

Important mathematical jargon: Sets

Very roughly, sets are the mathematicians' data structures. Informally, we will consider a set as a (well-defined, unordered) collection of mathematical objects, called the elements (or members) of the set.

Set membership

The symbol ' \in ' known as the *set membership* predicate is central to the theory of sets, and its purpose is to build statements of the form

$$x \in A$$

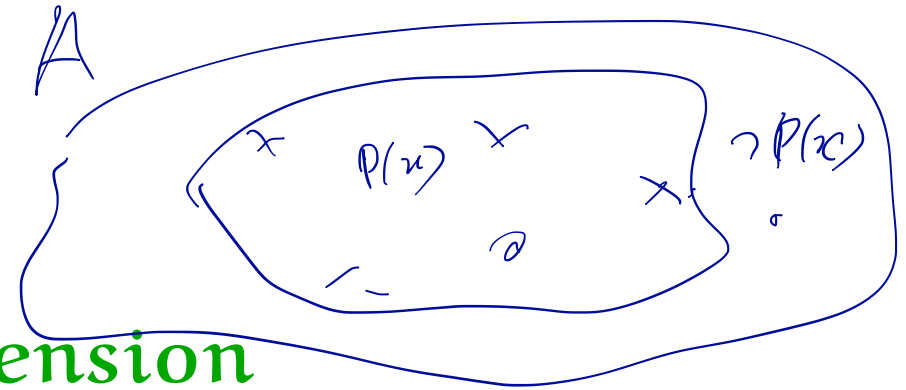
that are true whenever it is the case that the object x is an element of the set A , and false otherwise. *Equality of sets :*

$$A = B \quad \text{iff} \quad \forall x. \quad x \in A \Leftrightarrow x \in B.$$

Defining sets

The set	of even primes	is	{2}
	of booleans		{true, false}
	$[-2..3]$		$\{-2, -1, 0, 1, 2, 3\}$

\mathbb{N} the set of natural numbers
 $\{0, 1, 2, \dots, n, \dots\}$.



Set comprehension

The basic idea behind set comprehension is to define a set by means of a property that precisely characterises all the elements of the set.

Notations:

$$\{x \in A \mid P(x)\} \quad , \quad \{x \in A : P(x)\}$$

Greatest common divisor

Given a natural number n , the set of its *divisors* is defined by set comprehension as follows

$$D(n) = \{ d \in \mathbb{N} : d \mid n \} .$$

Example 53

1. $D(0) = \mathbb{N}$

2. $D(1224) = \left\{ \begin{array}{l} 1, 2, 3, 4, 6, 8, 9, 12, 17, 18, 24, 34, 36, 51, 68, \\ 72, 102, 136, 153, 204, 306, 408, 612, 1224 \end{array} \right\}$

Remark Sets of divisors are hard to compute. However, the computation of the greatest divisor is straightforward. :)

Going a step further, what about the *common divisors* of pairs of natural numbers? That is, the set

$$\text{CD}(m, n) = \{ d \in \mathbb{N} : d \mid m \wedge d \mid n \}$$

for $m, n \in \mathbb{N}$.

$$\text{CD}(n, m)$$

Example 54

$$\text{CD}(1224, 660) = \{ 1, 2, 3, 4, 6, 12 \}$$

Since $\text{CD}(n, n) = D(n)$, the computation of common divisors is as hard as that of divisors. But, what about the computation of the *greatest common divisor*?

$$\Rightarrow \text{hcf}$$

Lemma 56 (Key Lemma) Let m and m' be natural numbers and let n be a positive integer such that $m \equiv m' \pmod{n}$. Then,

$$\text{CD}(m, n) = \text{CD}(m', n) .$$

PROOF: Assume $m \equiv m' \pmod{n}$, i.e. $m' = m + k \cdot n$
for some $k \in \mathbb{Z}$. RTP. $\forall d \in \mathbb{N}$. $d | m$ & $d | n \Leftrightarrow d | m' & d | n$.
Let $d \in \mathbb{N}$.

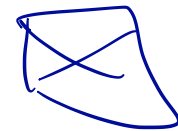
(\Rightarrow) Assume $d | m$ & $d | n$. Then

$d | m'$ because $d | (m + kn)$.

[Using $d | a$ & $d | b \Rightarrow d | (a + b)$].

So $d | m'$ & $d | n$.

(\Leftarrow) . Symmetrically



$$d \mid m \ \& \ n \mid m \ \Rightarrow \ d \mid n$$

Lemma 58 For all positive integers m and n ,

$$\text{by lemma 57} \quad \text{CD}(m, n) = \begin{cases} D(n) \checkmark & , \text{ if } n \mid m \\ \text{CD}(n, \text{rem}(m, n)) \checkmark & , \text{ otherwise} \end{cases}$$

$\text{CD}(\text{rem}(m, n), n)$

Lemma 58 For all positive integers m and n ,

$$\text{CD}(m, n) = \begin{cases} D(n) & , \text{ if } n \mid m \\ \text{CD}(n, \text{rem}(m, n)) & , \text{ otherwise} \end{cases}$$

Since a positive integer n is the greatest divisor in $D(n)$, the lemma suggests a recursive procedure:

$$\text{gcd}(m, n) = \begin{cases} n & , \text{ if } n \mid m \\ \text{gcd}(n, \text{rem}(m, n)) & , \text{ otherwise} \end{cases}$$

for computing the *greatest common divisor*, of two positive integers m and n . This is

Euclid's Algorithm

gcd

```
fun gcd( m , n )  
  = let  
    val ( q , r ) = divalg( m , n )  
  in  
    if r = 0 then n  
    else gcd( n , r )  
  end
```

Example 59 ($\gcd(13, 34) = 1$)

$$\begin{aligned}\gcd(13, 34) &= \gcd(34, 13) \\ &= \gcd(13, 8) \\ &= \gcd(8, 5) \\ &= \gcd(5, 3) \\ &= \gcd(3, 2) \\ &= \gcd(2, 1) \\ &= 1\end{aligned}$$

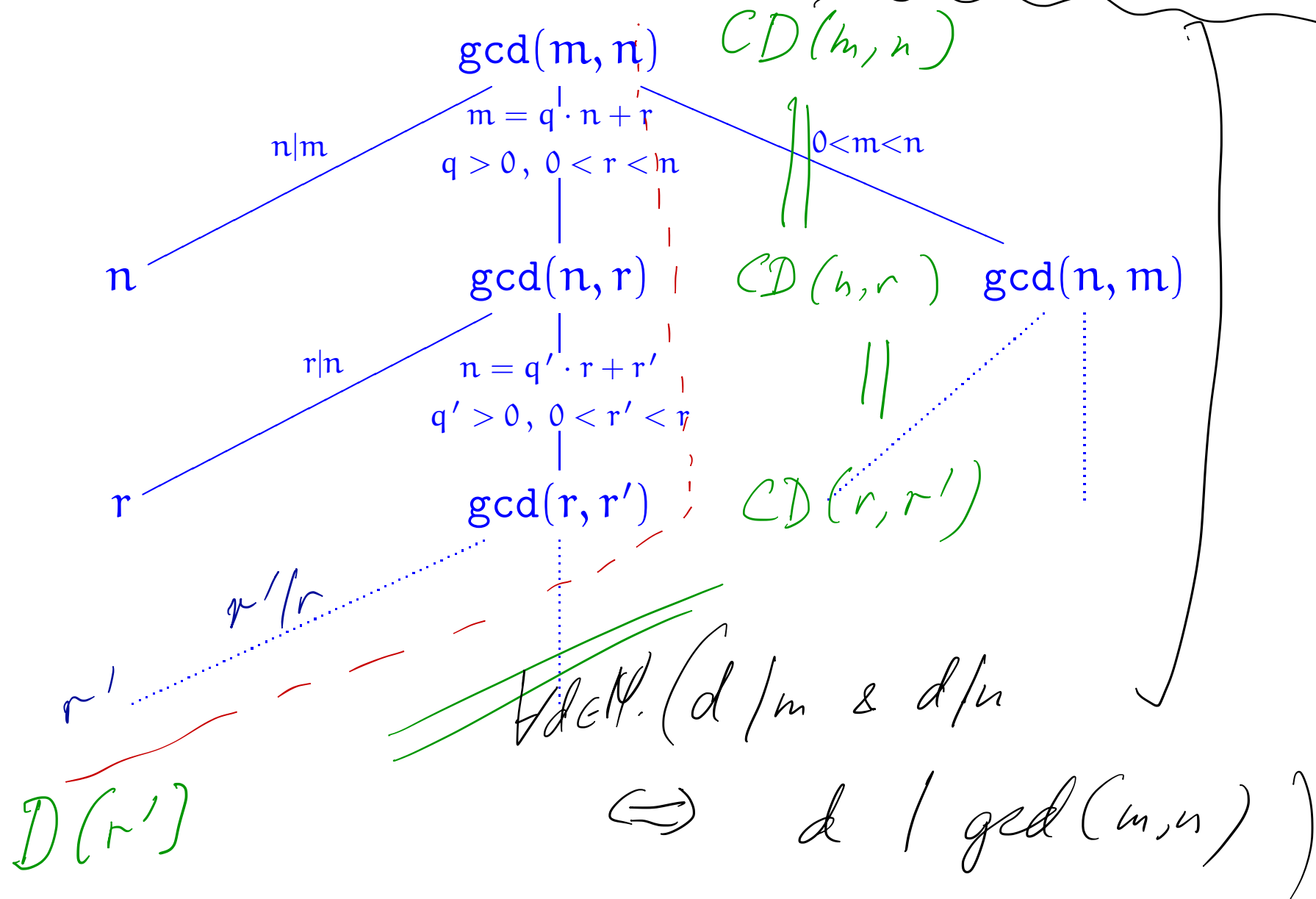
Theorem 60 *Euclid's Algorithm \gcd terminates on all pairs of positive integers and, for such m and n , $\gcd(m, n)$ is the greatest common divisor of m and n in the sense that the following two properties hold:*

- (i) both $\gcd(m, n) \mid m$ and $\gcd(m, n) \mid n$, and*
- (ii) for all positive integers d such that $d \mid m$ and $d \mid n$ it necessarily follows that $d \mid \gcd(m, n)$.*

PROOF:

Termination: it arg. decreases.

$$CD(m, n) \stackrel{?}{=} D(\gcd(m, n))$$



Fractions in lowest terms

```
fun lowterms( m , n )  
  = let  
    val gcdval = gcd( m , n )  
  in  
    ( m div gcdval , n div gcdval )  
  end
```

Some fundamental properties of gcds

Lemma 62 For all positive integers l , m , and n ,

1. **(Commutativity)** $\gcd(m, n) = \gcd(n, m)$,
2. **(Associativity)** $\gcd(l, \gcd(m, n)) = \gcd(\gcd(l, m), n)$,
3. **(Linearity)^a** $\gcd(l \cdot m, l \cdot n) = l \cdot \gcd(m, n)$.

PROOF: \uparrow (3) (1) $l \cdot \gcd(m, n) \mid \gcd(l \cdot m, l \cdot n)$
(2) $\gcd(l \cdot m, l \cdot n) \mid l \cdot \gcd(m, n)$
(1) Have $\gcd(m, n) \mid m, n \therefore l \cdot \gcd(m, n) \mid l \cdot m, l \cdot n$
 $\therefore l \cdot \gcd(m, n) \mid \gcd(l \cdot m, l \cdot n)$

^aAka (Distributivity).

$$\text{RTP. (2) } \text{gcd}(l.m, l.n) \mid l.\text{gcd}(m, n)$$

$$\text{Note } l \mid \text{gcd}(l.m, l.n).$$

[Because $l \mid l.m, l.n$].

$$\therefore l.k = \text{gcd}(l.m, l.n) \quad (1)$$

for some $k \in \mathbb{N}$. Because $g \mid l$,

$$l.k \mid l.m, l.n$$

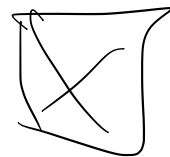
$$\therefore k \mid m, n.$$

$$\therefore k \mid \text{gcd}(m, n)$$

$$\therefore l.k \mid l.\text{gcd}(m, n)$$

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$$\text{gcd}(l.m, l.n)$$



Euclid's Theorem

Theorem 63 For positive integers k , m , and n , if $k \mid (m \cdot n)$ and $\gcd(k, m) = 1$ then $k \mid n$.

PROOF: