



Social and Technological Network Data Analytics

Lecture 6: Network Robustness and Applications

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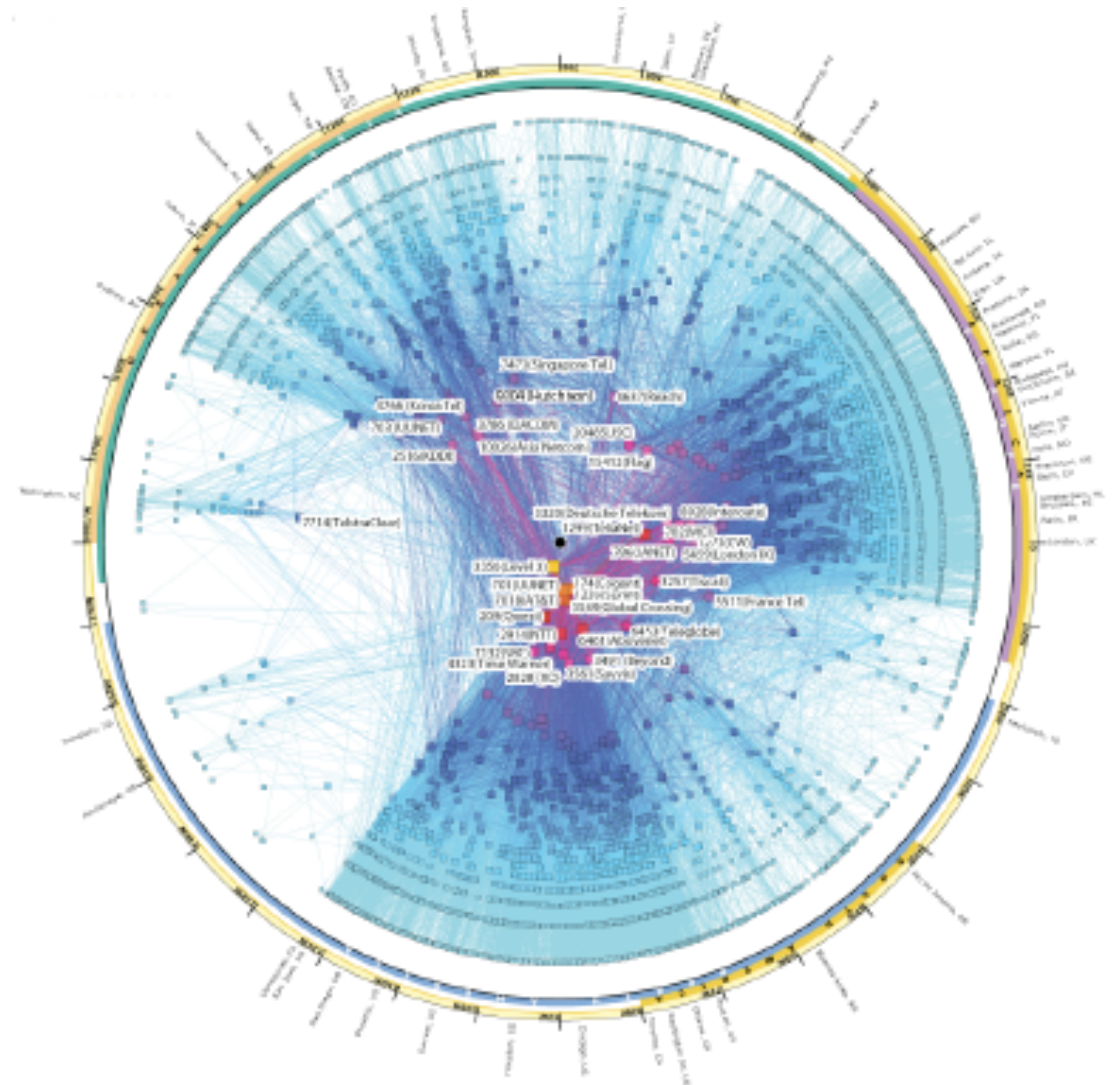
In This Lecture

- We revisit power-law networks and define the concept of robustness
- We show the effect of random and targeted attacks on power law networks versus random networks
- We discuss applications to various networks

Internet AS topology



- Autonomous System (AS): a collection of networks under the same administration
- 2009: 25,000 ASs in the Internet

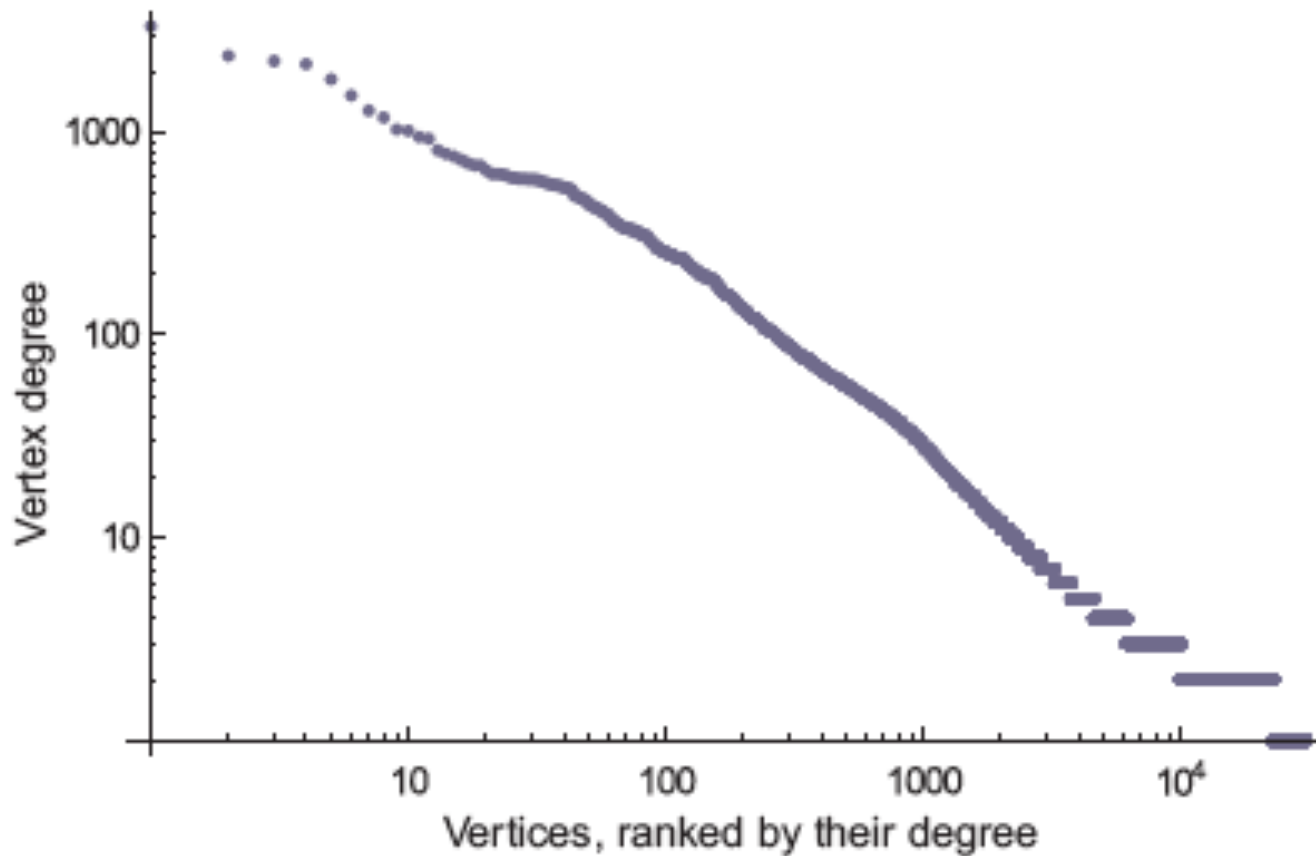


Topology Information



- By reading the routing tables of some gateways connected ASs, Internet topology information could be gathered
- October 08:
 - Over 30,000 ASs (including repeated entries)
 - Over 100,000 edges

Degree distribution of ASs: A scale free network!





Properties

- The top AS is connected to almost 10% of all ASs
- This connectedness drops rapidly
- Very high clustering coefficient for top 1000 hubs: an almost complete graph
- Most paths no longer than 3-4 hops
- Most ASs separated by shortest paths of maximum length of 6

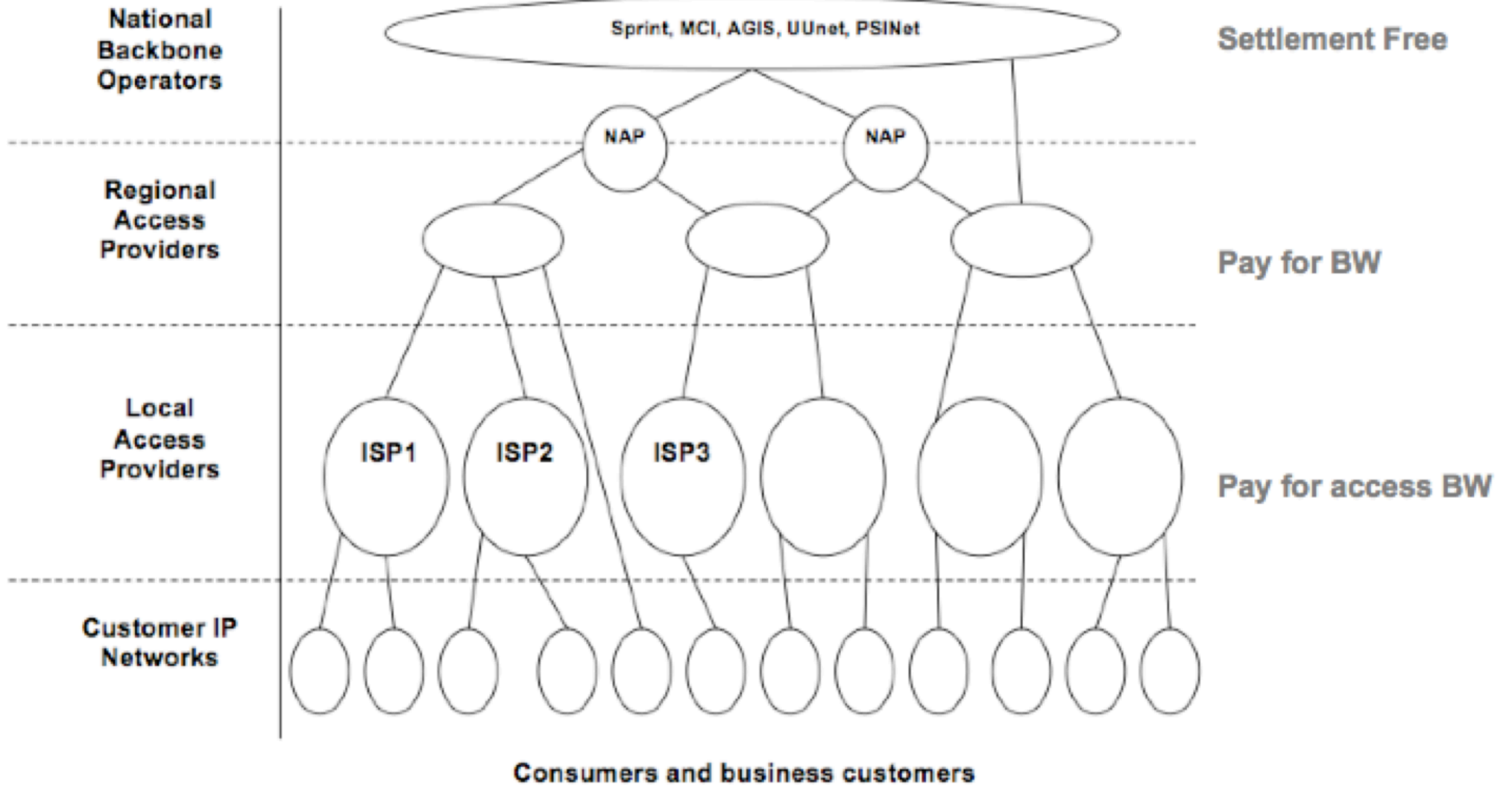
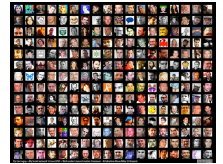
Rank:	1	2	3	4	5	6	7	8	9	10
Degree:	3309	2371	2232	2162	1816	1512	1273	1180	1029	1012

The Internet Now [Sigcomm10]

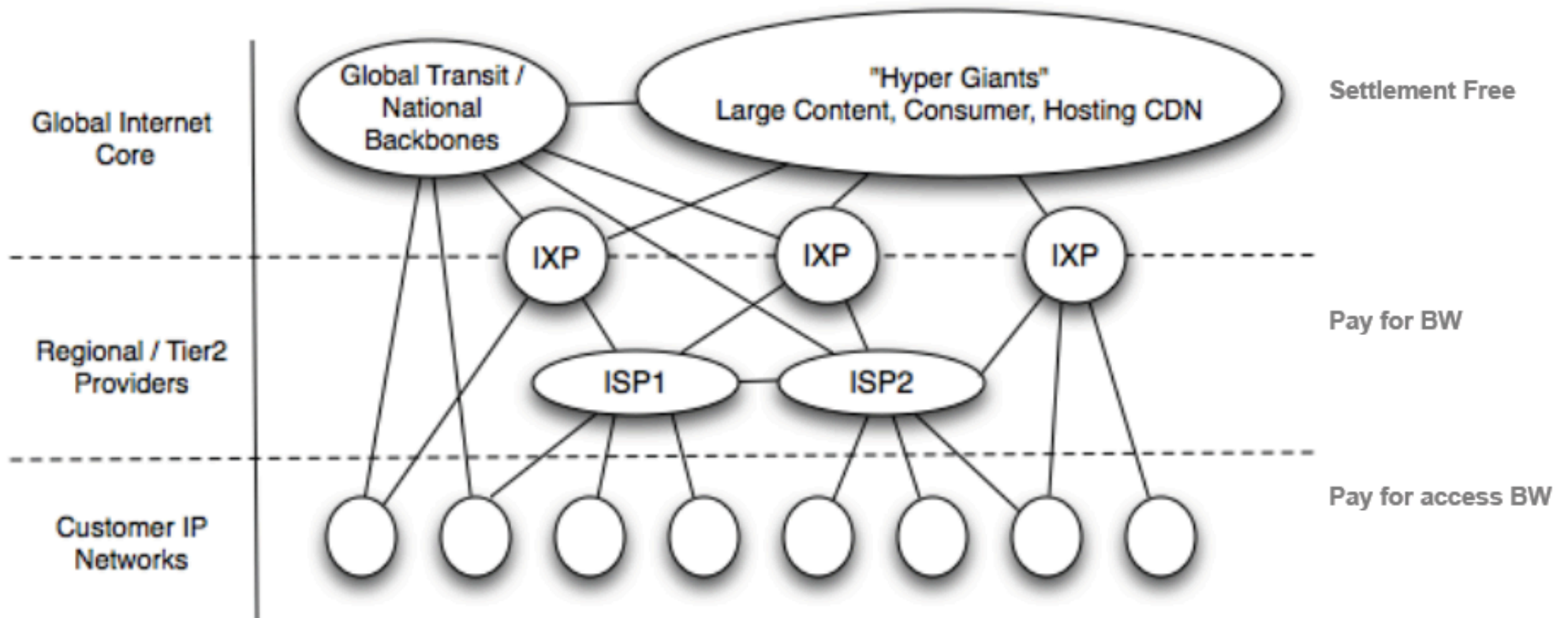


- They monitored inter-domain traffic for **2 years**
 - 3095 Routers
 - 110 ISPs
 - 18 Global
 - 38 Regional
 - 42 Consumer
 - 12 Terabits per second
 - 200 Exabytes total (200,000,000,000,000,000,000)
 - ~25% all inter-domain traffic.
- Inspect packets and classify them.

Internet 2007

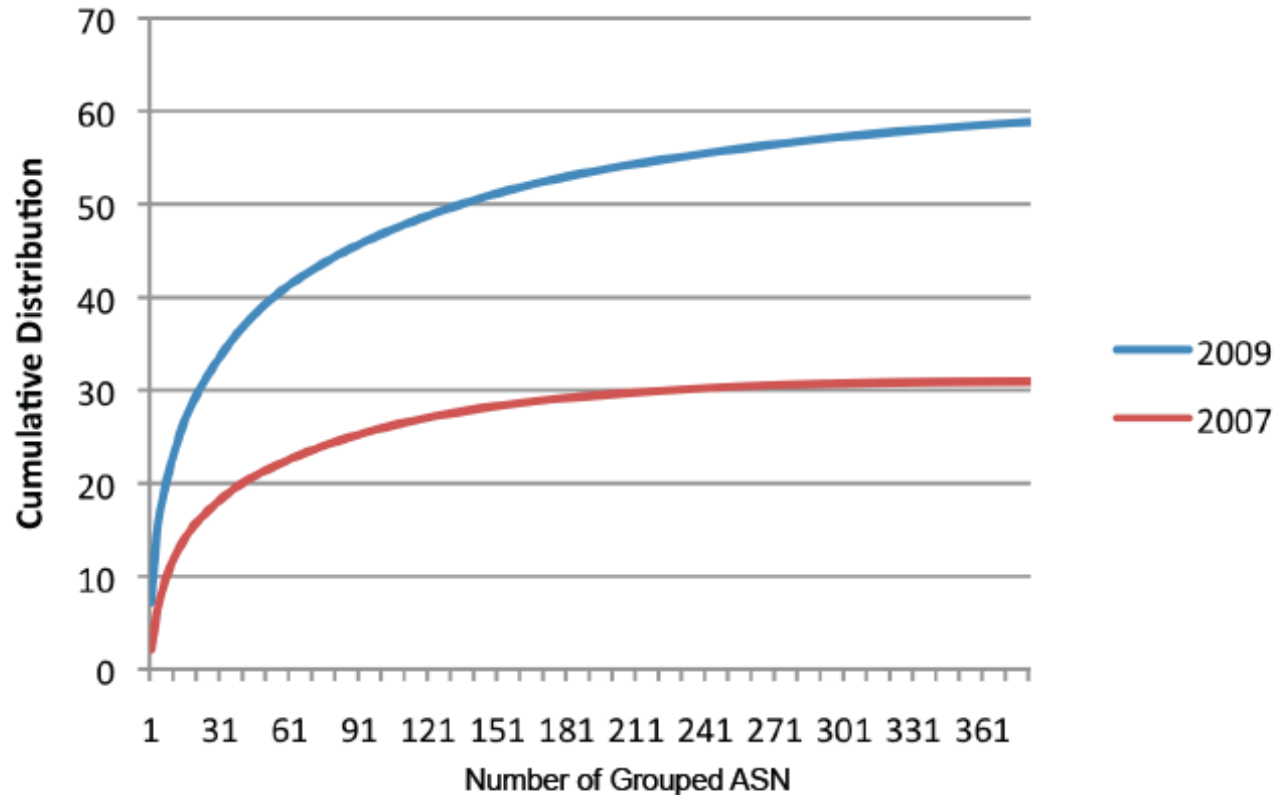


Internet 2009

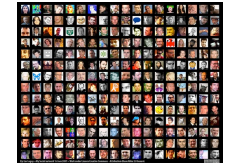


- Flatter and much more densely interconnected Internet
- Disintermediation between content and "eyeball" networks
- New commercial models between content, consumer and transit

Internet traffic: responsibility to few



- In 2007, thousands of ASNs contributed 50% of content
- In 2009, 150 ASNs contribute 50% of all Internet traffic



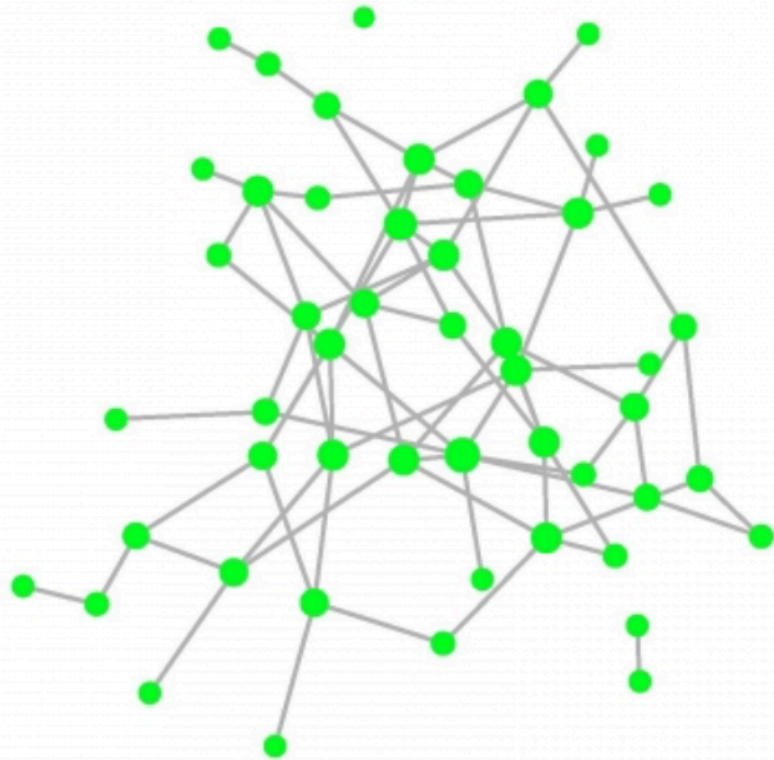
Robustness

- If a fraction of nodes or edges are removed:
 - How large are connected components?
 - What is the average distance between nodes in the components?
- This is related to *Percolation*
 - each edge/node is removed with probability $(1-p)$
 - Corresponds to random failure
 - Targeted attacks: remove nodes with high degree, or edges with high betweenness.
- The formation or dissolution of a giant component defines the percolation threshold

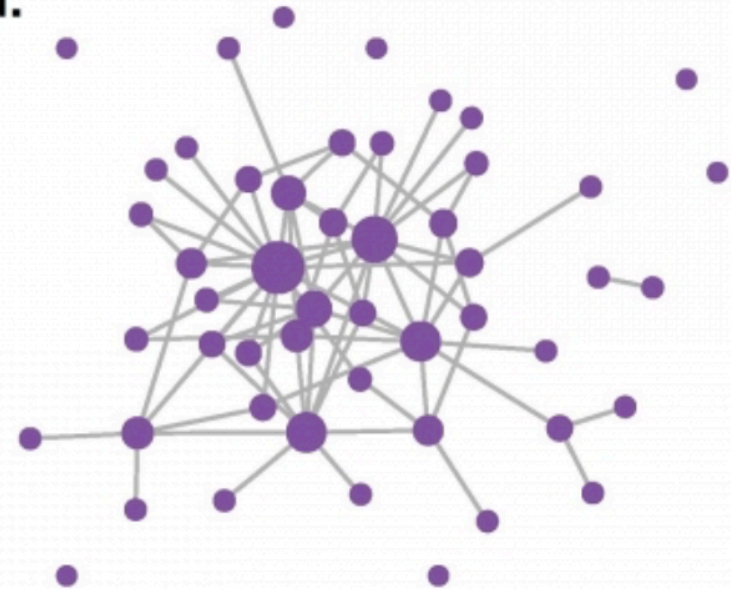
How Robust are These?



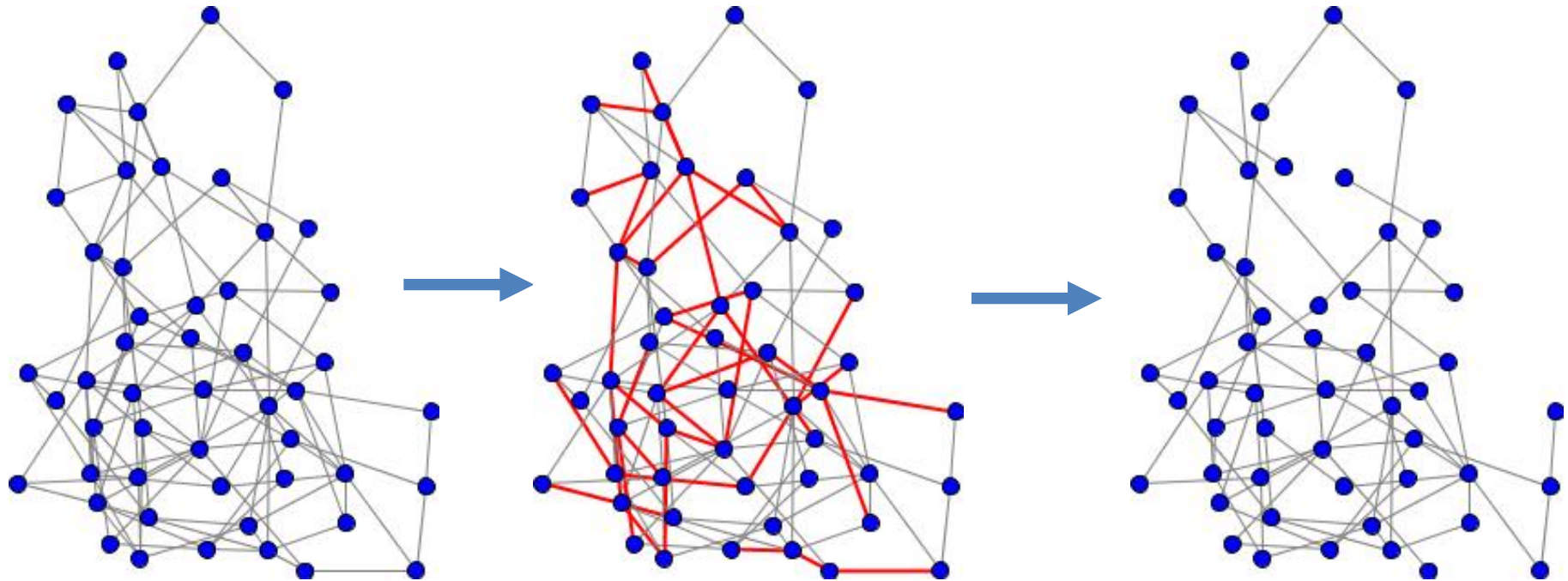
c.



d.

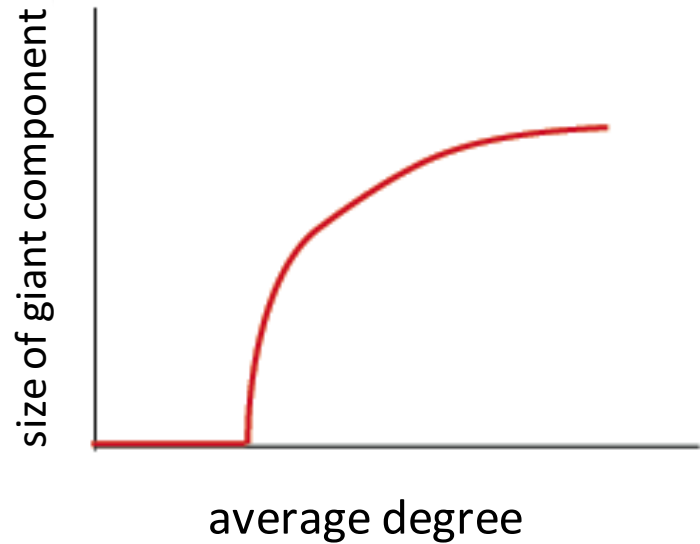


Edge (or Bond) Percolation



- 50 nodes, 116 edges, average degree 4.64
- after 25% edge removal
- 76 edges, average degree 3.04 – still well above percolation threshold

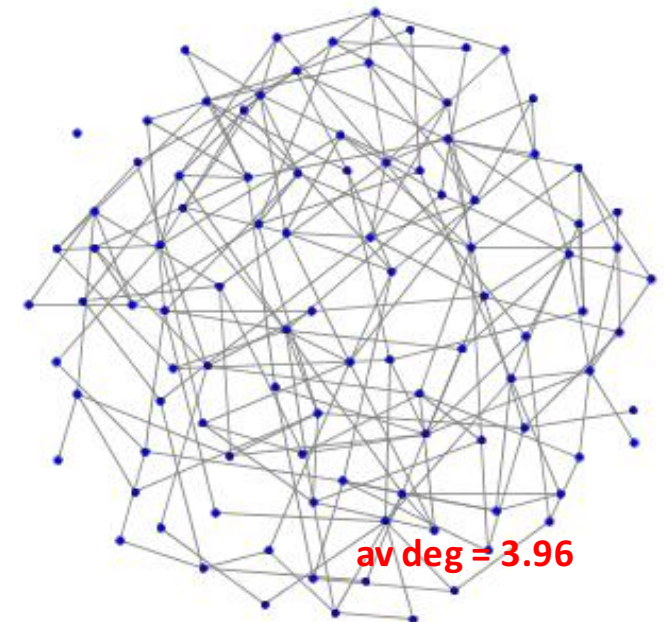
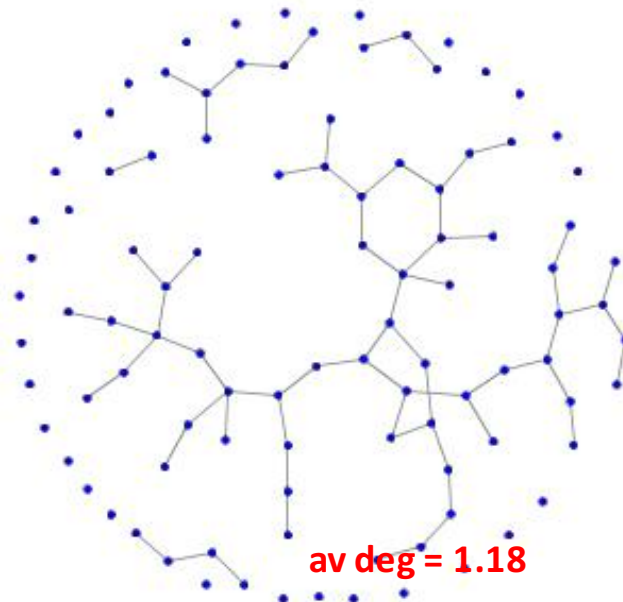
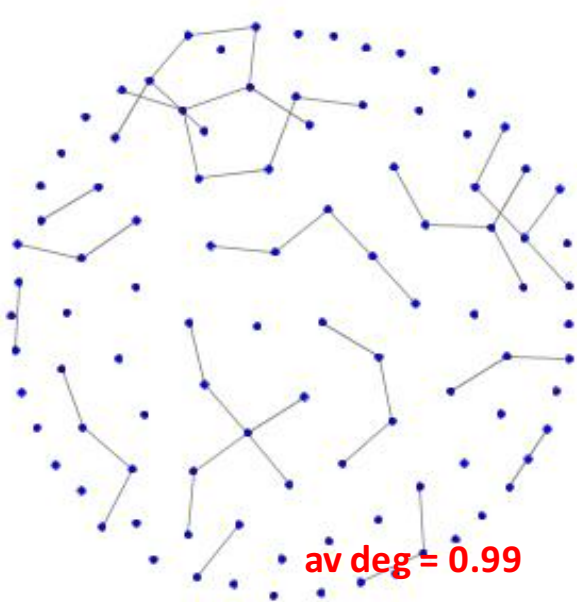
Percolation threshold in Random Graphs



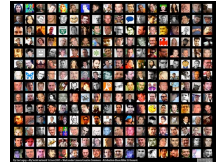
Percolation threshold: how many edges have to be removed before the giant component disappears?

As the average degree increases to 1, a giant component suddenly appears

Edge removal is the opposite process – at some point the average degree drops below 1 and the network becomes disconnected



Barabasi-Yeong-Albert's study (2000)

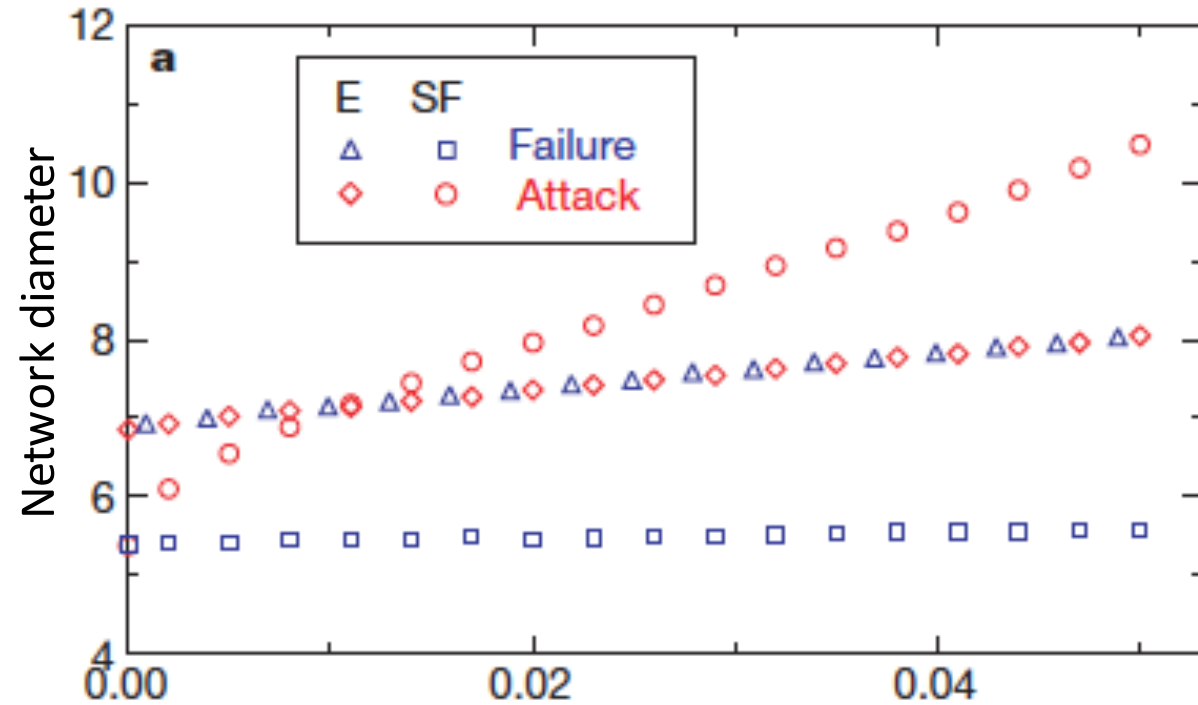


- Given 2 networks (one exponential one scale free) with same number of nodes and links
- Remove a small number of nodes and study changes in average shortest path to see if information communication has been disrupted and how much.

Let's look at the blue lines



- Random graph: increasing monotonically
- SF: remains unchanged until at least 5%

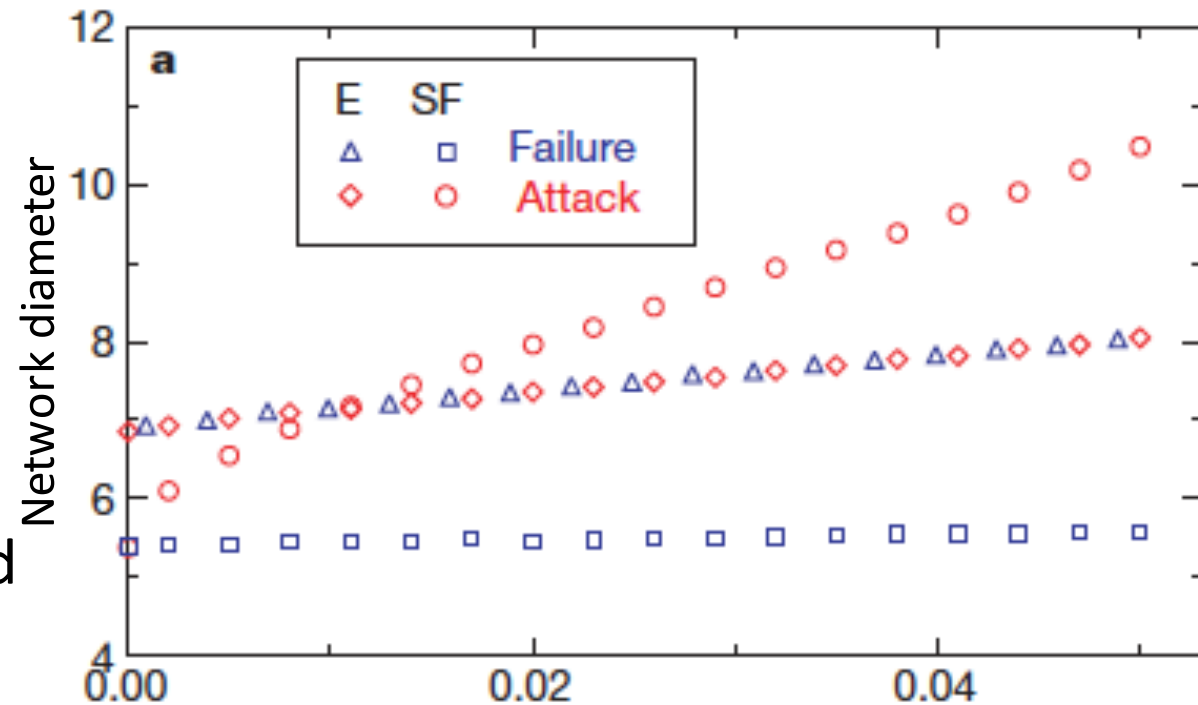


Fraction of deleted nodes

Let's look at the red lines

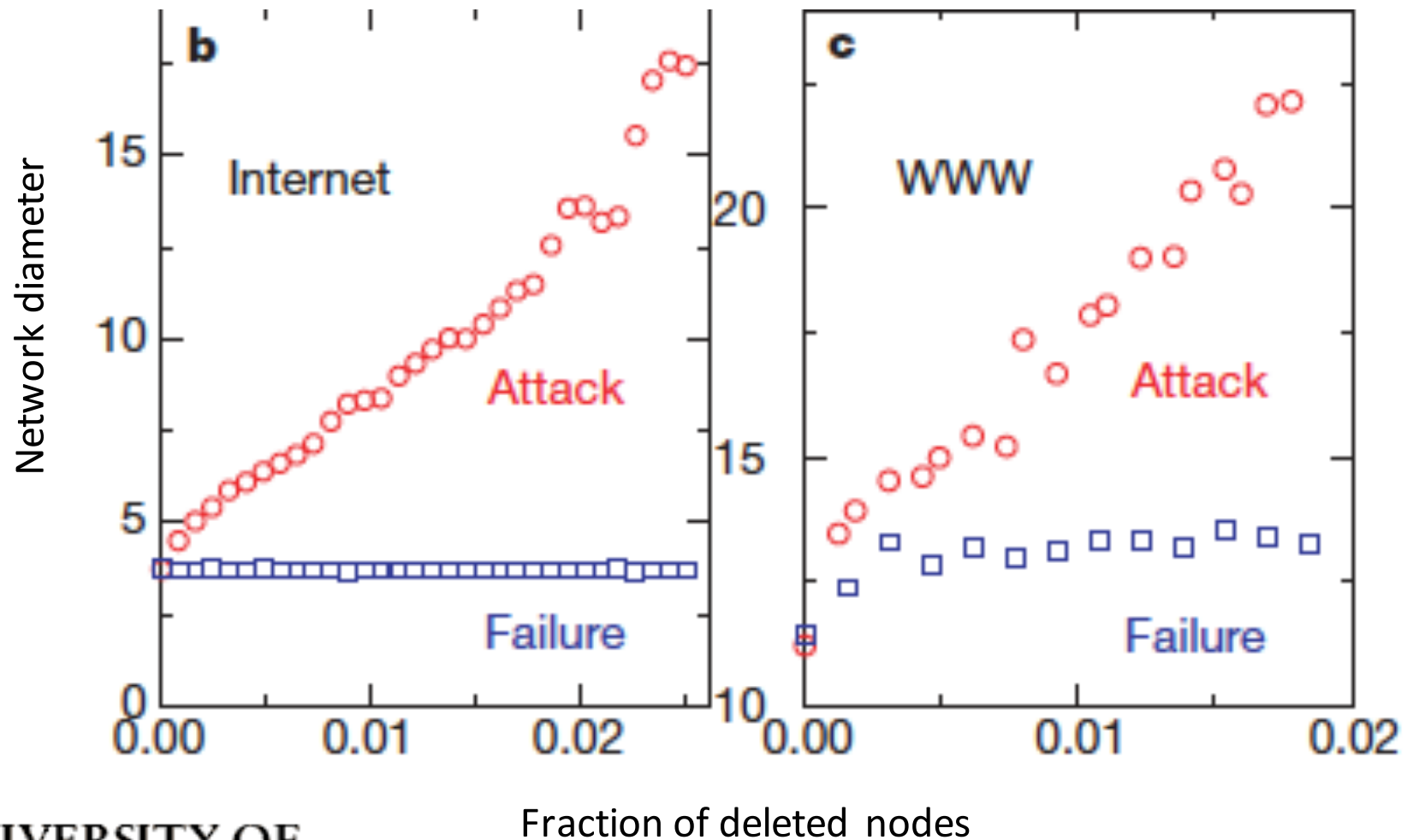
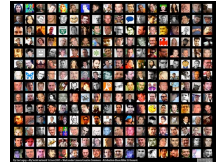


- Random graph: same behaviour if nodes with most links are chosen first
- SF: with 5% nodes removed the diameter is doubled

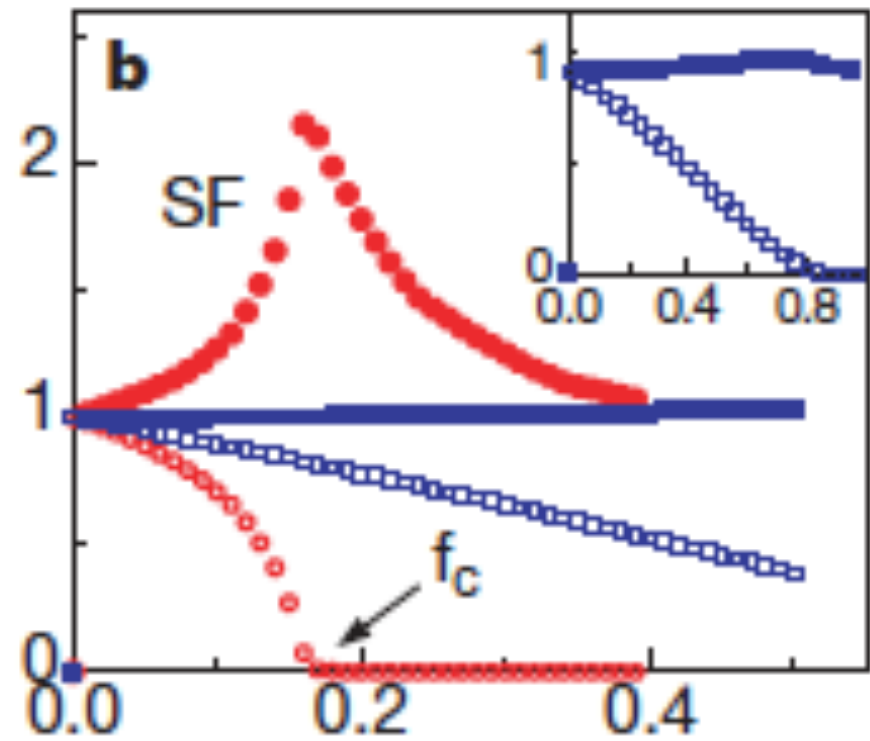
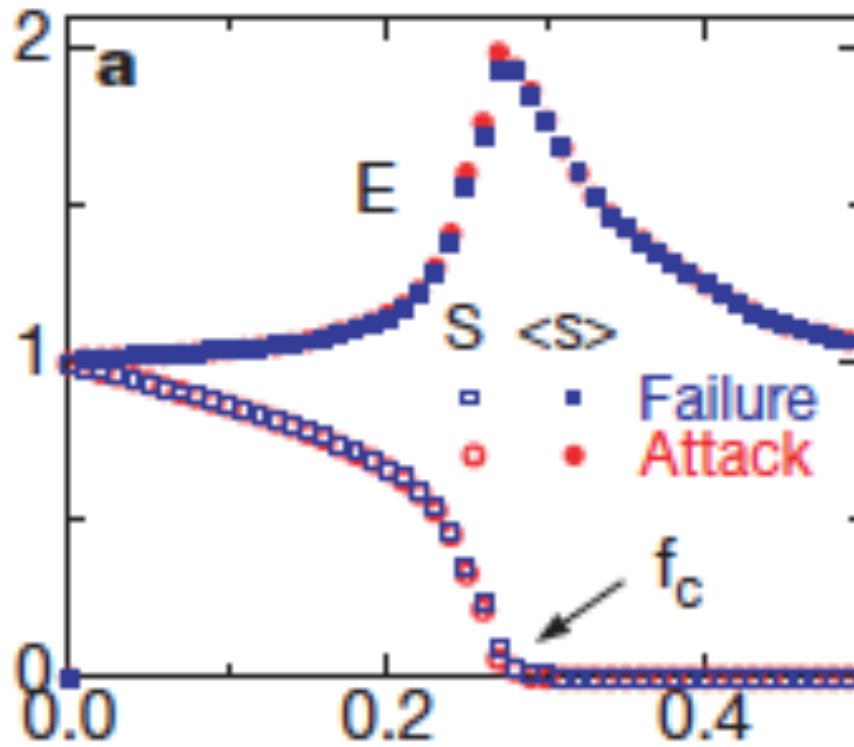


Fraction of deleted nodes

Effect of attacks and failure on WWW and Internet

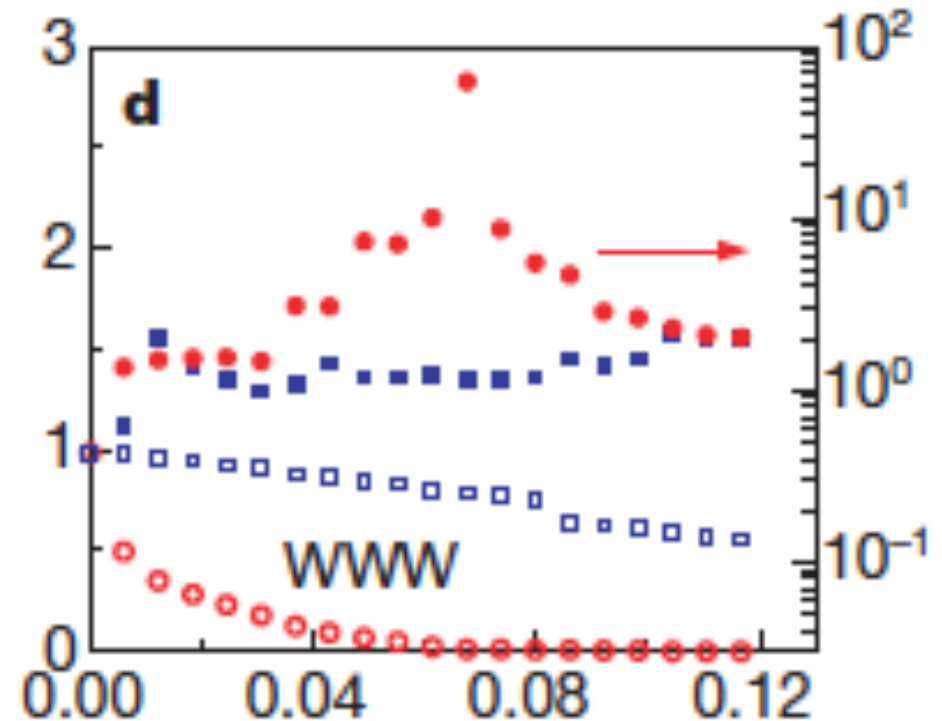
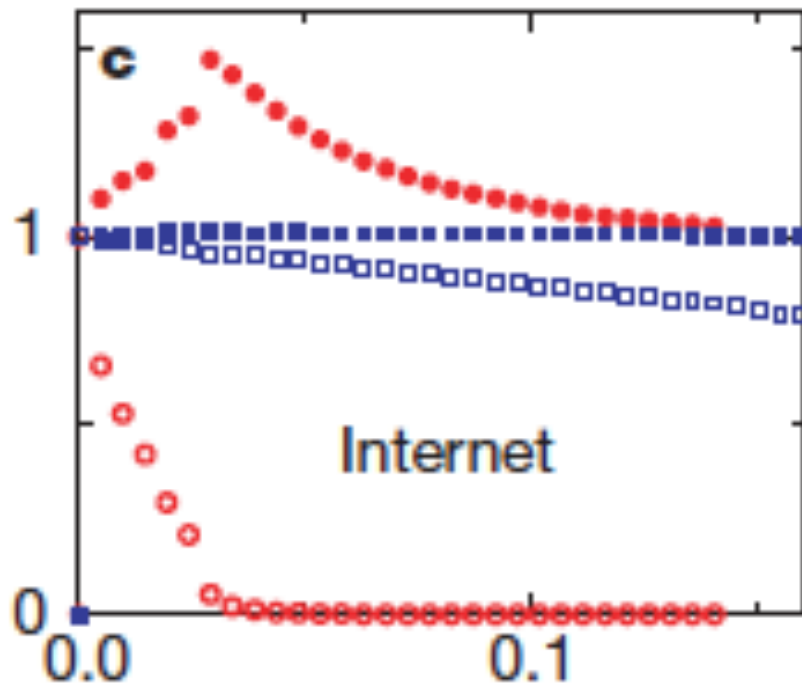
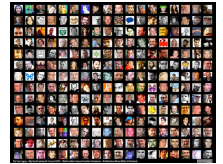


Effect on Giant Component



Fraction of deleted nodes

Internet and WWW: Effect on Giant Component

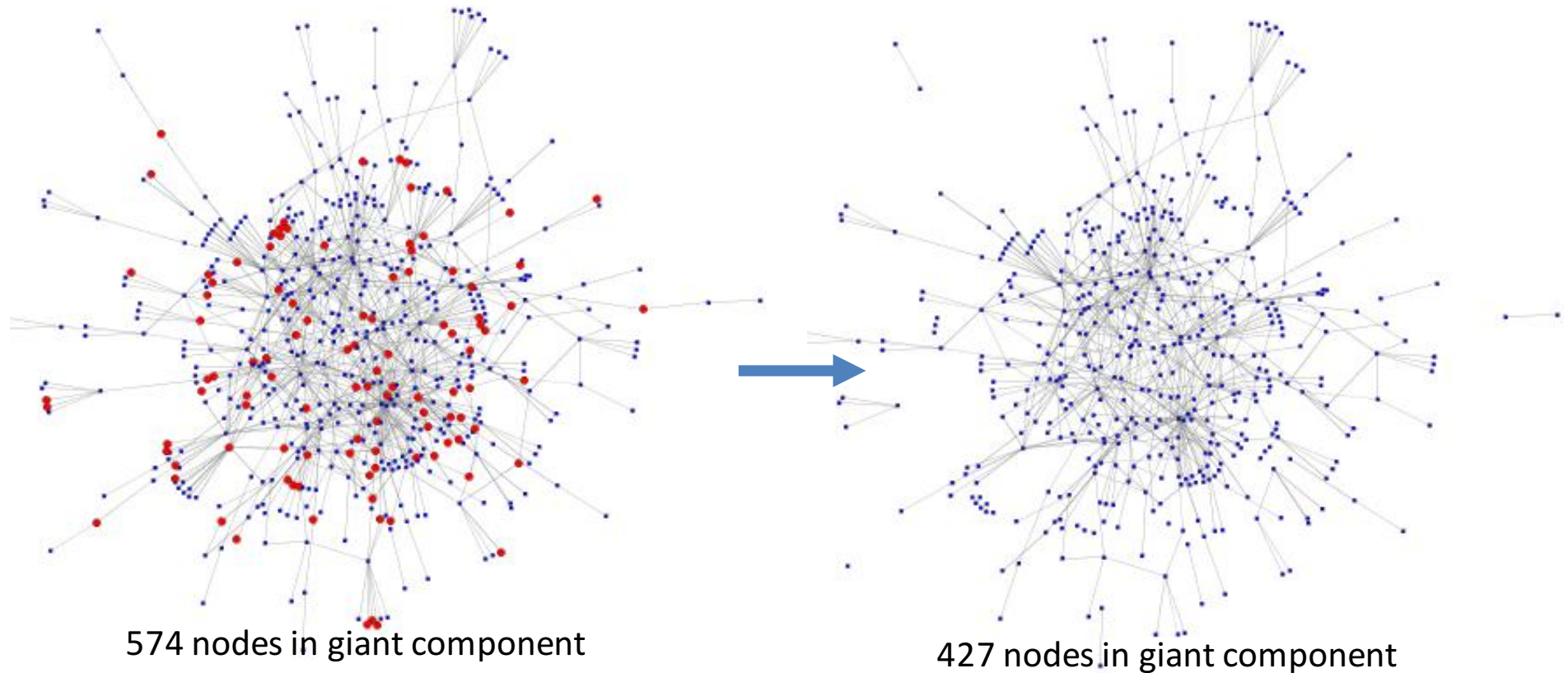


Fraction of deleted nodes

Scale-free networks are resilient with respect to random attack



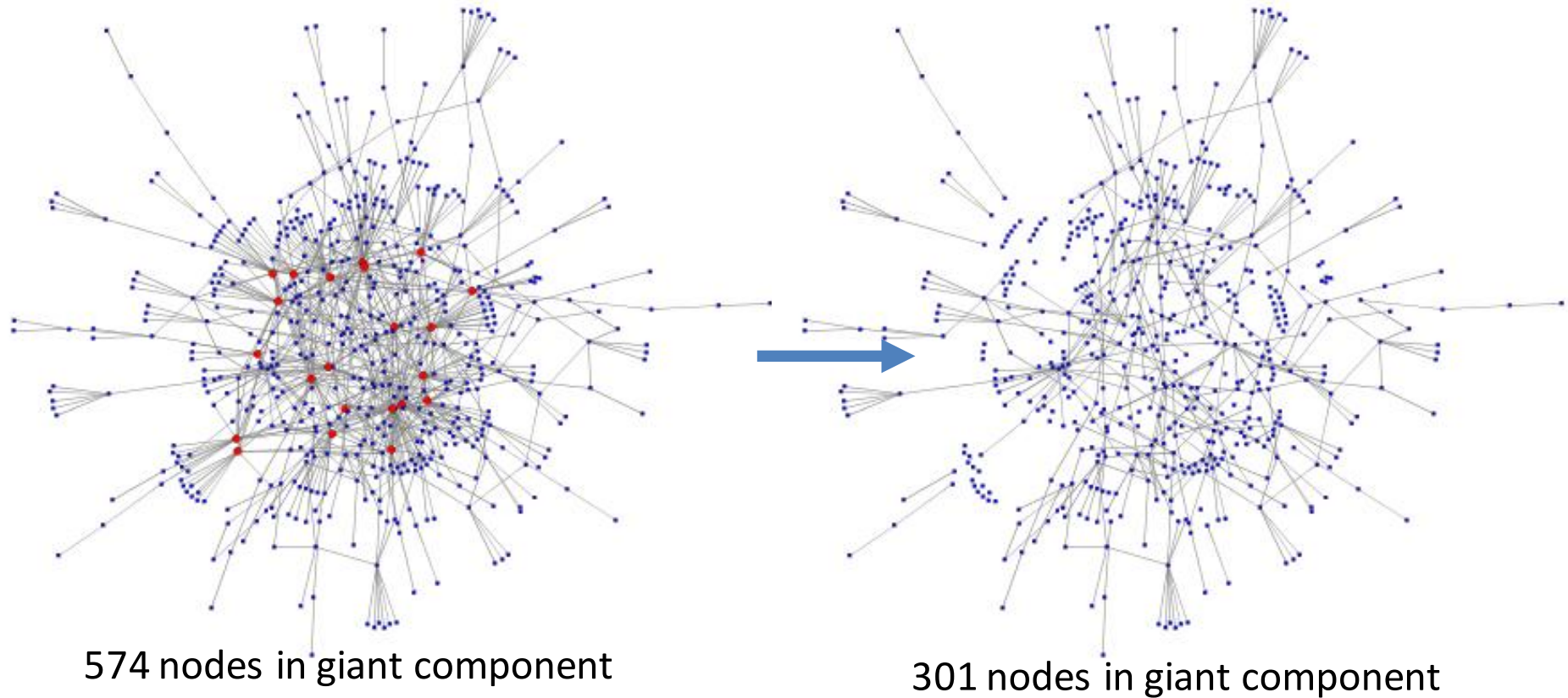
- Example: Gnutella network, 20% of nodes



Targeted attacks are effective against scale-free networks



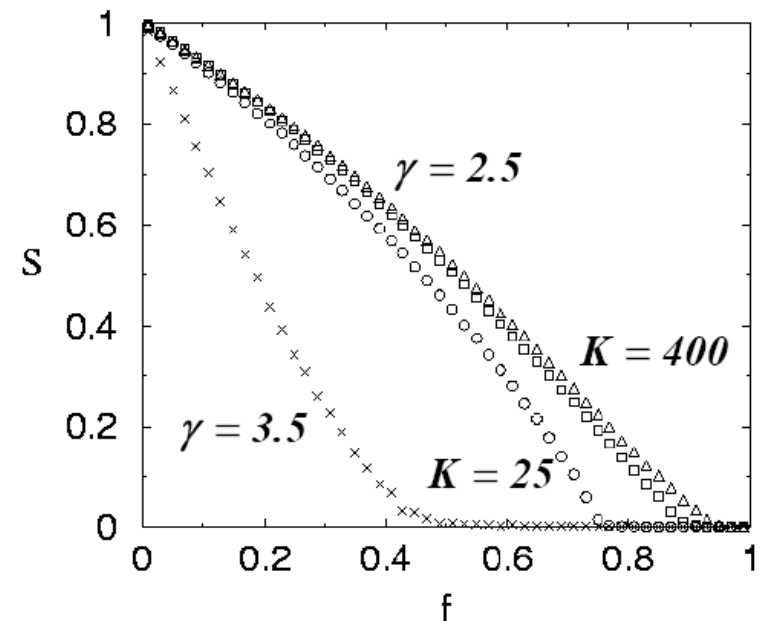
- Example: same Gnutella network, 22 most connected nodes removed (2.8% of the nodes)



Another study of power-laws



- Graph shows fraction of GC size over fraction of nodes randomly removed.
- Robustness of the Internet (γ is the exponent of PL).
 - $\gamma = 2.5$ Virtually no threshold exists which means a GC is always present
 - For $\gamma = 3.5$ there is a threshold around .0.4



Skewness of power-law networks and effects and targeted attack

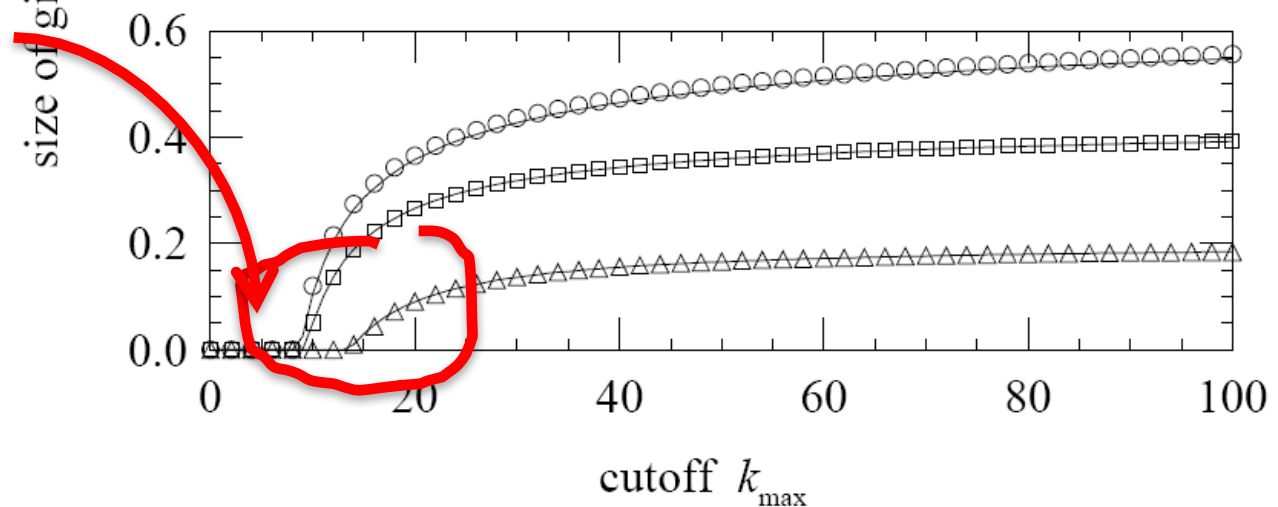
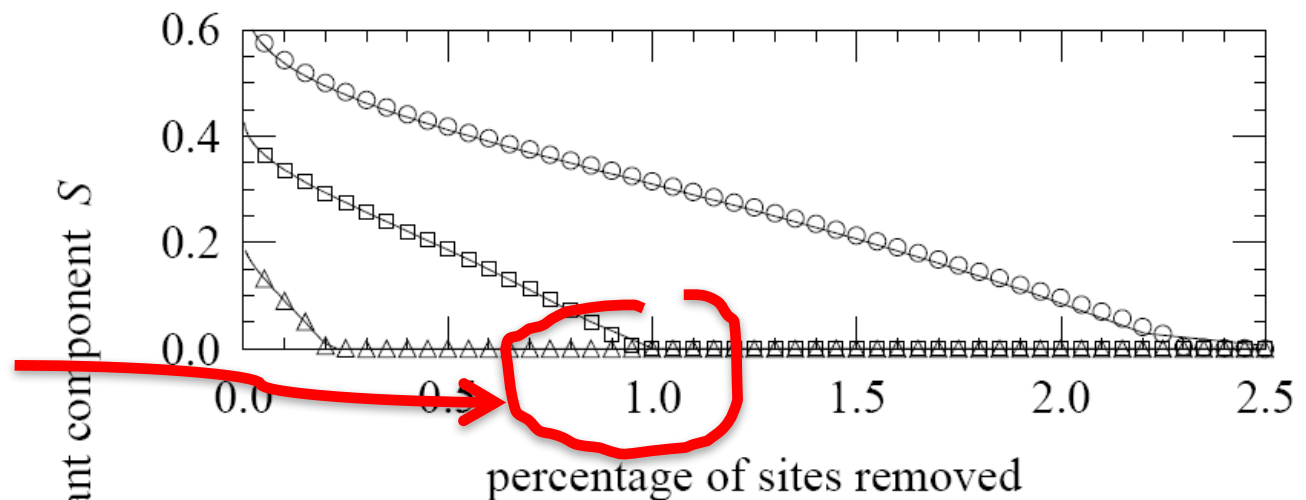


% of nodes removed,
from highest to lowest
degree

$\gamma = 2.7$ only 1% nodes
removed leads to no GC

k_{\max} needs to be very
low (10) to destroy the GC

k_{\max} is the highest
degree among the
remaining nodes



Percolation: let's get formal

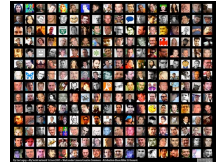


- Percolation process:
 - Occupation probability ϕ = number of nodes in the network [ie not removed]
- It can be proven that the critical threshold depends on the degree:

$$\phi_c = \frac{\langle k \rangle}{\langle k^2 \rangle - \langle k \rangle}$$

- This tells us the minimum fraction of nodes which must exist for a GC to exist.

Threshold for Random Graphs



- For Random networks $\varphi_{\text{critical}}=1/c$ where c is the mean degree
 - If c is large the network can withstand the loss of many vertices
 - $c=4$ then $\frac{1}{4}$ of vertices are enough to have a GC [3/4 of the vertices need to be destroyed to destroy the GC]

Threshold for Scale Free Networks



- For the Internet and Scale Free networks with $2 < \alpha < 3$
 - Finite mean $\langle k \rangle$ however $\langle k^2 \rangle$ diverges (in theory)
 - Then $\varphi_{\text{critical}}$ **is zero**: no matter how many vertices we remove there will always be a GC
 - In practice $\langle k^2 \rangle$ is never infinite for a finite network, although it can be very large, resulting in very small $\varphi_{\text{critical}}$, so still highly robust networks

Non random removal



- The threshold models we have presented hold for random node removal but not for targeted attacks [ie removal of high degree nodes first]
- The equation for non random removal cannot be solved analytically

Robustness Study and Improvements



- A method to improve network resilience
- Percolation threshold q ignores situation when the network is very damaged but not collapsing.

- Robustness:

$$R = \frac{1}{N} \sum_{Q=1}^N s(Q)$$

$S(Q)$ = fraction of nodes in the connected component after removing $Q=qN$ nodes

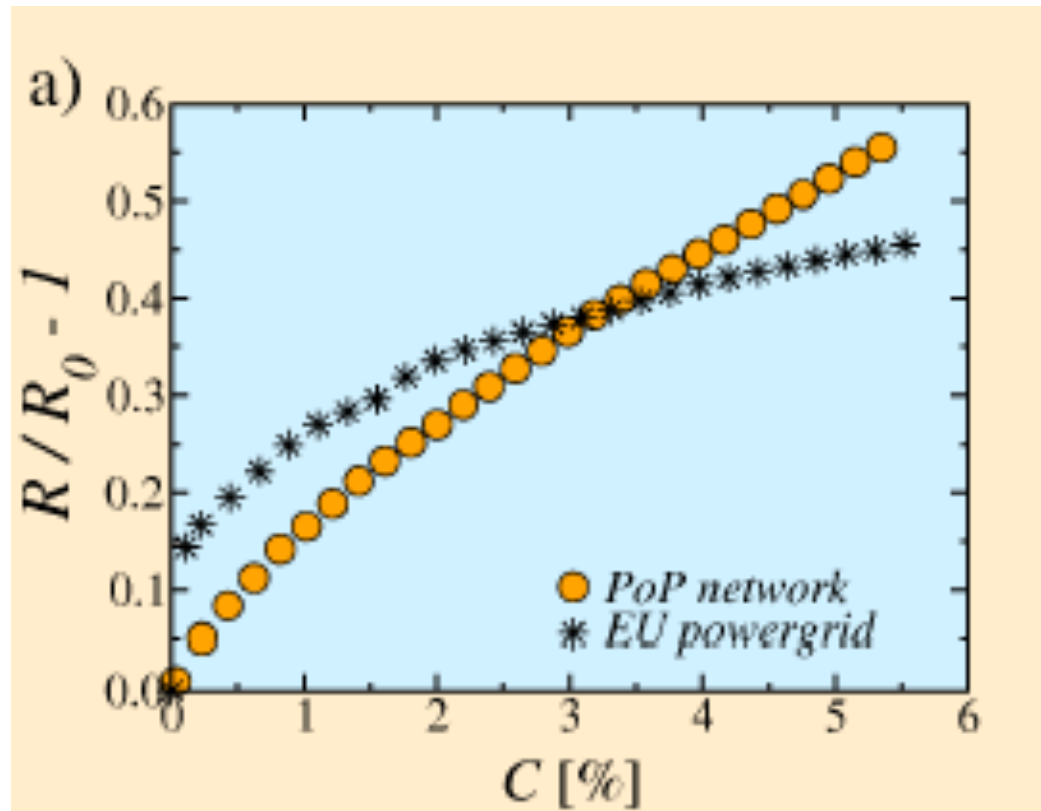
- R ranges values from star and fully connected graph.



Improve Robustness

- Add links until network is fully connected: not practical.
- Swap edges of 2 random nodes so that $R' > R$
 - Repeat until no substantial improvement (a value delta);
 - Some additional constraints could be introduced (limit the geographical length of new edges for economic reasons).

Robustness Improvement over edge changes

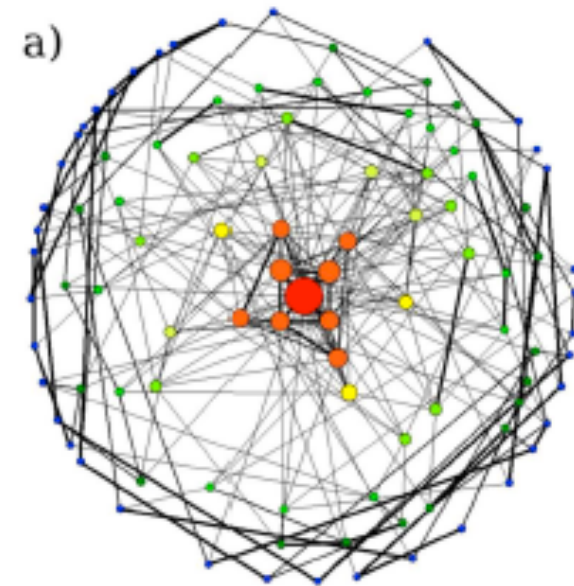


Robustness improved by 55-45% with ~5% link change. Percolation threshold remains unchanged.

Best Network for Robustness



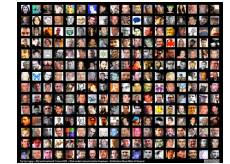
- How do we design a robust network with a fixed degree distribution?
- Scale free $N=100$ edges=300, exponent=2.5
- Onion-like structure!



Robustness of Technological and Social Network



- Targeted attacks on high degree nodes are lethal to a technological and a biological as well as transport network.
- However as seen in Lecture 2, for social systems it is the bridges and weak ties which make a difference...



References

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