## **Optimising Compilers**

Briefly summarize the main concepts of strictness analysis including the kind of languages to which it applies, and the way in which both system-provided and user-defined functions f yield strictness properties  $f^{\#}$  (relate the types of f and  $f^{\#}$ ). [6 marks]

Give the strictness functions corresponding to the following ternary functions:

$$(a)$$
 f1(x,y,z) = x\*y + z

(c) 
$$f3(x,y,z) = pif x=9$$
 then y else z

where  $pif e_1$  then  $e_2$  else  $e_3$  is the *parallel conditional*: it behaves similarly to the standard conditional in that if  $e_1$  evaluates to true or false then it yields  $e_2$ or  $e_3$  as appropriate; however, evaluation of  $e_2$  and  $e_3$  occurs concurrently with  $e_1$ to allow the pif construct also to terminate with the value of  $e_2$  when  $e_2$  and  $e_3$ both terminate with equal values (even if  $e_1$  computes forever).

Comment briefly how your strictness property for f1 would change if the multiplication returned zero without evaluating the other argument in the event that one argument were zero. [7 marks]

Let g,  $h_1$  and  $h_2$  be binary functions and recall the definition of function composition:

$$g \circ \langle h_1, h_2 \rangle = \lambda(x, y) \cdot g(h_1(x, y), h_2(x, y)).$$

Define three such functions in an ML-like syntax (whose arguments and results are integers) and which have the property that

$$(g \circ \langle h_1, h_2 \rangle)^{\#} \neq g^{\#} \circ \langle h_1^{\#}, h_2^{\#} \rangle.$$

[Hint: you might find it helpful to think of a solution where g may ignore one of its arguments but always does when composed with  $\langle h_1, h_2 \rangle$ .] Comment whether this inequality means that  $g^{\#} \circ \langle h_1^{\#}, h_2^{\#} \rangle$  fails to be a safe strictness property for  $g \circ \langle h_1, h_2 \rangle$ . [7 marks]