Complexity Theory
Easter 2016
Suggested Exercises 1

1. In the lecture, a proof was sketched showing a $\Omega(n \log n)$ lower bound on the complexity of the sorting problem. It was also stated that a similar analysis could be used to establish the same bound for the Travelling Salesman Problem. Give a detailed sketch of such an argument. Can you think of a way to improve the lower bound?

2. Say we are given a set $V = \{v_1, \ldots, v_n\}$ of vertices and a cost matrix $c : V \times V \rightarrow \mathbb{N}$. For a set $S \subseteq V$, let $t_{S,i}$ denote the cost of the shortest path that starts at v_1, visits all vertices in S and ends at v_i. Describe a *dynamic programming* algorithm that computes $t_{S,i}$ for all sets S and all i. Show that your algorithm can be used to solve the Travelling Salesman Problem in $O(n^2 2^n)$.

3. Consider the language $\mathsf{Unary\text{-}Prime}$ in the one letter alphabet $\{a\}$ defined by
$$\mathsf{Unary} - \mathsf{Prime} = \{a^n \mid n \text{ is prime}\}.$$
Show that this language is in P.

4. We say that a propositional formula ϕ is in $\mathsf{2CNF}$ if it is a conjunction of clauses, each of which contains exactly 2 literals. The point of this problem is to show that the satisfiability problem for formulas in $\mathsf{2CNF}$ can be solved by a polynomial time algorithm.

   First note that any clause with 2 literals can be written as an implication in exactly two ways. For instance $(p \vee \neg q)$ is equivalent to $(q \rightarrow p)$ and $(\neg p \rightarrow \neg q)$, and $(p \vee q)$ is equivalent to $(\neg p \rightarrow q)$ and $(\neg q \rightarrow p)$.

   For any formula ϕ, define the directed graph G_ϕ to be the graph whose set of vertices is the set of all literals that occur in ϕ, and in which there is an edge from literal x to literal y if, and only if, the implication $(x \rightarrow y)$ is equivalent to one of the clauses in ϕ.

   (a) If ϕ has n variables and m clauses, give an upper bound on the number of vertices and edges in G_ϕ.

   (b) Show that ϕ is *unsatisfiable* if, and only if, there is a literal x such that there is a path in G_ϕ from x to $\neg x$ *and* a path from $\neg x$ to x.

   (c) Give an algorithm for verifying that a graph G_ϕ satisfies the property stated in (b) above. What is the complexity of your algorithm?

(d) From (c) deduce that there is a polynomial time algorithm for testing whether or not a 2CNF propositional formula is satisfiable.

(e) Why does this idea not work if we have 3 literals per clause?

5. A clause (i.e. a disjunction of literals) is called a *Horn* clause, if it contains *at most one* positive literal. Such a clause can be written as an implication: $(x \vee (\neg y) \vee (\neg w) \vee (\neg z))$ is equivalent to $((y \wedge w \wedge z) \rightarrow x))$. HORNSAT is the problem of deciding whether a given Boolean expression that is a conjunction of Horn clauses is satisfiable.

   (a) Show that there is a polynomial time algorithm for solving HORNSAT. (Hint: if a variable is the only literal in a clause, it must be set to `true`; if all the negative variables in a clause have been set to `true`, then the positive one must also be set to `true`. Continue this procedure until a contradiction is reached or a satisfying truth assignment is found).

   (b) In the proof of the NP-completeness of SAT it was shown how to construct, for every nondeterministic machine M, integer k and string x a Boolean expression ϕ which is satisfiable if, and only if, M accepts x within n^k steps. Show that, if M is deterministic, than ϕ can be chosen to be a conjunction of Horn clauses.

   (c) Conclude from (b) that the problem HORNSAT is P-complete under L-reductions.

6. We define the complexity class of *quasi-polynomial-time* problems Quasi-P by:

$$\mathsf{Quasi\text{-}P} = \bigcup_{k=1}^{\infty} \mathsf{Time}(n^{(\log n)^k}).$$

   Show that if $L_1 \leq_P L_2$ and $L_2 \in \mathsf{Quasi\text{-}P}$, then $L_1 \in \mathsf{Quasi\text{-}P}$.

7. In general k-colourability is the problem of deciding, given a graph $G = (V, E)$, whether there is a colouring $\chi : V \rightarrow \{1, \ldots, k\}$ of the vertices such that if $(u, v) \in E$, then $\chi(u) \neq \chi(v)$. That is, adjacent vertices do not have the same colour.

   (a) Show that there is a polynomial time algorithm for solving 2-colourability.

   (b) Show that, for each k, k-colourability is reducible to $k+1$-colourability. What can you conclude from this about the complexity of 4-colourability?