

10 Optimising Compilers (AM)

- (a) Annotate the uses of a given variable with ‘*D*’ (definition), ‘*R*’ (reference) and ‘*U*’ (undefinition—e.g. entering or leaving scope). Explain why three pairs of such uses may be regarded as *anomalous* data flow, giving brief example programs illustrating how each of these anomalies can represent a programmer error. Your answer should make use of the concept of path in a dataflow graph. [3 marks]
- (b) Give a single program containing all three of the above anomalies and justify it as *not* containing any programmer error. [3 marks]
- (c) By analogy with “on all paths” and “on some path” in dataflow analysis refine your definition of dataflow anomaly in part (a) into that of ‘must-anomaly’ and ‘may-anomaly’ indicating which, if either, of these corresponds to the previous definition. How many must-anomalies does your program in part (b) contain? [3 marks]
- (d) Give *two* sets of dataflow equations, in the style used to define “live variables” or “available expressions”, which each calculate a set of variables for every program point. Explain how these can respectively be used to issue compiler warning messages of the form “*variable ‘x’ may be read before being set*” and “*variable ‘x’ is definitely read before being set*”. For both forms, state the initialisation of these sets of variables to be used when solving the dataflow equations. [7 marks]
- (e) Now suppose the language supports indirect assignment to address-taken variables as in C. Explain briefly the refinements necessary to the above analyses. [2 marks]
- (f) The wording (‘is definitely read’) of the second message in part (d) may be criticised when considering the program

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{ int x,y=0; while (...) y=f(y); print x; }
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Comment briefly on this claim. [2 marks]