Compiler Construction

(a) Explain how a parse tree representing an expression can (i) be converted into stack-oriented intermediate code and then (ii) be translated into simple machine code for a register-oriented architecture (e.g. ARM or IA32) on an instruction-by-instruction basis. Also indicate how this code might be improved to remove push–pop pairs introduced by (ii). Your answer need only consider expression forms encountered in the expression:

\[ h(a, g(b), c) \ast 3 + d \]

[12 marks]

(b) In Java, expressions are evaluated strictly left-to-right. Consider compiling the function \( f \) in the following Java class definition:

```java
class A {
    static int a,b;
    void f() { ... <<C>> ... }
    int g(int x) { ... a++; ... }
};
```

Indicate what both the intermediate code and (improved as above) target code might be for \( <<C>> \) for the cases where \( <<C>> \) is:

(i) \( b = g(7) + a; \)

(ii) \( b = a + g(7); \)

(iii) \( b = (-g(7)) + a; \)

(iv) \( b = a - g(7); \)

Comment on any inherent differences in efficiency at both the intermediate code and target code levels.

[8 marks]