

# Some fundamental properties of gcds

**Lemma 62** For all positive integers  $l$ ,  $m$ , and  $n$ ,

1. **(Commutativity)**  $\gcd(m, n) = \gcd(n, m)$ ,
2. **(Associativity)**  $\gcd(l, \gcd(m, n)) = \gcd(\gcd(l, m), n)$ ,
3. **(Linearity)<sup>a</sup>**  $\gcd(l \cdot m, l \cdot n) = l \cdot \gcd(m, n)$ .

PROOF:

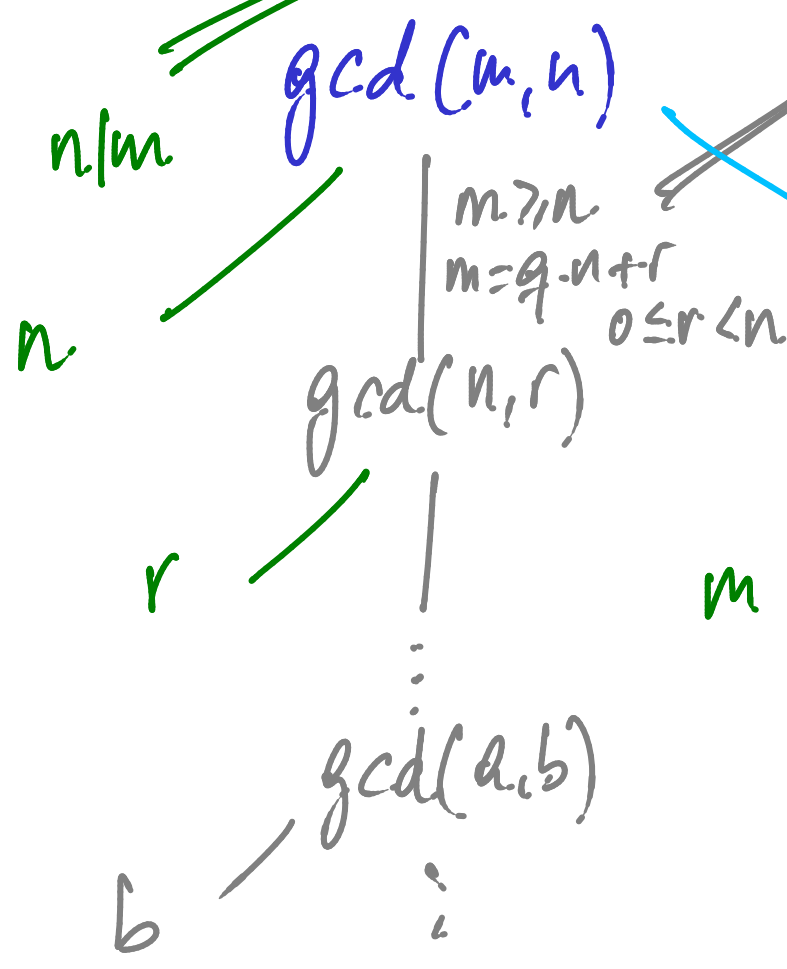
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<sup>a</sup>Aka (Distributivity).

Given  $m, n \rightarrow l, n, l, m$

$$l_m = q(l_n) + l_r$$

$$0 \leq l_r < l_n$$



output  
 $\gcd(m, n)$

$$l. \gcd(m, n) = \gcd(l_m, l_n)$$

# Euclid's Theorem

**Theorem 63** For positive integers  $k$ ,  $m$ , and  $n$ , if  $k \mid (m \cdot n)$  and  $\gcd(k, m) = 1$  then  $k \mid n$ .

PROOF: Consider integers  $k, m, n$ .

Assume: ①  $k \mid (m \cdot n)$  and ②  $\gcd(k, m) = 1$

From ②,  $n \cdot \gcd(k, m) = n$   
// by linearity

$\gcd(nk, mn)$

// by ①

$\gcd(nk, lk)$  for some int  $l$

//

by linearity

$k \cdot \gcd(n, l)$

$k \mid n \quad \square$

**Corollary 64 (Euclid's Theorem)** For positive integers  $m$  and  $n$ , and prime  $p$ , if  $p \mid (m \cdot n)$  then  $p \mid m$  or  $p \mid n$ .

Now, the second part of Fermat's Little Theorem follows as a corollary of the first part and Euclid's Theorem.

PROOF: Let  $m, n$  be pos. int. Let  $p$  be a prime

Assume  $p \mid (m \cdot n)$

RTP:  $p \mid m \vee p \mid n$

We argue by cases:

① If  $p \mid m$  then we are done.

② If  $\neg(p \mid m)$  then  $\gcd(p, m) = 1$  and by the previous then  $p \mid n$ .



Recall

for  $p$  prime

$$p \mid \binom{p}{m}$$

$$0 < m < p$$

$$\binom{p}{m} = p \frac{(p-1)!}{m!(p-m)!}$$

$$\Rightarrow [m!(p-m)!] \cdot \binom{p}{m} = p \cdot (p-1)!$$

$$\Rightarrow p \mid [m!(p-m)!] \cdot \binom{p}{m} \wedge p \nmid [m!(p-m)!]$$

$$\Rightarrow p \mid \binom{p}{m}$$



$$\text{FLT} : i^p \equiv i \pmod{p}$$

$$? \downarrow i^{p-1} \equiv 1 \pmod{p} \quad i \text{ not a multiple of } p \quad \Leftrightarrow \quad i \not\equiv 0 \pmod{p}$$

## Fields of modular arithmetic

**Corollary 66** For prime  $p$ , every non-zero element  $i$  of  $\mathbb{Z}_p$  has  $[i^{p-2}]_p$  as multiplicative inverse. Hence,  $\mathbb{Z}_p$  is what in the mathematical jargon is referred to as a field.

$$i^p \equiv i \pmod{p} \Rightarrow p \mid (i^p - i) = (i^{p-1} - 1) \cdot i$$

$$\text{If } p \nmid i \text{ then } p \mid i^{p-1} - 1 \Leftrightarrow i^{p-1} \equiv 1 \pmod{p}.$$

## Extended Euclid's Algorithm

### Example 67

		<i>quotients</i>		<i>remainders.</i>	
gcd(34, 13)	34 = 2 · 13 + 8				
= gcd(13, 8)	13 = 1 · 8 + 5				
= gcd(8, 5)	8 = 1 · 5 + 3				
= gcd(5, 3)	5 = 1 · 3 + 2				
= gcd(3, 2)	3 = 1 · 2 + 1				
= gcd(2, 1)	2 = 2 · 1 + 0				
= 1					

# Extended Euclid's Algorithm

## Example 67

$\gcd(34, 13)$	$34 = 2 \cdot 13 + 8$
$= \gcd(13, 8)$	$13 = 1 \cdot 8 + 5$
$= \gcd(8, 5)$	$8 = 1 \cdot 5 + 3$
$= \gcd(5, 3)$	$5 = 1 \cdot 3 + 2$
$= \gcd(3, 2)$	$3 = 1 \cdot 2 + 1$
$= \gcd(2, 1)$	$2 = 2 \cdot 1 + 0$
$= 1$	

*Reminders*

$8 = 34 - 2 \cdot 13$
$5 = 13 - 1 \cdot 8$
$3 = 8 - 1 \cdot 5$
$2 = 5 - 1 \cdot 3$
$1 = 3 - 1 \cdot 2$

*integer  
linear  
combinations*



$$\begin{array}{l}
\text{gcd}(34, 13) \\
= \text{gcd}(13, 8) \\
= \text{gcd}(8, 5) \\
= \text{gcd}(5, 3) \\
= \text{gcd}(3, 2)
\end{array}
\left| \begin{array}{l}
8 = 34 - 2 \cdot 13 \\
5 = 13 - 1 \cdot 8 \\
3 = 8 - 1 \cdot 5 \\
2 = 5 - 1 \cdot 3 \\
1 = 3 - 1 \cdot 2
\end{array} \right.$$

$$\text{gcd}(34, 13)$$

$$= \text{gcd}(13, 8)$$

$$= \text{gcd}(8, 5)$$

$$= \text{gcd}(5, 3)$$

$$= \text{gcd}(3, 2)$$

$$8 =$$

$$5 =$$

$$=$$

$$=$$

$$3 =$$

$$2 =$$

$$1 =$$

$$34$$

$$13$$

$$13$$

$$-1 \cdot 34 + 3 \cdot 13$$

$$8$$

$$5$$

$$3$$

$$-2 \cdot$$

$$-1 \cdot$$

$$-1 \cdot$$

$$-1 \cdot$$

$$-1 \cdot$$

$$-1 \cdot$$

$$13$$

$$8$$

$$(34 - 2 \cdot 13)$$

$$5$$

$$3$$

$$2$$

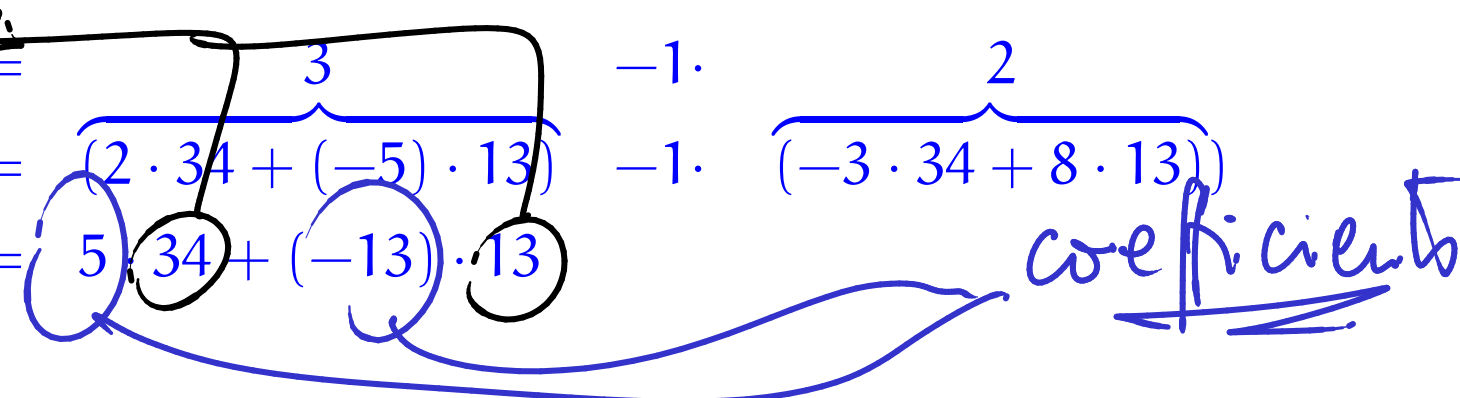


$\gcd(34, 13)$	$8 =$	$34$	$-2 \cdot$	$13$
$= \gcd(13, 8)$	$5 =$	$13$	$-1 \cdot$	$8$
	$=$	$13$	$-1 \cdot$	$\underbrace{(34 - 2 \cdot 13)}$
	$=$	$-1 \cdot 34 + 3 \cdot 13$		
$= \gcd(8, 5)$	$3 =$	$8$	$-1 \cdot$	$5$
	$=$	$\underbrace{(34 - 2 \cdot 13)}$	$-1 \cdot$	$\underbrace{(-1 \cdot 34 + 3 \cdot 13)}$
	$=$	$2 \cdot \textcircled{34} + (-5) \cdot \textcircled{13}$		
$= \gcd(5, 3)$	$2 =$	$5$	$-1 \cdot$	$3$
$= \gcd(3, 2)$	$1 =$	$3$	$-1 \cdot$	$2$

$\text{gcd}(34, 13)$	$8 =$	$34$	$-2 \cdot$	$13$
$= \text{gcd}(13, 8)$	$5 =$	$13$	$-1 \cdot$	$\underbrace{8}_{(34 - 2 \cdot 13)}$
	$=$	$13$	$-1 \cdot$	
	$=$	$-1 \cdot 34 + 3 \cdot 13$		
$= \text{gcd}(8, 5)$	$3 =$	$\underbrace{8}_{(34 - 2 \cdot 13)}$	$-1 \cdot$	$\underbrace{5}_{(-1 \cdot 34 + 3 \cdot 13)}$
	$=$	$\underbrace{2 \cdot 34 + (-5) \cdot 13}_6$		
$= \text{gcd}(5, 3)$	$2 =$	$\underbrace{-1 \cdot 34 + 3 \cdot 13}_5$	$-1 \cdot$	$\underbrace{3}_{(2 \cdot 34 + (-5) \cdot 13)}$
	$=$	$\underbrace{-3 \cdot 34 + 8 \cdot 13}_2$	$-1 \cdot$	
$= \text{gcd}(3, 2)$	$1 =$	$3$	$-1 \cdot$	$2$

Fact:  $\gcd(m, n)$  is an integer linear combination of  $m$  and  $n$ .

$\gcd(34, 13)$	$8 =$	$34$	$-2 \cdot$		$13$
$= \gcd(13, 8)$	$5 =$	$13$	$-1 \cdot$		$8$
	$=$	$13$	$-1 \cdot$		$\overbrace{(34 - 2 \cdot 13)}$
	$=$	$-1 \cdot 34 + 3 \cdot 13$			
$= \gcd(8, 5)$	$3 =$	$8$	$-1 \cdot$		$5$
	$=$	$\overbrace{(34 - 2 \cdot 13)}$	$-1 \cdot$		$\overbrace{(-1 \cdot 34 + 3 \cdot 13)}$
	$=$	$2 \cdot 34 + (-5) \cdot 13$			
$= \gcd(5, 3)$	$2 =$	$5$	$-1 \cdot$		$3$
	$=$	$\overbrace{-1 \cdot 34 + 3 \cdot 13}$	$-1 \cdot$		$\overbrace{(2 \cdot 34 + (-5) \cdot 13)}$
	$=$	$-3 \cdot 34 + 8 \cdot 13$			
$= \gcd(3, 2)$	$1 =$	$3$	$-1 \cdot$		$2$
	$=$	$\overbrace{(2 \cdot 34 + (-5) \cdot 13)}$	$-1 \cdot$		$\overbrace{(-3 \cdot 34 + 8 \cdot 13)}$
	$=$	$5 \cdot 34 + (-13) \cdot 13$			





Suppose  $\gcd(m, n) = 1$  [ $m$  and  $n$  are coprime]

$\Rightarrow s \cdot m + t \cdot n = 1$  with  $s$  and  $t$  efficiently

$\Rightarrow t \cdot n \equiv 1 \pmod{m}$  <sup>computed</sup> and so  $t$  is a multiplicative

**Theorem 69** For all positive integers  $m$  and  $n$ ,

inverse of  $n$  in  $\mathbb{Z}_m$ .

1.  $\gcd(m, n)$  is a linear combination of  $m$  and  $n$ , and
2. a pair  $lc_1(m, n), lc_2(m, n)$  of integer coefficients for it, i.e. such that

$$\begin{bmatrix} lc_1(m, n) & lc_2(m, n) \end{bmatrix} \cdot \begin{bmatrix} m \\ n \end{bmatrix} = \gcd(m, n) ,$$

can be efficiently computed.

**Proposition 70** For all integers  $m$  and  $n$ ,

1.  $\begin{bmatrix} ?_1 & ?_2 \end{bmatrix} \cdot \begin{bmatrix} m \\ n \end{bmatrix} = m \wedge \begin{bmatrix} ?_1 & ?_2 \end{bmatrix} \cdot \begin{bmatrix} m \\ n \end{bmatrix} = n ;$



**Proposition 70** For all integers  $m$  and  $n$ ,

1.  $\begin{bmatrix} ?_1 & ?_2 \end{bmatrix} \cdot \begin{bmatrix} m \\ n \end{bmatrix} = m \wedge \begin{bmatrix} ?_1 & ?_2 \end{bmatrix} \cdot \begin{bmatrix} m \\ n \end{bmatrix} = n ;$

2. for all integers  $s_1, t_1, r_1$  and  $s_2, t_2, r_2$ ,

$$\begin{bmatrix} s_1 & t_1 \end{bmatrix} \cdot \begin{bmatrix} m \\ n \end{bmatrix} = r_1 \wedge \begin{bmatrix} s_2 & t_2 \end{bmatrix} \cdot \begin{bmatrix} m \\ n \end{bmatrix} = r_2$$

implies  $s_1 + s_2, t_1 + t_2$

$$\begin{bmatrix} ?_1 & ?_2 \end{bmatrix} \cdot \begin{bmatrix} m \\ n \end{bmatrix} = r_1 + r_2 ;$$

**Proposition 70** For all integers  $m$  and  $n$ ,

1.  $\begin{bmatrix} ?_1 & ?_2 \end{bmatrix} \cdot \begin{bmatrix} m \\ n \end{bmatrix} = m \wedge \begin{bmatrix} ?_1 & ?_2 \end{bmatrix} \cdot \begin{bmatrix} m \\ n \end{bmatrix} = n ;$

2. for all integers  $s_1, t_1, r_1$  and  $s_2, t_2, r_2$ ,

$$\begin{bmatrix} s_1 & t_1 \end{bmatrix} \cdot \begin{bmatrix} m \\ n \end{bmatrix} = r_1 \wedge \begin{bmatrix} s_2 & t_2 \end{bmatrix} \cdot \begin{bmatrix} m \\ n \end{bmatrix} = r_2$$

implies

$$\begin{bmatrix} ?_1 & ?_2 \end{bmatrix} \cdot \begin{bmatrix} m \\ n \end{bmatrix} = r_1 + r_2 ;$$

3. for all integers  $k$  and  $s, t, r$ ,

$$\begin{bmatrix} s & t \end{bmatrix} \cdot \begin{bmatrix} m \\ n \end{bmatrix} = r \text{ implies } \begin{bmatrix} ?_1 & ?_2 \end{bmatrix} \cdot \begin{bmatrix} m \\ n \end{bmatrix} = k \cdot r .$$


## gcd

```
fun gcd( m , n )
= let
  fun gcditer(  $[s_1 \ t_1]$  r1 , c as  $[s_2 \ t_2]$  r2 )
  = let
    val (q,r) = divalg(r1,r2)    (* r = r1-q*r2 *)
  in
    if r = 0
    then c
    else gcditer( c ,  $[s_1 - q*s_2 \ t_1 - q*t_2]$  r )
  end
in
  gcditer(  $[1 \ 0]$  m ,  $[0 \ 1]$  n )
end
```

## egcd

```
fun egcd( m , n )
= let
  fun egcditer( ((s1,t1),r1) , lc as ((s2,t2),r2) )
  = let
    val (q,r) = divalg(r1,r2)    (* r = r1-q*r2 *)
  in
    if r = 0
    then lc
    else egcditer( lc , ((s1-q*s2,t1-q*t2),r) )
  end
in
  egcditer( ((1,0),m) , ((0,1),n) )
end
```

```
fun gcd( m , n ) = #2( egcd( m , n ) )
```

```
fun lc1( m , n ) = #1( #1( egcd( m , n ) ) )
```

```
fun lc2( m , n ) = #2( #1( egcd( m , n ) ) )
```

# Multiplicative inverses in modular arithmetic

**Corollary 74** *For all positive integers  $m$  and  $n$ ,*

1.  $n \cdot \text{lc}_2(m, n) \equiv \text{gcd}(m, n) \pmod{m}$ , and
2. whenever  $\text{gcd}(m, n) = 1$ ,

$[\text{lc}_2(m, n)]_m$  is the multiplicative inverse of  $[n]_m$  in  $\mathbb{Z}_m$  .