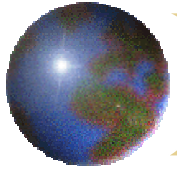


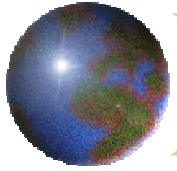
Digital Communications II: IPv6

Richard Black



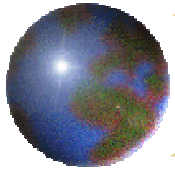
Outline

- What is wrong with IPv4
 - Some IPv4 technology
- Key differences in IPv6
 - Basic facts
 - Plus facts to understand ...
 - Basic Principles

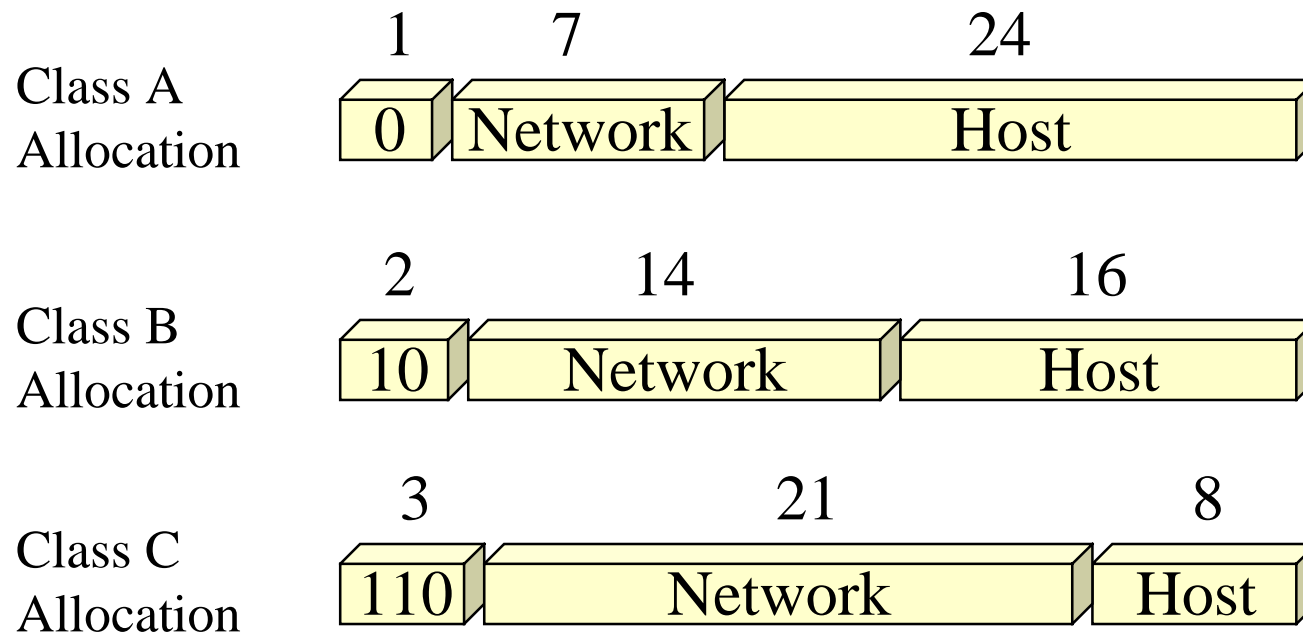


Problems with IPv4

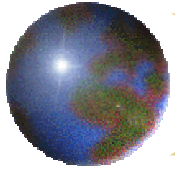
- ❁ Address space running out
 - ❑ 32bits, poor density, growth
 - ❑ Mobile Phones
- ❁ Routeing tables getting too big
 - ❑ Core routers are default-free
 - ❑ Need an entry per network
 - ❑ Insufficiently hierarchical
- ❁ Difficulty of configuration



IPv4 Address structure

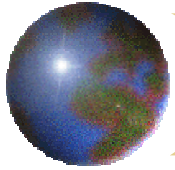


- 2^{21} means only two million organisations and core routers need to know about them all

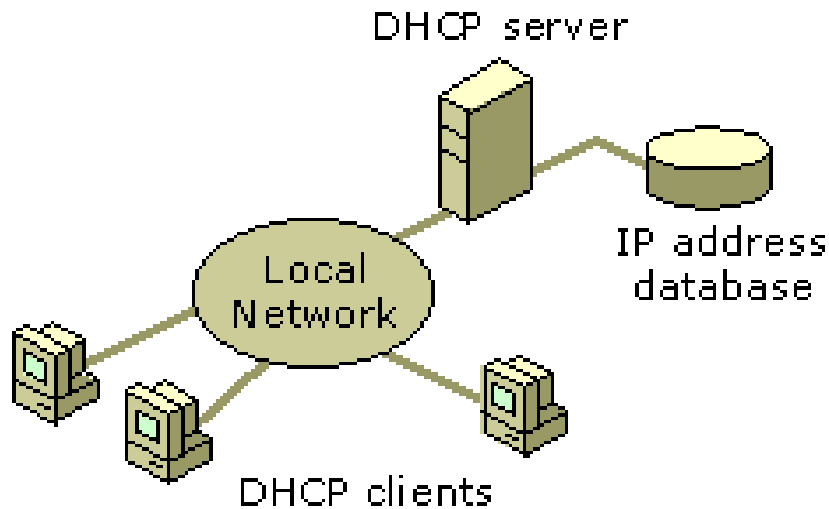


Workarounds with IPv4

- RFC1918 private network addresses
 - 10/8, 172.16/12, 192.168/16
- CIDR (classless inter-domain routing)
 - Too little too late. Consider JANET:
 - Single AS786. 237 different routing entries
 - Still 70000+ routing entries (see p65)
- NAT
- DHCP

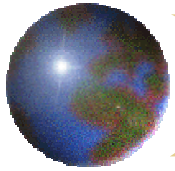


DHCP in one slide

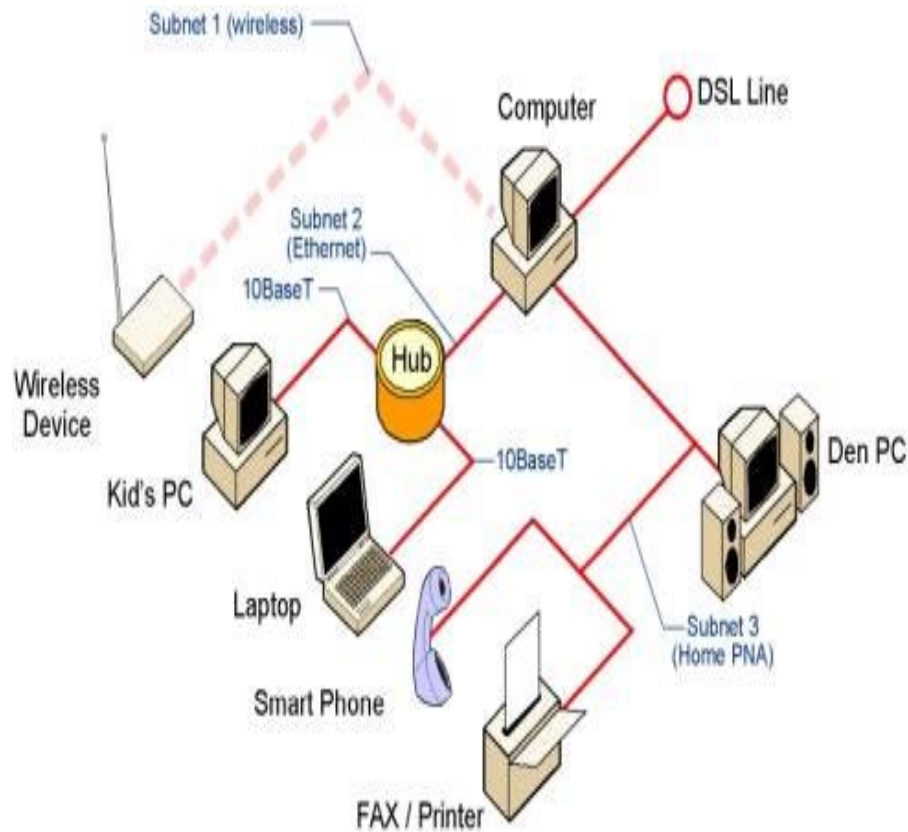


- ➊ Avoid static IP configuration
 - ▣ Assign dynamically

- ➋ Hosts request details
 - ▣ Identified by LL address
- ➌ Server replies
 - ▣ IP address and mask
 - ▣ Default router
 - ▣ DNS servers
 - ▣ Etc.
- ➍ IP address on lease
 - ▣ Client renews
- ➎ Naming Problems
 - ▣ Both DNS and SAP

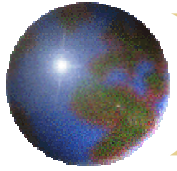


NAT in one slide



- DHCP+NAT = "Internet connection sharing"

- NAT box is gateway
 - It has a real address
 - Its address gets shared
- Translates:
 - (local,port) (its, port')
 - Maintains a mapping
 - Has to understand some protocols (ftp)
- Also APP level relay
 - Web proxy / cache
- Naming problems

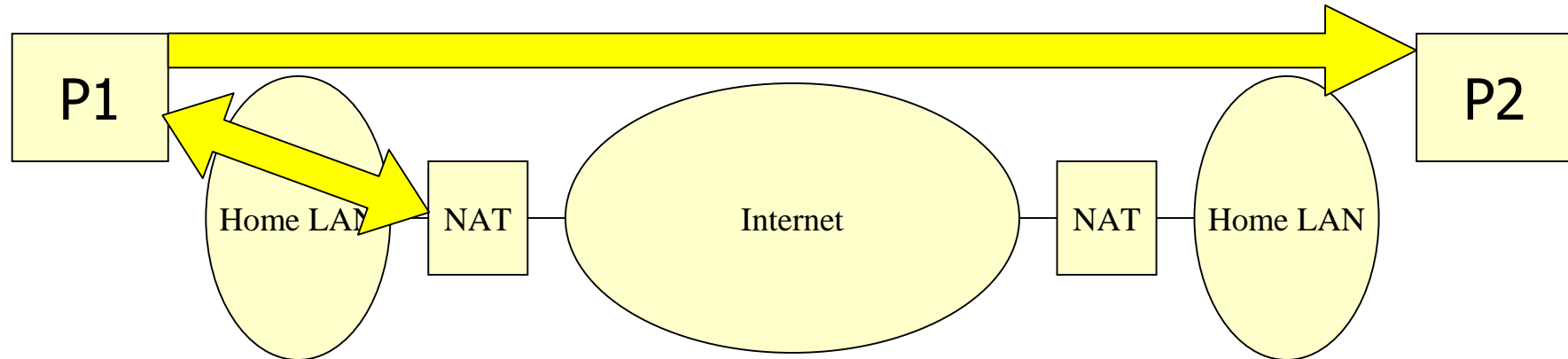


Advantages of IPv6

- ⊕ Enormous address space: 128 bits
- ⊕ Strict topology-based allocation
- ⊕ More scaleable routeing
- ⊕ Easier configuration
- ⊕ Better security
- ⊕ 15 years of experience



Problem 1: Peer-to-peer RTP audio example



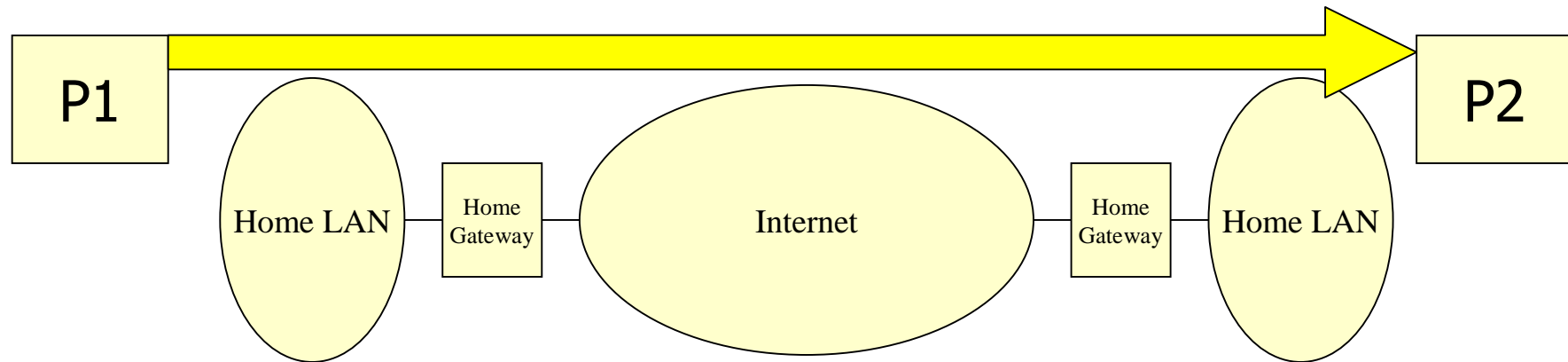
With NAT:

- Need to learn the address "outside the NAT"
- Provide that address to peer
- Need either NAT-aware application, or application-aware NAT
- May need a third party registration server to facilitate finding peers



Solution 1: Peer-to-peer RTP

audio example

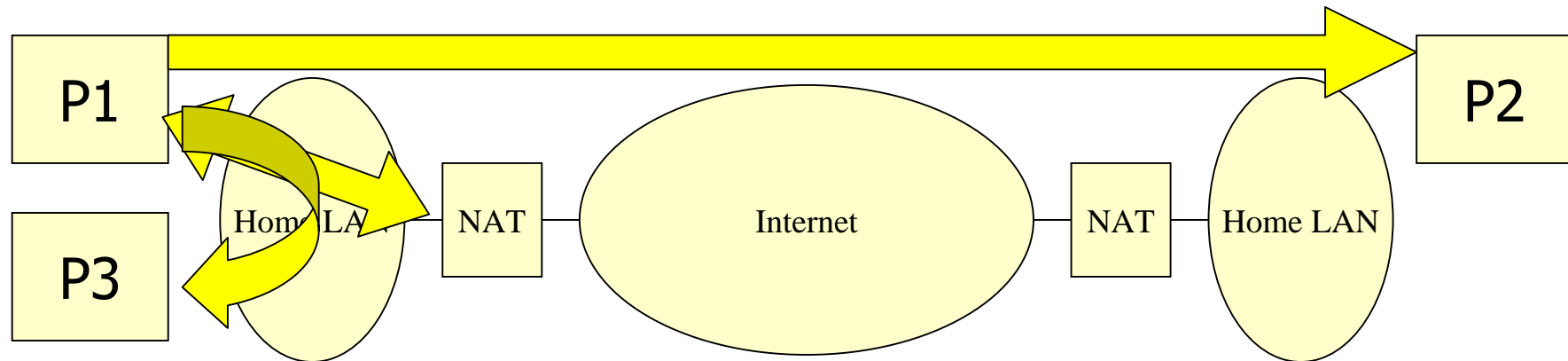


✪ With IPv6:

▣ Just use IPv6 address



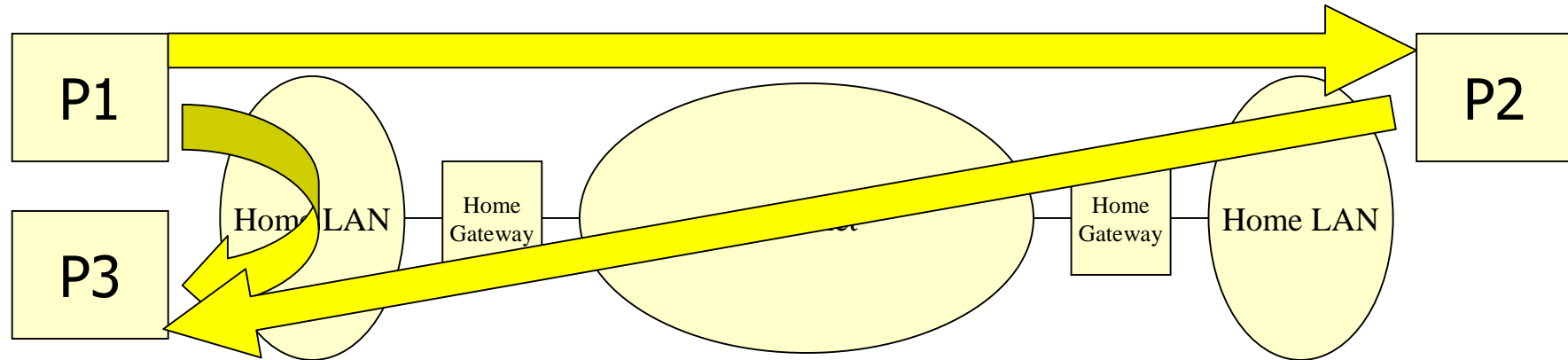
Problem 2: Multiparty Conference Example



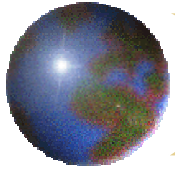
- ✿ With NAT, complex and brittle software:
 - ❑ 2 Addresses, inside and outside
 - ❑ P1 provides "inside address" to P3, "outside address" to P2
 - ❑ Need to recognize inside, outside
 - ❑ P1 does not know outside address of P3 to inform P2



Solution 2: Multiparty IPv6 Conference Example

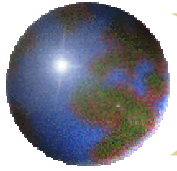


- With IPv6:
 - Just use IPv6 addresses



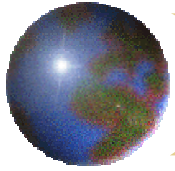
What does an address look like

- 128 bits broken up into 16 bit groups
 - Each 16 bit group written in hex
 - No 0x, remove leading zeros
 - Groups separated by ':' colon
- A single run of 0: entries elided with ::
- Last two groups allow IPv4 notation
- Heavy use of prefix '/' notation
 - 2001:618:1::2:2a0 ::/0
 - 2002:80e8:10:0:/64 fe80::/96

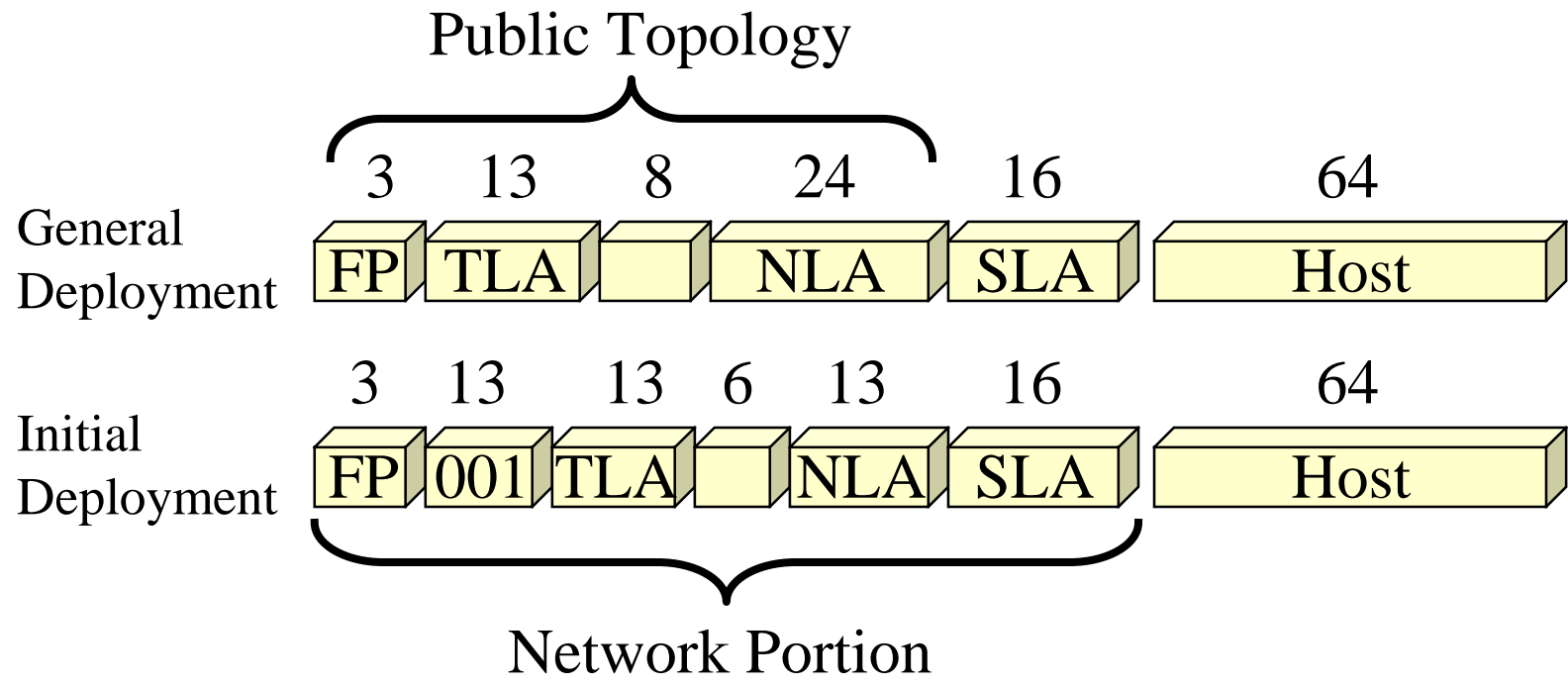


Topology-based allocation

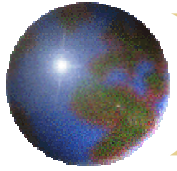
- Network Service Providers get allocations from registries
 - Backbone routing tables are small
- NSPs make allocations to ISPs
 - ISPs have small routing tables
- ISPs make allocations to companies
 - Addresses depend on service provider
 - May have to change addresses with provider
 - Multi-provider sites need research



IPv6 Address structure

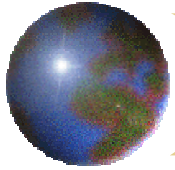


- | | | | |
|-----|---------------------|-----|----------------------|
| FP | Format Prefix | NLA | Next Level Aggregate |
| TLA | Top Level Aggregate | SLA | Site Level Aggregate |



More scalable routing

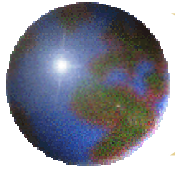
- There are more network bits
- But any router has only to look at a few
 - Usually only one of TLA or NLA or SLA
- Compare JANET again
 - One TLA for JANET
 - One NLA per institution
 - Institution uses SLA like current “subnets”



Packet Differences

- Designed to be easier on routers
 - Header is fixed size
 - No fragmentation at routers
- Extension headers
 - Either fixed size or have own length

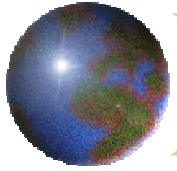
Base Header	Extension Header 1	...	Extension Header N	Data
-------------	--------------------	-----	--------------------	------



Base Header

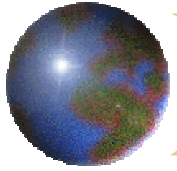
vers	class	flow label	
length		next	hop lim
Source Address			
Destination Address			

- Compare with IPv4 (page 3)
 - Bigger
 - Simpler
- Extension for anything else



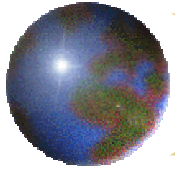
Extension Headers

- Included in specific order
 - Hop by Hop
 - Destination (for intermediates)
 - Routing
 - Fragment
 - Authentication
 - Encapsulating Security Payload
 - Destination (for final destination)



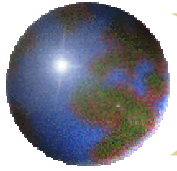
Path MTU discovery

- Recommended in IPv4
- Necessary in IPv6
- Relies on “too big” ICMP
- Another example of “soft state”



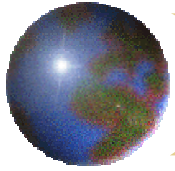
IPv6 Neighbour Discovery

- Replaces ARP, ICMP redirect etc.
- Used for:
 - Router discovery
 - Parameter and prefix discovery
 - Address resolution
 - Address auto-configuration
- Send Neighbour solicitation for:...x:yz to ff02::1:ff x:yz

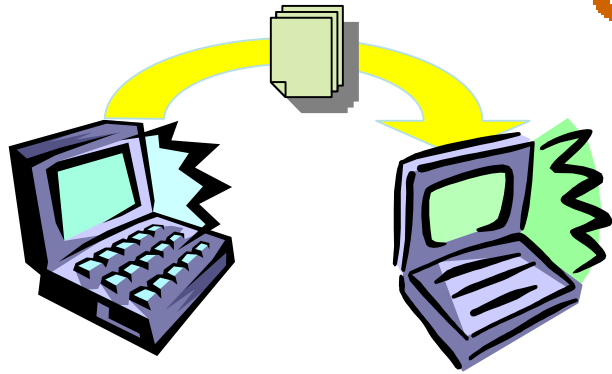


Address auto-configuration

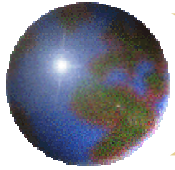
- Router advertises prefix (/64)
- Machines chose an address using MAC
- Look for duplicate before using
 - Send Neighbour solicitation for intended addr
- Also have link local addresses (fe80::)
- Possible privacy concerns
 - An advantage to NAT after all?
 - Work on anonymous addresses



Problem 3: Ad-hoc networking

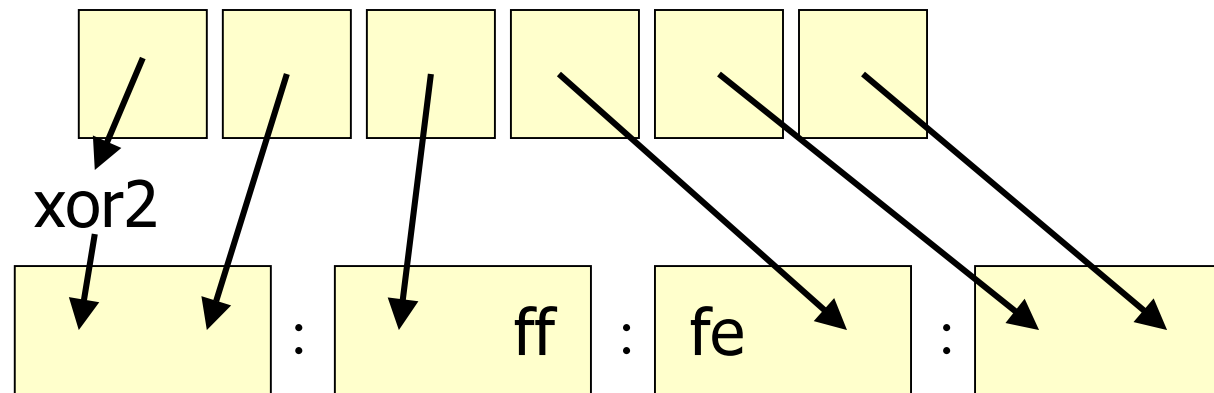


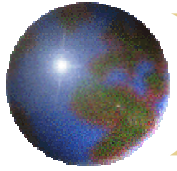
- IPv4: media lock + 63 sec.
 - Try DHCP
 - Wait for timeout
 - Select Link Local address
 - From 169.254/16
 - See draft-ietf-zeroconf-ipv4-linklocal
 - Conflict detect
- IPv6: media lock + 1 sec.
 - Configure using MAC
 - Conflict detect



IPv6 on Ethernet (802.2)

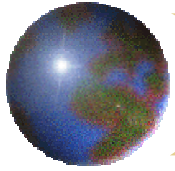
- New Ethernet type (0x86dd)
- IPv6 multicast packets FF.....::wx:yz sent to Ethernet 33-33-w-x-y-z
- Auto-configuration from Ethernet address



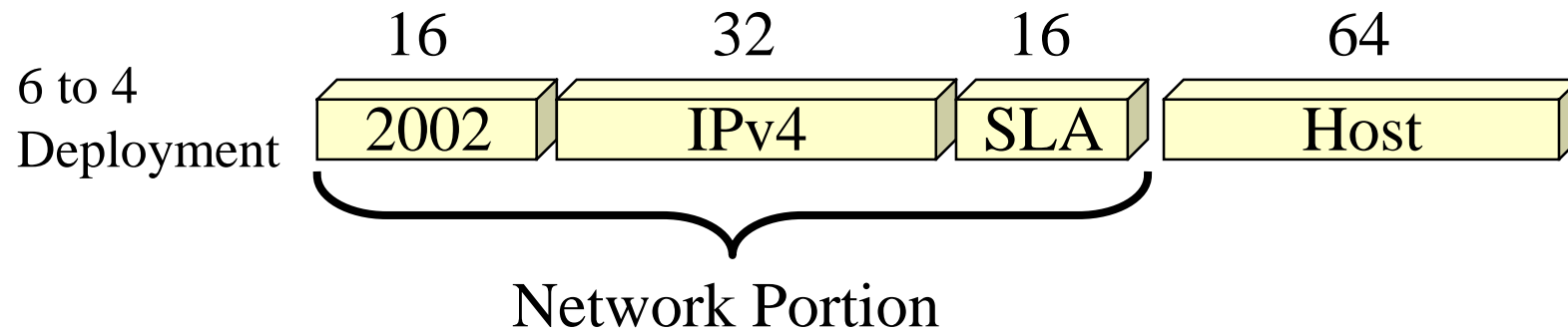


Transition Mechanisms

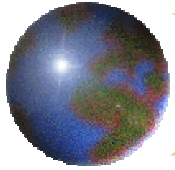
- Tunnel IPv6 over IPv4 (see p78)
- Static/Configured tunnels
 - Especially by ISPs and 6bone
 - Tunnel broker service
- 6over4
 - Use an IPv4 multicast enabled cloud as a single IPv6 link
- 6to4



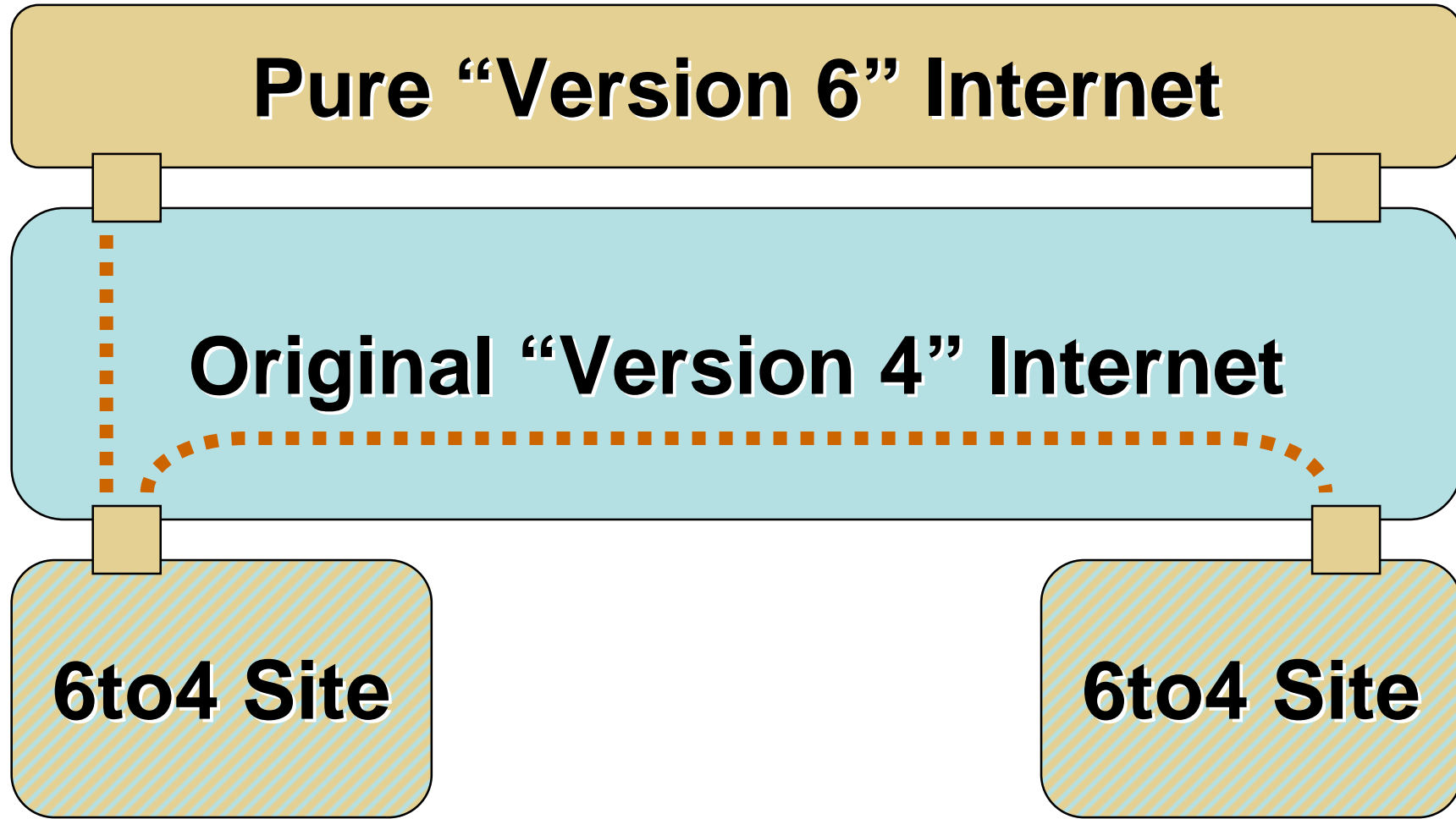
6to4 Address allocation

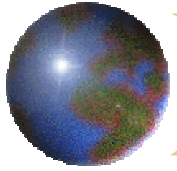


- ❁ Stateless tunnel over the IPv4 network
 - ❁ Assume IPv4 address of “lower layer”
 - ❁ Entire campus fits behind one IPv4 address
 - ❁ No dependence on IPv6 core



A picture of 6to4





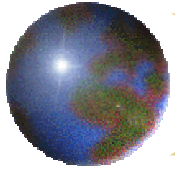
Mobile IP

☉ Each mobile host:

- ☐ Has a unique permanent address
- ☐ Adopts a “home agent” on home network
- ☐ Sends location updates to agent.
- ☐ Acquires “care-of address” on foreign networks

☉ Home agent:

- ☐ Intercepts packets intended for mobile host
- ☐ Forwards to mobile host’s care-of-address

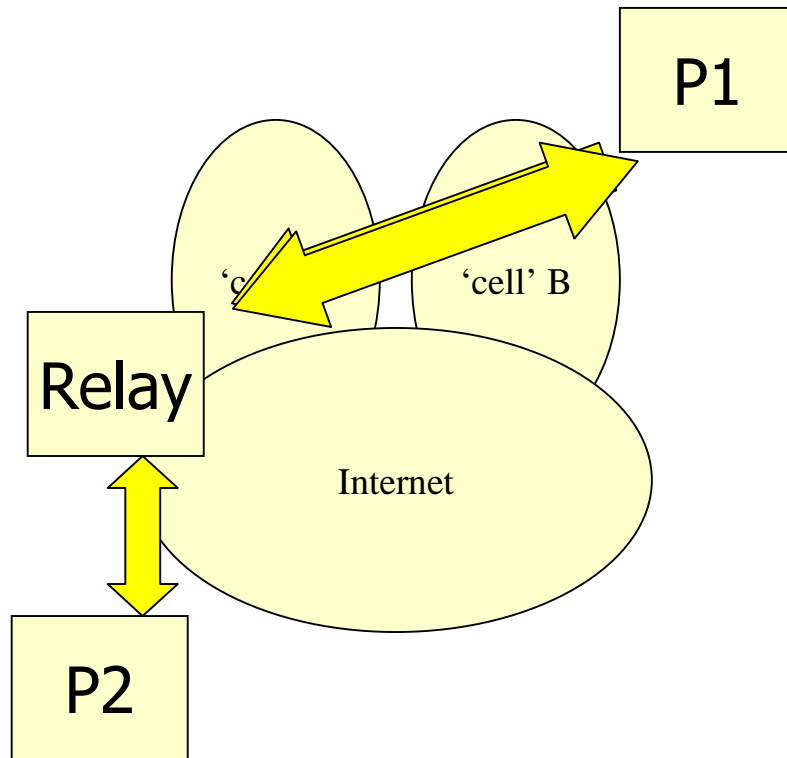


Mobile IPv6

- Less complex than Mobile IPv4
 - Requires ***NO EXTRA FUNCTIONALITY*** in off site routers
- Reuses many standard IPv6 services
- Improved support for low latency handoff
 - Stateless Address Configuration
- Fast convergence to optimal route
- Improved security/authentication mechanisms



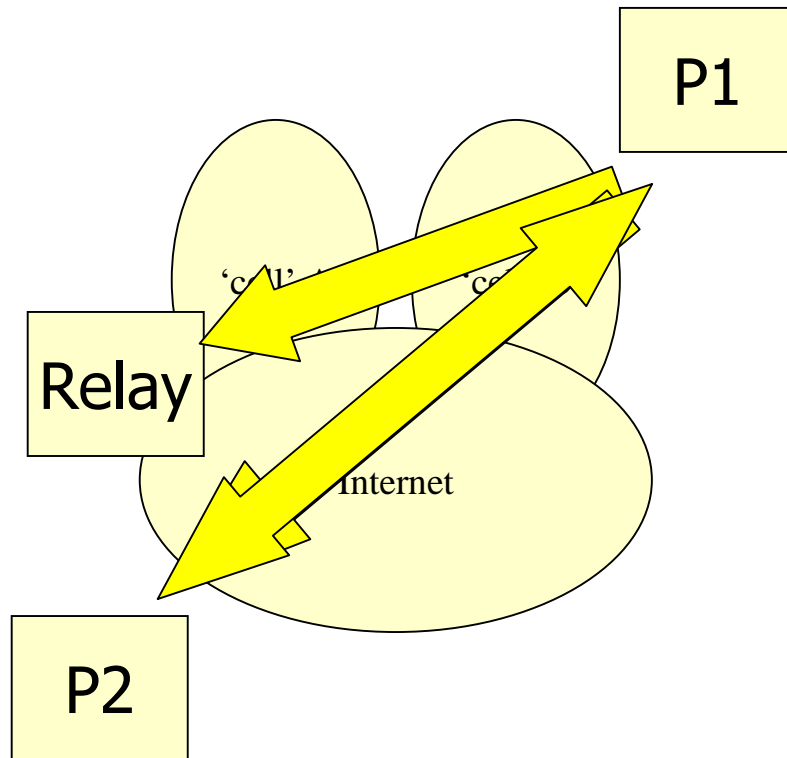
Problem 4: Move from “cell” to “cell”



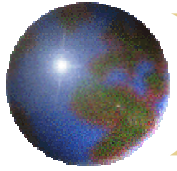
- IPv4:
 - Tell server,
 - Packets are relayed through the server



Solution 4: Move from “cell” to “cell” with IPv6

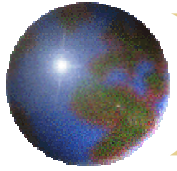


- IPv6:
 - Tell server + peer
 - Packets take direct path

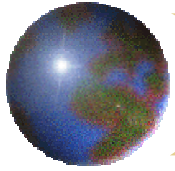


Where to Get More Information

- <http://www.microsoft.com/windows2000/library/technologies/communications/>
- <http://msdn.microsoft.com/downloads/sdk/s/platform/tpipv6.asp>



The End



Overview

- Give the big picture of the subject
- Explain how all the individual topics fit together

