## COMPUTER SCIENCE TRIPOS Part IA - 2018 - Paper 1

## 10 Algorithms (RKH-DJW)

(a) Let dijkstra_path $(g, a, b)$ be an implementation of Dijkstra's shortest path algorithm that returns the shortest path from node $a$ to node $b$ in a graph $g$. Prove that the implementation can safely terminate when it first encounters node $b$.
(b) Consider all paths in a graph from $a$ to $b$, ordered from shortest to longest. Assuming $p=$ dijkstra_path $(g, a, b)$ is the first path in this collection, an algorithm to find the second path considers deviations from the vertices of $p$. An algorithm to do this is given below.

```
function second_path(Graph g, Vertex a, Vertex b):
    p = dijkstra_path(g,a,b)
    best_so_far = []
    for i = 1 to len(p)-1:
        t = p[:i] # First i elements of p
        c = g.get_edge_weight(p[i], p[i+1])
        g.set_edge_weight(p[i], p[i+1], infinity)
        t.append(dijkstra_path(g,p[i],b))
        if (len(best_so_far) == 0 or
            cost(t) < cost(best_so_far)):
            best_so_far = t
        g.set_edge_weight(p[i], p[i+1], c)
    return best_so_far
```

(i) Show the steps of this algorithm on the following graph, from A to B .

[5 marks]
(ii) What is the asymptotic complexity of this algorithm in terms of the number of edges, $E$, and the number of vertices, $V$ ? Assume the implementation of Dijkstra's algorithm uses a priority queue based on a Fibonacci heap.
[4 marks]
(iii) Show how to adapt this algorithm to find the top- $k$ shortest paths in the collection. State the complexity of the adapted algorithm.

