## COMPUTER SCIENCE TRIPOS Part Ib - 2014 - Paper 5

## 8 Concurrent and Distributed Systems (RNW)

(a) Monitors are a programming primitive linking data with two synchronization types: mutual exclusion and condition synchronisation. Which is provided implicitly; which is provided explicitly?
(b) Describe two ways in which Monitors and Conditional Critical Regions differ.
[2 marks]
(c) The object-oriented programming style encouraged by Monitors has many benefits as the number of data types and locks increases in the system.
(i) Placing all data in a single Monitor may improve program correctness. Explain why this might have undesirable performance effects. [1 mark]
(ii) One problem that can arise when using multiple locks is deadlock, which can be prevented by imposing a partial order on locks. Describe the implications this has for code structure when using Monitors.
(iii) Explain why Java's Monitor feature does not necessarily impose this code structure.
(d) Condition variables allow condition satisfaction to be signalled between threads. Explain the difference between Hoare's signal-and-wait and Mesa's signal-andcontinue in terms of mutual exclusion and scheduling.
(e) Consider the (incorrect) pseudocode on the next page:
(i) Describe and justify minimal modifications to this code, referencing line numbers, in order to make it correct in the presence of Hoare signal-and-wait semantics.
(ii) Describe and justify minimal modifications to this code, referencing line numbers, in order to make it correct in the presence of Mesa signal-and-continue semantics.
[4 marks]
[continued ...]

```
monitor ProducerConsumer {
    int in, out, buf[N];
    condition notfull, notempty;
    procedure produce(item) {
        if ((in-out) == N)
            wait(notfull);
        buf[in % N] = item;
        if ((in-out) == 0)
            signal(notempty);
        in = in + 1;
    }
    procedure int consume() {
        if ((in-out) == 0)
            wait(notempty);
        item = buf[out % N];
        if ((in-out) == N)
            signal(notfull);
        out = out + 1;
    }
    /* init */ { in = out = 0; }
}
```

