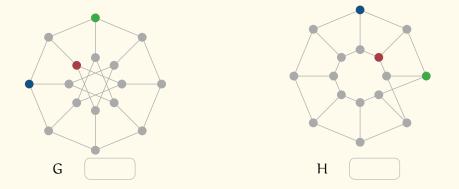
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Spoiler Duplicator places the red pebble in G on an element in X



Spoiler Duplicator places the red pebble in G on an element in X



Counting game characterises $C^k \rightsquigarrow$ game method for IFPC

Stronger logics for polynomial time

Basic problem not expressible in IFPC Solvability of linear equations over a finite field \rightsquigarrow Gaussian elimination

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Extend fixed-point logic with operators for expressing matrix rank over finite fields \rightsquigarrow IFPR

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 $\begin{array}{ll} \mbox{formula} & \phi(x,y) \\ \mbox{graph} & G = (V\!, E^G) \end{array}$

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 $\begin{array}{lll} \mbox{formula} & \phi(x,y) \\ \mbox{graph} & G = (V, E^G) & \leadsto & \begin{tabular}{c} M^G_\phi \\ \mbox{(over GF_2)} \end{array}$

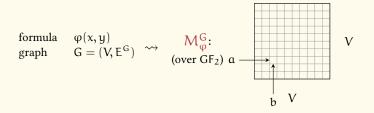


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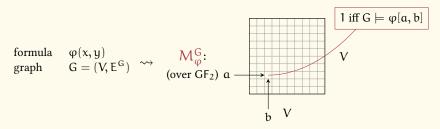
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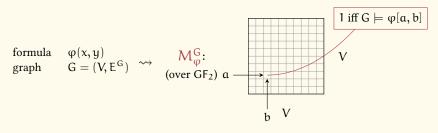
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Example

 $\phi(x,y):=Exy \ \rightsquigarrow \ M^G_\phi=adjacency\ matrix\ of\ G$

 R^k - extension of L^k with rank quantifiers: $rk^{\geq i}(x, y) \cdot \phi(x, y)$

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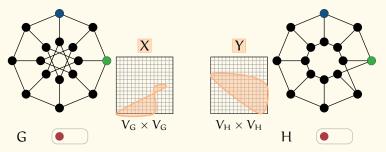
 $\rightsquigarrow \ rank \ of \ M^G_{\phi} \ over \ GF_2$

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First attempt: Extend the Immerman-Lander cardinality game

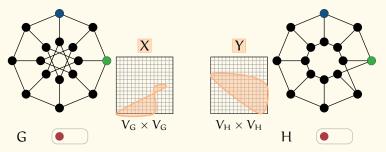
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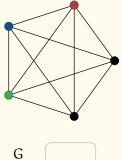
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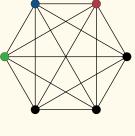


Problem: rank is not monotone!

Partition game for L^k - illustrates main idea

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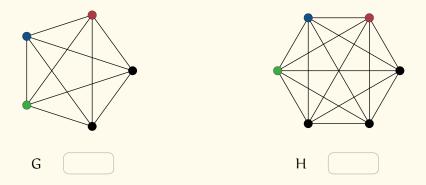




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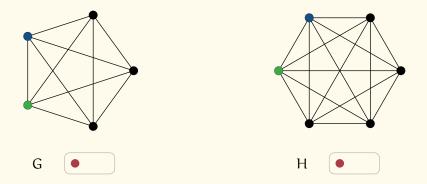
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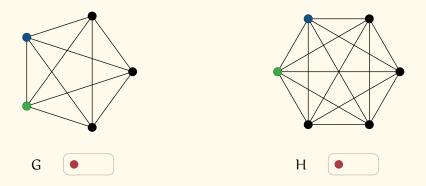
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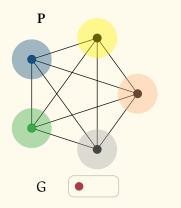
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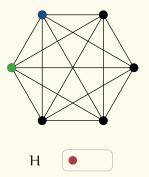
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Spoiler Duplicator gives a partition P of V_G, \ldots

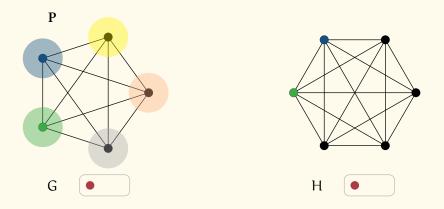
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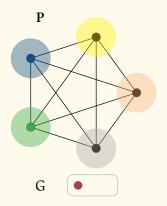


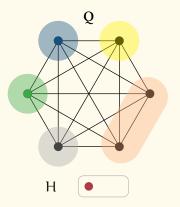


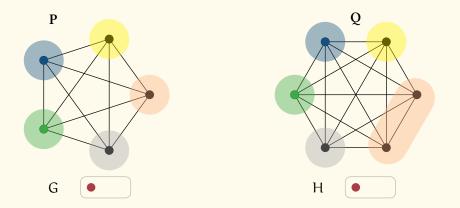
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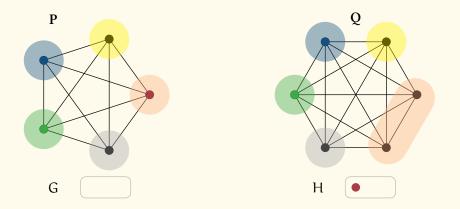
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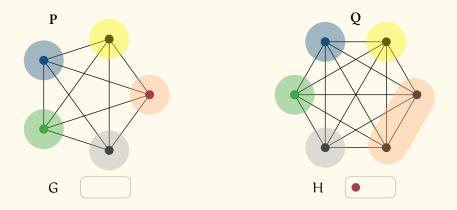


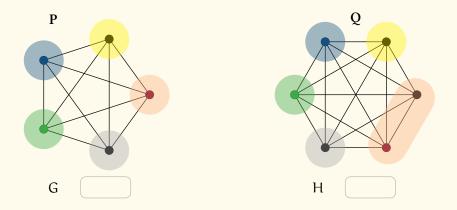




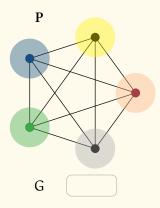


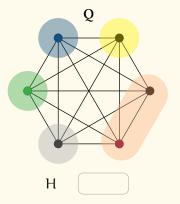




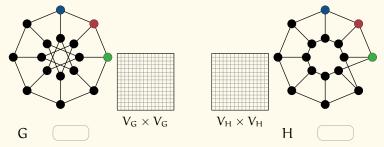


A new pebble-game protocol: partition games To get a partition game for $C^k \rightsquigarrow$ for all $X \in P$, ||X|| = ||f(X)||

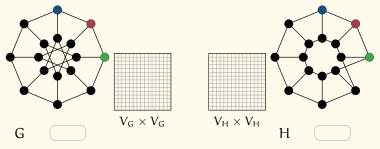




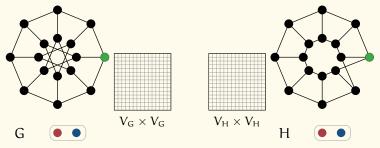
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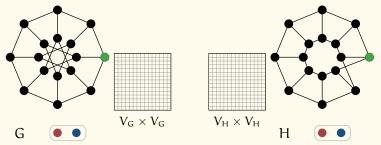
Spoiler Duplicator



Spoiler removes two pairs of corresponding pebbles Duplicator

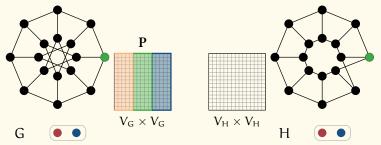


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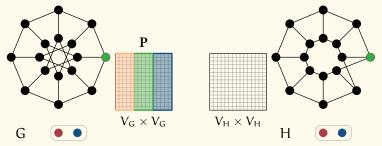
Spoiler

Duplicator gives a partition **P** of $V_G \times V_G$, ...

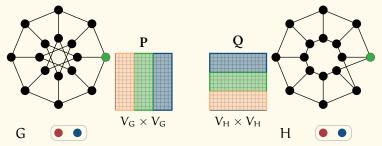


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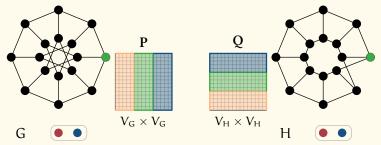
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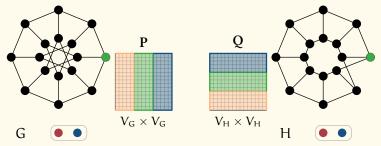


Spoiler

Duplicator

a partition Q of $V_H \times V_H$, and $f : P \rightarrow Q$, such that for all $X \subseteq P$:

$$\operatorname{rk}\left(\sum_{M\in X} P\right) = \operatorname{rk}\left(\sum_{M\in X} f(M)\right)$$



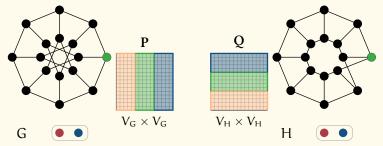
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Can we decide in polynomial time who wins the game? ~ "yes" for standard pebble and cardinality games

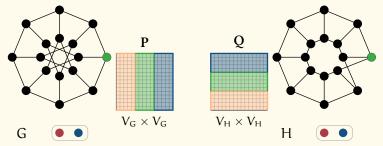


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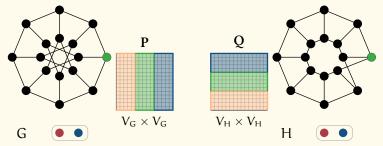


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Duplicator

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 $\sum_{M \in X} P$ and $\sum_{M \in X} f(M)$ are equivalent

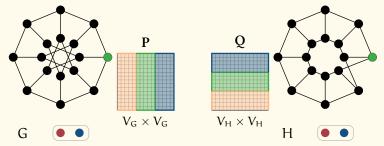


Spoiler

Duplicator

a partition Q of $V_H \times V_H$, and $f : P \rightarrow Q$, such that for all $X \subseteq P$:

$$\sum_{M \in X} P$$
 and $\sum_{M \in X} f(M)$ are similar

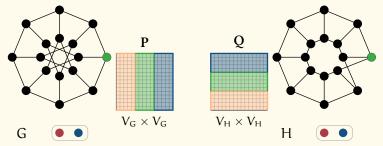


Spoiler

Duplicator

a partition Q of $V_{\mathsf{H}} \times V_{\mathsf{H}},$ and $\mathsf{f}: P \to Q,$ such that:

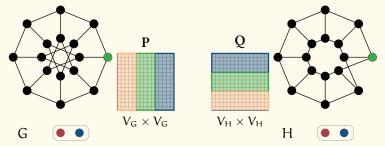
 $(M)_{M \in \mathbf{P}}$ and $(f(M))_{M \in \mathbf{P}}$ are simultaneously similar



Spoiler

> $(M)_{M \in \mathbf{P}}$ and $(f(M))_{M \in \mathbf{P}}$ are simultaneously similar

 \leadsto we can decide who wins this game in polynomial time



Spoiler

> $(M)_{M \in \mathbf{P}}$ and $(f(M))_{M \in \mathbf{P}}$ are simultaneously similar

→ we can decide who wins this game in polynomial time
→ Application: family of algorithms for testing graph isomorphism

From logics to games – and back again?

- Does the "simultaneous-similarity game" correspond to a natural logic?
- ► Duplicator wins the simultaneous-similarity game ⇒ Duplicator wins the matrix-rank game. Converse?