

## 2003 Paper 9 Question 7

### Optimising Compilers

(a) Summarise briefly the principles of *strictness* analysis, including descriptions of:

- (i) the space of values used for analysis-time representation of a  $k$ -argument, 1-result function in the source language;
- (ii) how a built-in function is given an abstract meaning;
- (iii) how a recursive user-defined function is given an abstract meaning (it is acceptable to do this part by example);
- (iv) the machine-level benefit of the associated optimisation.

[8 marks]

(b) A problem amenable to similar treatment is that of *escape analysis*. Here we have a call-by-value language with `cons` and the question to be answered is “whether a value containing a cons-node passed as argument to a function *may* be returned (‘escape’) as part of the function’s result”.

- (i) Choose (and state clearly) an appropriate set of abstract values and abstractions of functions to formalise the problem of escape analysis for a simple first-order language with integers and simple integer lists (but not lists of lists). Also give abstract interpretations of `if-then-else`, `+`, `cons`, `hd` and `tl`.

[Hint: to manage this system without using static types, you might best assume that *nil* is treated as 0, and that any type-error (dynamically detected) such as `cons(1,nil)+3`, `tl(3)` and even (because of the ‘no-lists-of-lists’ rule) `cons(cons(1,nil),nil)` gives result 0.] [8 marks]

- (ii) Give without proof abstract meanings (resulting from your system) of the following functions:

```
f(x,y,z) = cons(hd(tl(x)), if hd(x) then y else tl(z))
g(x,y) = if x=0 then 0 else cons(hd x, g(tl x, y))
h(x,y) = if x=0 then x else cons(hd x, h(tl x, y))
k(x,y) = if x=0 then y else cons(hd x, k(tl x, y))
```

[4 marks]