

Binary products

In a category \mathcal{C} , a **product** for objects $X, Y \in \mathcal{C}$ is a diagram $X \xleftarrow{\pi_1} P \xrightarrow{\pi_2} Y$ with the universal property:

for all $X \xleftarrow{f} Z \xrightarrow{g} Y$
there is a unique $h: Z \rightarrow P$

such that

commutes

Binary products

In a category \mathcal{C} , a **product** for objects

$X, Y \in \mathcal{C}$ is a diagram $X \xleftarrow{\pi_1} P \xrightarrow{\pi_2} Y$

satisfying

(P, π_1, π_2) is terminal in the category with
- objects (Z, f, g) where $X \xleftarrow{f} Z \xrightarrow{g} Y$ in \mathcal{C}

- morphisms $h: (Z, f, g) \rightarrow (Z', f', g')$ are

$h \in \mathcal{C}(Z, Z')$ such that $f' \circ h = f$ & $g' \circ h = g$

- composition & identities as in \mathcal{C}

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so if they exist, products are unique up to (unique) isomorphism.

Binary products - notation

Usual notation for product of X & Y is

$$X \xleftarrow{\pi_1} X \times Y \xrightarrow{\pi_2} Y$$

and, given $X \xleftarrow{f} Z \xrightarrow{g} Y$, the
unique $h: Z \rightarrow X \times Y$ with $\begin{cases} \pi_1 \circ h = f \\ \pi_2 \circ h = g \end{cases}$

will be written

$$\langle f, g \rangle : Z \rightarrow X \times Y$$

Examples of products

In Set : $X \times Y = \{(x, y) \mid x \in X \wedge y \in Y\}$

$$\pi_1(x, y) = x$$

$$\pi_2(x, y) = y$$

because ...

Examples of products

$$\text{In Pre} : (P, \leq) \times (Q, \leq)$$

$$= (P \times Q, \leq) \quad \text{product in Set}$$

$$(p_1, q_1) \leq (p_2, q_2) \triangleq p_1 \leq p_2 \wedge q_1 \leq q_2$$

in P in Q

$$\pi_1(p, q) = p$$

$$\pi_2(p, q) = q$$

} are monotone functions

Examples of products

In Mon : $(M, \cdot, e) \times (N, \cdot, e)$

$$= (M \times N, \cdot, (e, e))$$

product
in Set

$$(m_1, n_1) \cdot (m_2, n_2) \stackrel{\Delta}{=} (m_1 \cdot m_2, n_1 \cdot n_2)$$

in M

in N

unit for this is (e, e)

$$\pi_1(m, n) = m$$

$$\pi_2(m, n) = n$$

} give monoid homomorphisms

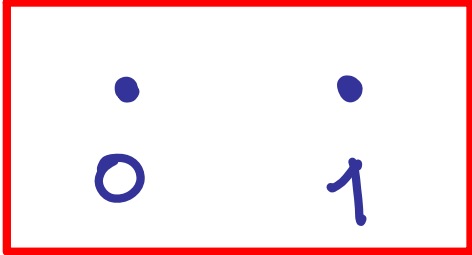
Examples of products

In a pre-ordered set (P, \leq) , regarded as a Category, the product of $p, q \in P$ is a greatest lower bound (glb, or meet) $p \wedge q$:

$$p \wedge q \leq p \quad \& \quad p \wedge q \leq q$$

$$(\forall r \in P) \quad r \leq p \quad \& \quad r \leq q \quad \Rightarrow \quad r \leq p \wedge q$$

Non-example

The poset , that is

$(\{0, 1\}, \leq)$ where $0 \leq 0$ & $1 \leq 1$

does not possess a product (= meet)
for 0 & 1 .

Duality

Binary product in \mathcal{C}^{op} is called
binary **Coproduct** in \mathcal{C} .

E.g. coproduct of $X, Y \in \text{Set}$ is

$$X \xrightarrow{i_1} X \amalg Y \xleftarrow{i_2} Y$$

$$i_1(x) \triangleq (x, 0)$$

$$i_2(y) \triangleq (y, 1)$$

$$\triangleq \{(x, 0) \mid x \in X\} \cup \{(y, 1) \mid y \in Y\}$$