

# Divisibility and congruence

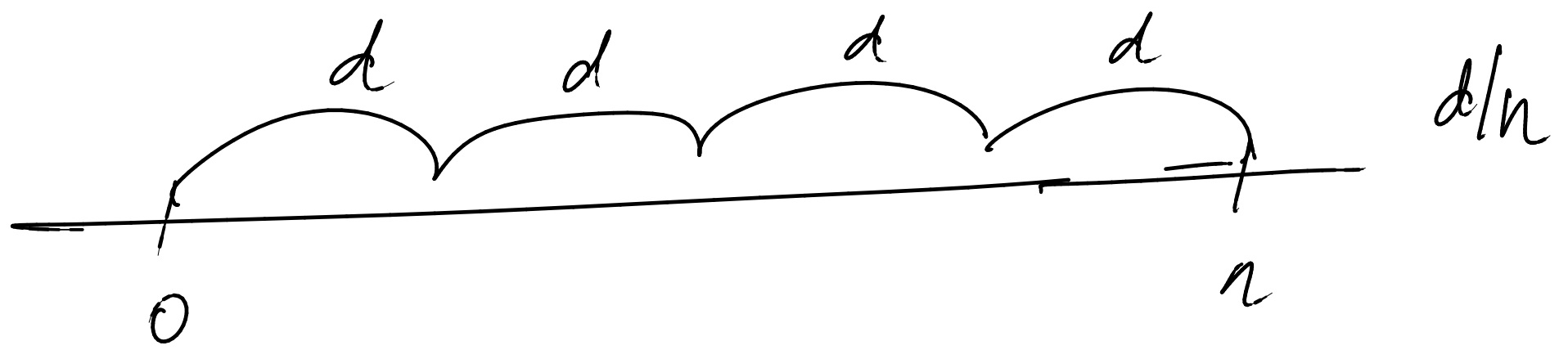
**Definition 13** Let  $d$  and  $n$  be integers. We say that  $d$  divides  $n$ , and write  $d \mid n$ , whenever there is an integer  $k$  such that  $n = k \cdot d$ .

**Example 14** The statement  $2 \mid 4$  is true, while  $4 \mid 2$  is not.

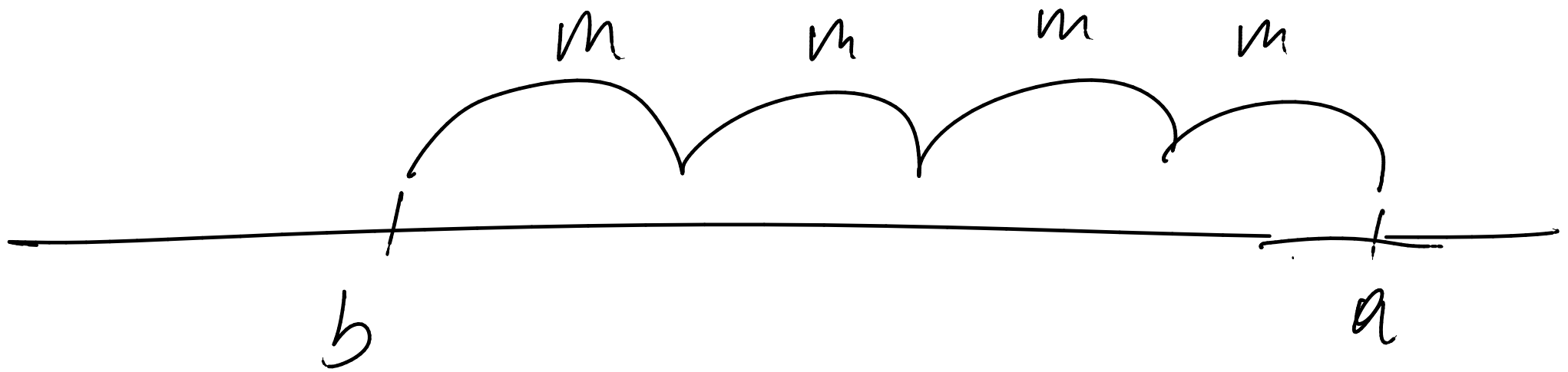
**Definition 15** Fix a positive integer  $m$ . For integers  $a$  and  $b$ , we say that  $a$  is congruent to  $b$  modulo  $m$ , and write  $a \equiv b \pmod{m}$ , whenever  $m \mid (a - b)$ .

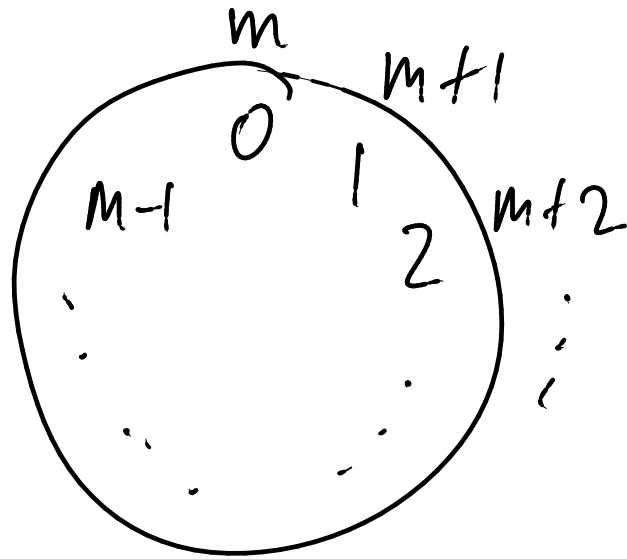
**Example 16**

1.  $18 \equiv 2 \pmod{4}$
2.  $2 \equiv -2 \pmod{4}$
3.  $18 \equiv -2 \pmod{4}$



$$a \equiv b \pmod{m}$$





$$a \equiv b \pmod{m}$$
$$b \equiv c \pmod{m}$$

then  $a \equiv c \pmod{m}$

# Universal quantification

Universal statements are of the form

**for all** individuals  $x$  of the universe of discourse,  
the property  $P(x)$  holds

or, in other words,

no matter what individual  $x$  in the universe of discourse  
one considers, the property  $P(x)$  for it holds

or, in symbols,

$\forall x. P(x)$

*equivalent to*  
 $\forall y. P(y)$

## Example 18

2. For every positive real number  $x$ , if  $x$  is irrational then so is  $\sqrt{x}$ .
3. For every integer  $n$ , we have that  $n$  is even iff so is  $n^2$ .

## The main proof strategy for universal statements:

To prove a goal of the form

$$\forall x. P(x)$$

let  $x$  stand for an arbitrary individual and prove  $P(x)$ .

## Proof pattern:

In order to prove that

$$\forall x. P(x)$$

1. **Write:** Let  $x$  be an arbitrary individual.

2. Show that  $P(x)$  holds.

Example

$$\forall x. \forall y. \forall z. P(x, y, z) \rightsquigarrow$$

Let  $x, y, z$   
be arbitrary  
individuals

### Proof pattern:

In order to prove that

$$\forall x. P(x)$$

$\rightsquigarrow$

1. **Write:** Let  $x$  be an arbitrary individual.

$P(x, y, z)$

**Warning:** Make sure that the variable  $x$  is new (also referred to as fresh) in the proof! If for some reason the variable  $x$  is already being used in the proof to stand for something else, then you must use an unused variable, say  $y$ , to stand for the arbitrary individual, and prove  $P(y)$ .

2. Show that  $P(x)$  holds.



## Scratch work:

Before using the strategy

Assumptions

⋮

Goal

$\forall x. P(x)$

After using the strategy

Assumptions

⋮

Goal

$P(x)$  (for a new (or fresh)  $x$ )

Assumptions

...

$\forall x. P(x)$

$P(a)$

$P(b)$

...

Goal  
G

## The use of universal statements:

To use an assumption of the form  $\forall x. P(x)$ , you can plug in any value, say  $a$ , for  $x$  to conclude that  $P(a)$  is true and so further assume it.

This rule is called *universal instantiation*.

**Proposition 19** Fix a positive integer  $m$ . For integers  $a$  and  $b$ , we have that  $a \equiv b \pmod{m}$  if, and only if, for all positive integers  $n$ , we have that  $n \cdot a \equiv n \cdot b \pmod{n \cdot m}$ .

PROOF:  $\forall a, b$  integers.

$$a \equiv b \pmod{m} \iff (\forall \text{ positive int. } n. n \cdot a \equiv n \cdot b \pmod{n \cdot m})$$

Let  $a$  and  $b$  be arbitrary int.

$(\implies)$  Assume  $a \equiv b \pmod{m}$   $\sim a - b = k \cdot m$  for int  $k$

Need show  $\forall \text{ positive int. } n. n \cdot a \equiv n \cdot b \pmod{n \cdot m}$

Let  $n$  be an arbitrary positive int.

Need show  $n \cdot a \equiv n \cdot b \pmod{n \cdot m}$

$$\begin{aligned} n \cdot a - n \cdot b &= n \cdot (a - b) \\ &= n \cdot (k \cdot m) \\ &= (n \cdot k) \cdot m \end{aligned}$$

for an int  $l$

Since  $a-b = km$ , it follows that

$$n(a-b) = nkm$$

$$\begin{array}{ccc} \parallel & & \parallel \\ na - nb & & k(nm) \end{array}$$

and so  $na \equiv nb \pmod{nm}$ .

?	$a \equiv b \pmod{m}$
$\Downarrow$	?
	$a/k \equiv b/k \pmod{m/k}$

( $\Leftarrow$ ) Assume  $\forall$  pos. int  $n$ .  $na \equiv nb \pmod{nm}$ .

Required to prove  $a \equiv b \pmod{m}$ .  $\Downarrow$  for  $n=1$

1.  $a \equiv 1 \cdot b \pmod{1 \cdot m}$   $\square$   
is required

$$a \equiv b \pmod{m} \stackrel{?}{\Rightarrow} a + c \equiv b + c \pmod{m}$$

$$\stackrel{?}{\Rightarrow} a \cdot c \equiv b \cdot c \pmod{m}$$

$$\stackrel{?}{\Rightarrow} a^k \equiv b^k \pmod{m}$$

$$\stackrel{?}{\Rightarrow} c^a \equiv c^b \pmod{m}$$

## Equality axioms

Just for the record, here are the axioms for *equality*.

- ▶ Every individual is equal to itself.

$$\forall x. x = x$$

- ▶ For any pair of equal individuals, if a property holds for one of them then it also holds for the other one.

$$\forall x. \forall y. x = y \implies (P(x) \implies P(y))$$

**NB** From these axioms one may deduce the usual intuitive properties of equality, such as

$$\forall x. \forall y. x = y \implies y = x$$

and

$$\forall x. \forall y. \forall z. x = y \implies (y = z \implies x = z) .$$

However, in practice, you will not be required to formally do so; rather you may just use the properties of equality that you are already familiar with.

# Conjunction

Conjunctive statements are of the form

**P and Q**

or, in other words,

**both P and also Q hold**

or, in symbols,

**$P \wedge Q$**

or

**$P \& Q$**



## The proof strategy for conjunction:

To prove a goal of the form

$$P \wedge Q$$

first prove  $P$  and subsequently prove  $Q$  (or vice versa).

## Proof pattern:

In order to prove

$$P \wedge Q$$

1. **Write:** Firstly, we prove  $P$ . and provide a proof of  $P$ .
2. **Write:** Secondly, we prove  $Q$ . and provide a proof of  $Q$ .

## Scratch work:

Before using the strategy

Assumptions

⋮

Goal

$P \wedge Q$

After using the strategy

Assumptions

⋮

Goal

$P$

Assumptions

⋮

Goal

$Q$

Assumption

?

$P \wedge Q$

$P, Q$

Goal

$G$

## The use of conjunctions:

To use an assumption of the form  $P \wedge Q$ ,  
treat it as two separate assumptions:  $P$  and  $Q$ .

**Theorem 20** For every integer  $n$ , we have that  $6 \mid n$  iff  $2 \mid n$  and  $3 \mid n$ .

PROOF:  $\forall \text{int } n. 6 \mid n \Leftrightarrow (2 \mid n \wedge 3 \mid n)$

Let  $n$  be an arbitrary integer.

RTP:  $6 \mid n \Leftrightarrow (2 \mid n \wedge 3 \mid n)$

$(\Rightarrow)$  Assume  $6 \mid n$ ; that is,  $n = 6k$  for  $k \text{ int}$

Then  $n = 2(3k)$  and  $n = 3(2k)$

( $\Leftarrow$ ) Assume  $(2|n \wedge 3|n)$

Then  $2|n$  and also  $3|n$



$$n = 2k \quad (k \text{ int})$$

$$k = 3j$$



$$n = 3l \quad (l \text{ int})$$

$$l = 2j$$

$$j = 3j - 2j = k - l$$

PROOF  
IDEA

$$6|n \rightsquigarrow n = 6j \quad (j \text{ int})$$