

## 2011 Paper 9 Question 13

### Topics in Concurrency

This question concerns  $\text{CTL}^-$ , a variant of CTL in which assertions are of the form:

$$A := T \mid A_0 \vee A_1 \mid A_0 \wedge A_1 \mid \neg A \mid \langle \lambda \rangle A \mid \langle - \rangle A \mid \mathbf{AG} A \mid \mathbf{EG} A$$

where  $\lambda$  ranges over action names. Recall that a *path*  $\pi$  from state  $s$  is a maximal sequence of states  $\pi = (\pi_0, \pi_1, \pi_2, \dots)$  such that  $s = \pi_0$  and  $\pi_i \xrightarrow{\cdot} \pi_{i+1}$  for all  $i$ . The logical connectives have the standard interpretation and  $T$  represents “true”. The interpretation of the modalities is:

$$\begin{aligned} s \models \langle \lambda \rangle A & \text{ iff } \text{there exists } s' \text{ such that } s \xrightarrow{\lambda} s' \text{ and } s' \models A \\ s \models \langle - \rangle A & \text{ iff } \text{there exists } s' \text{ such that } s \xrightarrow{\cdot} s' \text{ and } s' \models A \\ s \models \mathbf{EG} A & \text{ iff } \text{for some path } \pi \text{ from } s, \text{ we have } \pi_i \models A \text{ for all } i \\ s \models \mathbf{AG} A & \text{ iff } \text{for all paths } \pi \text{ from } s, \text{ we have } \pi_i \models A \text{ for all } i \end{aligned}$$

- (a) What is the interpretation of the CTL modality  $\mathbf{E}[A_0 \mathbf{U} A_1]$ ? How can it be used to express the  $\text{CTL}^-$  modality  $\mathbf{AG} A$ ? [4 marks]
- (b) Consider the following three formulae:

$$A_1 : \mathbf{AG} \langle a \rangle T \quad A_2 : \neg \mathbf{AG} \neg \langle b \rangle T \quad A_3 : \neg \mathbf{EG} \langle b \rangle T$$

- (i) For each of the following two transition systems, state which of  $A_1$ ,  $A_2$  and  $A_3$  are satisfied in the initial state.



[4 marks]

- (ii) Draw a transition system with an initial state that satisfies  $A_1 \wedge A_2 \wedge A_3$ . [3 marks]

- (c) Give a modal- $\mu$  formula that corresponds to the  $\text{CTL}^-$  formula  $\mathbf{AG} \langle a \rangle T$ . [2 marks]

- (d) Let the function  $\varphi$  on sets of states of a transition system be defined as:

$$\varphi(X) \stackrel{\text{def}}{=} \langle a \rangle T \vee \langle - \rangle X$$

Show by induction on  $n \geq 1$  that

$$s \models \varphi^n(\emptyset) \quad \text{iff} \quad \text{there exists } m \leq n \text{ and states } s_1, \dots, s_m, s' \text{ such that } s = s_1 \xrightarrow{\cdot} \dots \xrightarrow{\cdot} s_m \xrightarrow{a} s'$$

Deduce that  $s \models \mu X. \langle a \rangle T \vee \langle - \rangle X$  in a finite-state transition system if, and only if,  $s \models \neg \mathbf{AG} \neg \langle a \rangle T$ . [7 marks]