COMPUTER SCIENCE TRIPOS Part IB – 2014 – Paper 5

8 Concurrent and Distributed Systems (RNW)

- (a) Monitors are a programming primitive linking data with two synchronization types: mutual exclusion and condition synchronisation. Which is provided implicitly; which is provided explicitly? [1 mark]
- (b) Describe two ways in which Monitors and Conditional Critical Regions differ. [2 marks]
- (c) The object-oriented programming style encouraged by Monitors has many benefits as the number of data types and locks increases in the system.
 - (i) Placing all data in a single Monitor may improve program correctness. Explain why this might have undesirable performance effects. [1 mark]
 - (ii) One problem that can arise when using multiple locks is deadlock, which can be prevented by imposing a partial order on locks. Describe the implications this has for code structure when using Monitors. [2 marks]
 - (iii) Explain why Java's Monitor feature does not necessarily impose this code structure. [2 marks]
- (d) Condition variables allow condition satisfaction to be signalled between threads. Explain the difference between Hoare's signal-and-wait and Mesa's signal-and-continue in terms of mutual exclusion and scheduling. [4 marks]
- (e) Consider the (incorrect) pseudocode on the next page:
 - (i) Describe and justify minimal modifications to this code, referencing line numbers, in order to make it correct in the presence of Hoare signal-and-wait semantics. [4 marks]
 - (ii) Describe and justify minimal modifications to this code, referencing line numbers, in order to make it correct in the presence of Mesa signal-and-continue semantics. [4 marks]

[continued ...]

```
1: monitor ProducerConsumer {
2:
       int in, out, buf[N];
       condition notfull, notempty;
3:
4:
5:
       procedure produce(item) {
6:
         if ((in-out) == N)
7:
           wait(notfull);
8:
         buf[in % N] = item;
         if ((in-out) == 0)
9:
10:
           signal(notempty);
         in = in + 1;
11:
12:
       }
13:
       procedure int consume() {
14:
15:
         if ((in-out) == 0)
16:
           wait(notempty);
17:
         item = buf[out % N];
18:
         if ((in-out) == N)
           signal(notfull);
19:
20:
        out = out + 1;
21:
22:
23:
       /* init */ { in = out = 0; }
24: }
```