## 2003 Paper 11 Question 2

## Compiler Construction

A programming language has commands $C$ and expressions $E$ which may involve the terminals $I$ (identifiers) and $N$ (integer constants). Its grammar $\mathcal{G}$ (with start symbol $S$ ) is given by

$$
\begin{aligned}
& S::=C \text { eof } \\
& C::=I=E \mid \text { if } E \text { then } C \mid \text { if } E \text { then } C \text { else } C \\
& E::=I|N| E+I
\end{aligned}
$$

Construct
(a) a recursive descent parser, and
(b) the characteristic finite state machine (CFSM) of a $\operatorname{LR}(k), \operatorname{SLR}(k)$ or $\operatorname{LALR}(k)$ parser,
explaining carefully whether there are any problems in the grammar $\mathcal{G}$ (and if so how you resolved them) and in adapting the grammar for recursive descent parsing and in adopting the CFSM so as to be suitable for $\operatorname{LR}(k), \operatorname{SLR}(k)$ or $\operatorname{LALR}(k)$ parsing. It is not necessary to consider how the corresponding parse tree is constructed.

