COMPUTER SCIENCE TRIPOS Part II – 2017 – Paper 8

15 Types (AMP)

 (a) (i) Explain the difference with respect to polymorphic occurrences of variables between the way let-bound and lambda-bound variables are treated in languages featuring ML-style polymorphism. Illustrate your answer by explaining why in the Mini-ML language one of the two expressions

$$\lambda g \left(\operatorname{let} f = \lambda x \left(\operatorname{let} u = g x \operatorname{in} x \right) \operatorname{in} f f \right)$$
$$\operatorname{let} f = \lambda x \left(\operatorname{let} u = x \operatorname{in} x \right) \operatorname{in} f f$$

is typeable and the other is not. As part of your answer you should define what it means for a type to be a specialisation of a type scheme and give the Mini-ML typing rules for variables, function abstractions, function application and let-expressions. [14 marks]

- (*ii*) What property of the type system is lost if one attempts to treat polymorphic occurrences of lambda-bound variables in the same way as let-bound ones? [1 mark]
- (b) Give the typing rules for expressions associated with ML-style reference types, namely reference creation ref M, reference look-up !M and assignment M := M' expressions. What is the *value-restriction* that is imposed on the typing rule for let-expressions in the presence of these forms of expression and what is its purpose? [5 marks]