Object Oriented Programming Dr Robert Harle

IA CST, PBS (CS) and NST (CS)
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The OOP Course

- So far you have studied functional programming (ML)
- Now we consider imperative programming (Java primarily but not exclusively).
- You have practicals in Java
 - This course complements the practicals
 - Some material appears only here
 - Some material appears only in the practicals
 - Some material appears in both: deliberately*!

^{*} Some material may be repeated unintentionally. If so I will claim it was deliberate.

Outline

- 1. Types, Objects and Classes
- 2. Pointers, References and Memory
- 3. Creating Classes
- 4. Inheritance
- 5. Polymorphism
- 6. Lifecycle of an Object
- 7. Error Handling
- 8. Copying Objects
- 9. Java Collections
- 10. Object Comparison
- 11. Design Patterns
- 12. Design Pattern (cont.)

Books and Resources I

OOP Concepts

- Look for books for those learning to first program in an OOP language (Java, C++, Python)
- Java: How to Program by Deitel & Deitel (also C++)
- Thinking in Java by Eckels
- Java in a Nutshell (O' Reilly) if you already know another OOP language
- Java specification book: http://java.sun.com/docs/books/jls/
- Lots of good resources on the web
- Design Patterns
 - Design Patterns by Gamma et al.
 - Lots of good resources on the web

Books and Resources II

- Also check the course web page
 - Updated notes (with annotations where possible)
 - Code from the lectures
 - Sample tripos questions

http://www.cl.cam.ac.uk/teaching/current/OOProg/

Lecture 1: Types, Objects and Classes

Types of Languages

- Declarative specify <u>what</u> to do, not <u>how</u> to do it. i.e.
 - E.g. HTML describes what should appear on a web page, and not how it should be drawn to the screen
 - E.g. SQL statements such as "select * from table" tell a program to get information from a database, but not how to do so
- Imperative specify <u>both</u> what and how
 - E.g. "double x" might be a declarative instruction that you want the variable x doubled somehow.
 Imperatively we could have "x=x*2" or "x=x+x"

ML as a Functional Language

- Functional languages are a subset of declarative languages
 - ML is a functional language
 - It may appear that you tell it how to do everything, but you should think of it as providing an explicit example of what should happen
 - The compiler may optimise i.e. replace your implementation with something entirely different but 100% equivalent.

Aside: Some things to note on Java

- Programs are compiled and then run (i.e. not interpreted 'live' as per ML).
 - 1) Write your .java code (e.g. Test.java)
 - 2) Compile using javac Test.java
 - 3) Run using java Test
- You need an explicit start point. This is called the main() function
- You have to name your .java files <u>exactly</u> the same as the class they define (this will make more sense later)

More Boilerplate than ML

```
HelloWorld.ml: print "Hello World";
```

```
HelloWorld.java:
public class HelloWorld {
   public static void main(String[] args) {
      System.out.println("Hello World!");
   }
}
```

Types and Variables

Most imperative languages don't have type inference

```
int x = 512;
int y = 200;
int z = x+y;
```

- The high-level language has a series of primitive (built-in) types that we use to signify what's in the memory
 - The compiler then knows what to do with them
 - E.g. An "int" is a primitive type in C, C++, Java and many languages. It's usually a 32-bit signed integer
- A variable is a name used in the code to refer to a specific instance of a type
 - x,y,z are variables above
 - They are all of type int

E.g. Primitive Types in Java

- "Primitive" types are the built in ones.
 - They are building blocks for more complicated types that we will be looking at soon.
- boolean 1 bit (true, false)
- char 16 bits
- char 16 bits

 byte 8 bits as a signed integer (-128 to 127)

 Two's

 complement
- short 16 bits as a signed integer
- int 32 bits as a signed integer
- long 64 bits as a signed integer
- float 32 bits as a floating point number
- double 64 bits as a floating point number



...vs ML

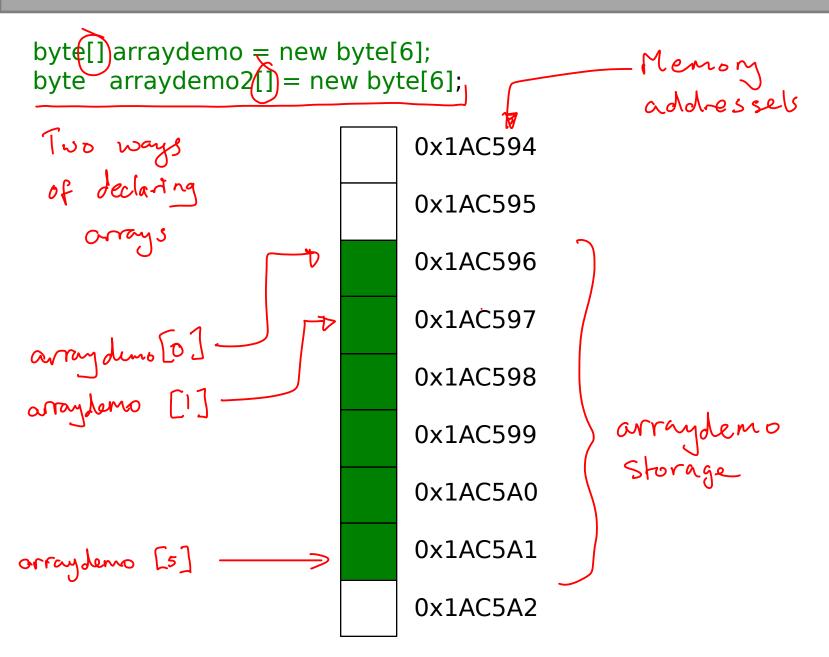
- bool 1 bit (true, false)
- int arbitrary-precision integer
- real arbitrary-precision floating point

Much simpler to work with but incurs a cost over working with fixed precision

Immutable to Mutable Data

```
ML
   - val x=5; ____
    > val x = 5: int
                        Compaison
   - x=7; _____
    > val it = false : bool
   - val x=9;
    > val x = 9: int
Java
   int x=5;
                      - assignment
   x=7; ____
   10t \times = 9
                                  (x = = 9) false (x = = 5) true
    wonst
    compile
```

Arrays



Function Prototypes

- Functions are made up of a prototype and a body
 - Prototype specifies the function name, arguments and possibly return type
 - Body is the actual function code

```
fun myfun(a,b) = ...;

int myfun(int a, int b) {...}

protobyp

return

type
```

Overloading Functions

- Same function name
- Different arguments
- Possibly different return type

```
int myfun(int a, int b) {...}
float myfun(float a, float b) {...}
double myfun(double a, double b) {...}
```

But <u>not</u> just a different return type

```
int myfun(int a, int b) {...}
float myfun(int a, int b) {...}
```

Function Side Effects

 Functions in imperative languages can use or alter larger system state → procedures

```
m(x,y) = xy
Maths:
                                                 Proper functions
(output depends
              fun m(x,y) = x*y;
ML:
              int m(int x, int y) = x*y;
Java:
              int y = 7;
              int m(x) {
                     y=y+1;
                      return x*y;
```

void Procedures

A void procedure returns nothing:

```
int count=0;

void addToCount() {
   count=count+1;
}

Meaning ful only
for procedures
(i.e. not in ml)
```

Control Flow: Decision Making

```
if (boolean_expression) {
    do_something()
}

if (boolean_expression) {
    do_something()
}
else {
    do_something_else()
}
```

Control Flow: Looping

```
for( initialisation; termination; increment )
                for (int i=0; i<8; i++) ...
                int j=0; for(; j<8; j++) ...
                for(int k=7; k>=0; (-) ...
                pan_expression)

post-increment equivalent to int i=0; while (i<8) \{(1++),...\}

i ; i=i+1;
while( boolean_expression )
                int j=7; while (j>=0) { j--; ...}
```

Control Flow: Looping Examples

```
Shorthand
for new int[5];
int arr[] = {1,2,3,4,5};
                                     Built in
length value
for arrays
for (int i=0; i<arr.length)i++) {
        System.out.println(arr[i]);
int i=0;
while (i<arr.length) {
        System.out.println(arr[i]);
        i=i+1;
```

Control Flow: Branching I

- Branching statements interrupt the current control flow
- return
 - Used to return from a function at any point

```
boolean linearSearch(int[] xs, int v) {
   for (int i=0;i<xs.length; i++) {
     if (xs[i]==v) return true;
   }
   return false;
}</pre>
```

Control Flow: Branching II

- Branching statements interrupt the current control flow
- break
 - Used to jump out of a loop

```
booleen
 void linearSearch(int[] xs, int v) {
     boolean found=false;
     for (int i=0;i<xs.length; i++) {</pre>
        if (xs[i]==v) {
              found=true;
              break; // stop looping
     return found; Allows us to keep just one return point
```

Control Flow: Branching III

Branching statements interrupt the current control flow

continue

Used to skip the current iteration in a loop

```
void printPositives(int[] xs) {
   for (int i=0;i<xs.length; i++) {
     if (xs[i]<0) continue;
     System.out.println(xs[i]);
   }
}</pre>
```

Custom Types

```
datatype 'a seq = Nil
               | Cons of 'a * (unit -> 'a seq);
              Custom type
public class Vector3D {
 float x;
float y;
float z;
```

State and Behaviour

```
datatype 'a seq = Nil

| Cons of 'a * (unit -> 'a seq);

fun hd (Cons(x,_)) = x;
```

```
public class Vector3D {
  float x;
  float y;
  float z;

  void add(float vx, float vy, float vz) {
     x=x+vx;
     y=y+vy;
     z=z+vz;
  }
}
```

grouped

Classes, Instances and Objects

- Classes can be seen as templates for representing various concepts
- We create *instances* of classes in a similar way. e.g.

```
MyCoolClass m = new MyCoolClass(); a mycoolClass n = new MyCoolClass(); object
```

makes two instances of class MyCoolClass.

An instance of a class is called an **object**

of an object

Loose Terminology (again!)

State
Fields
Instance Variables
Properties
Variables
Members

Behaviour
Functions
Methods
Procedures

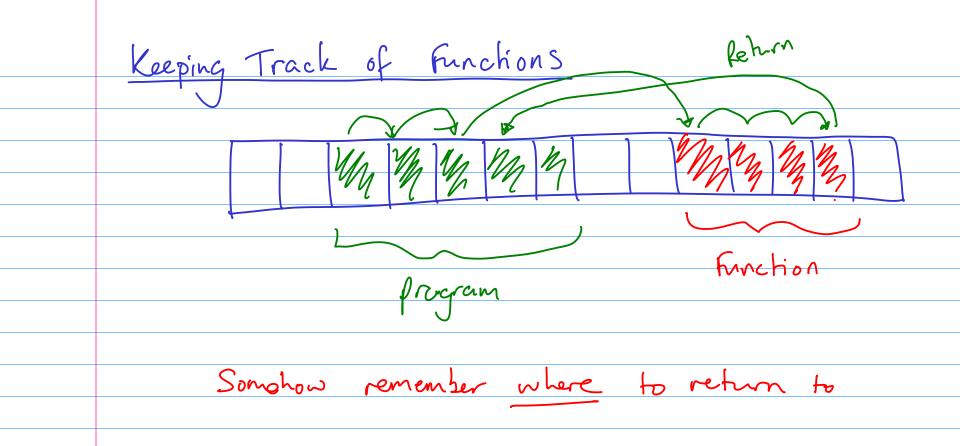
Parameterised Classes

 ML's polymorphism allowed us to specify functions that could be applied to multiple types

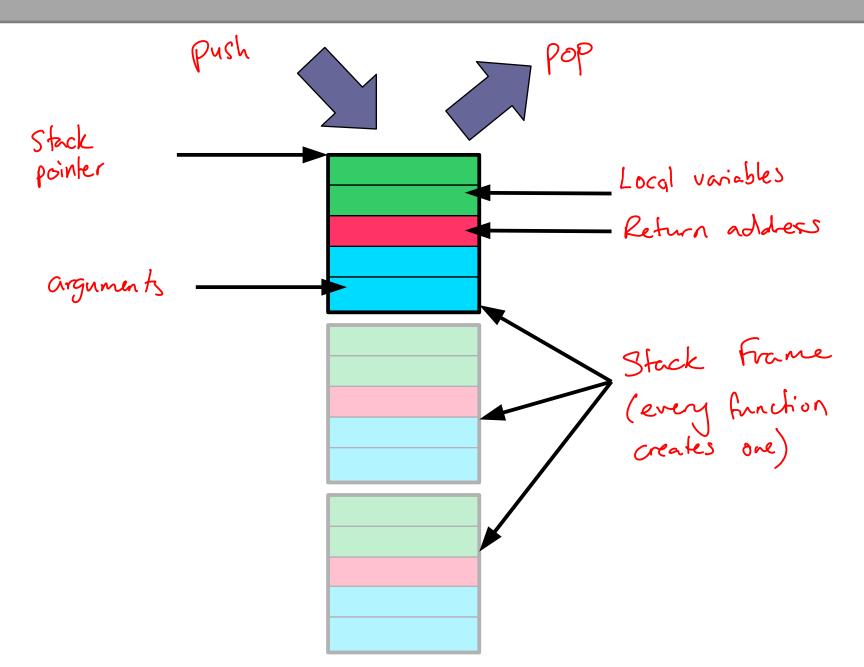
```
> fun self(x)=x;
val self = fn : 'a -> 'a
```

- In Java, we can achieve something similar through Generics; C++ through templates
 - Classes are defined with placeholders (see later lectures)
 - We fill them in when we create objects using them LinkedList<Integer> = new LinkedList<Integer>()
 LinkedList<Double> = new LinkedList<Double>()

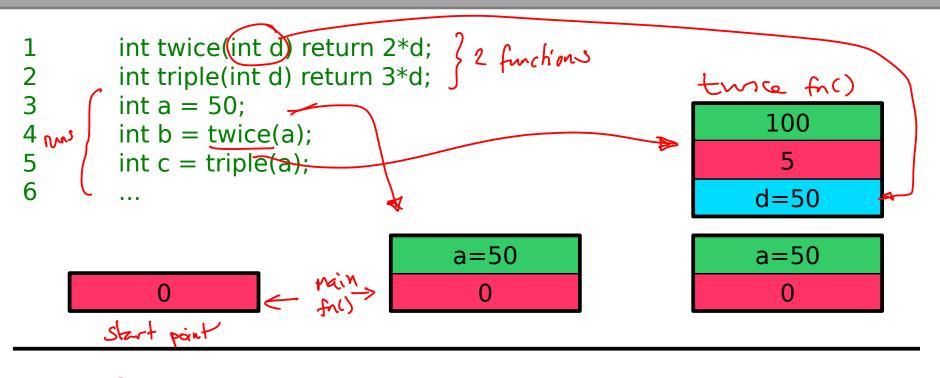
Lecture 2: Pointers, References and Memory



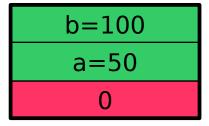
The Call Stack



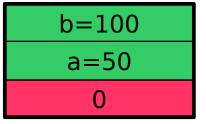
The Call Stack: Example





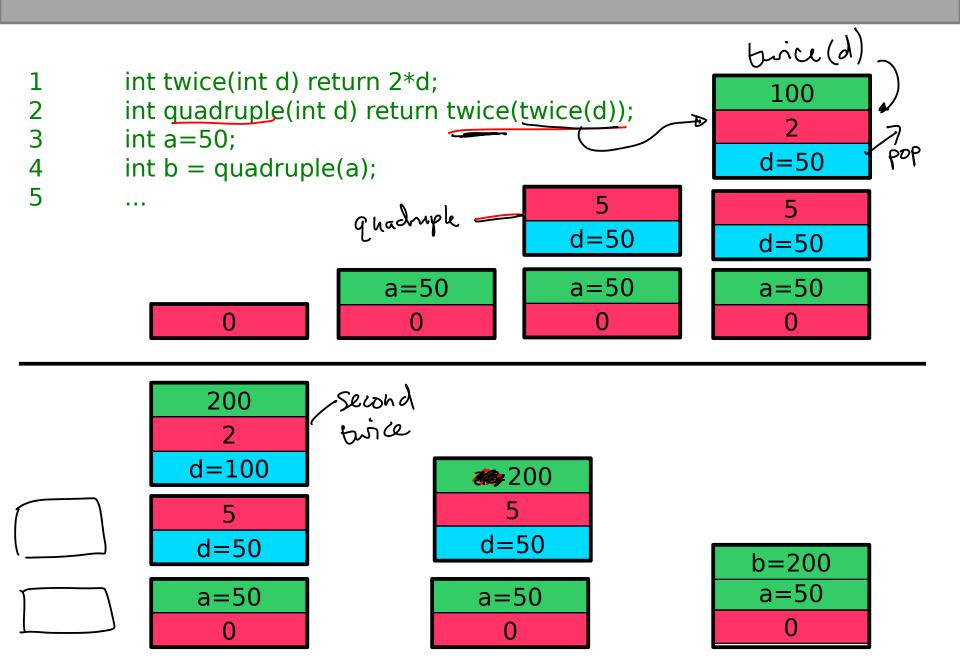


150 6 d=50



c=150 b=100 a=50

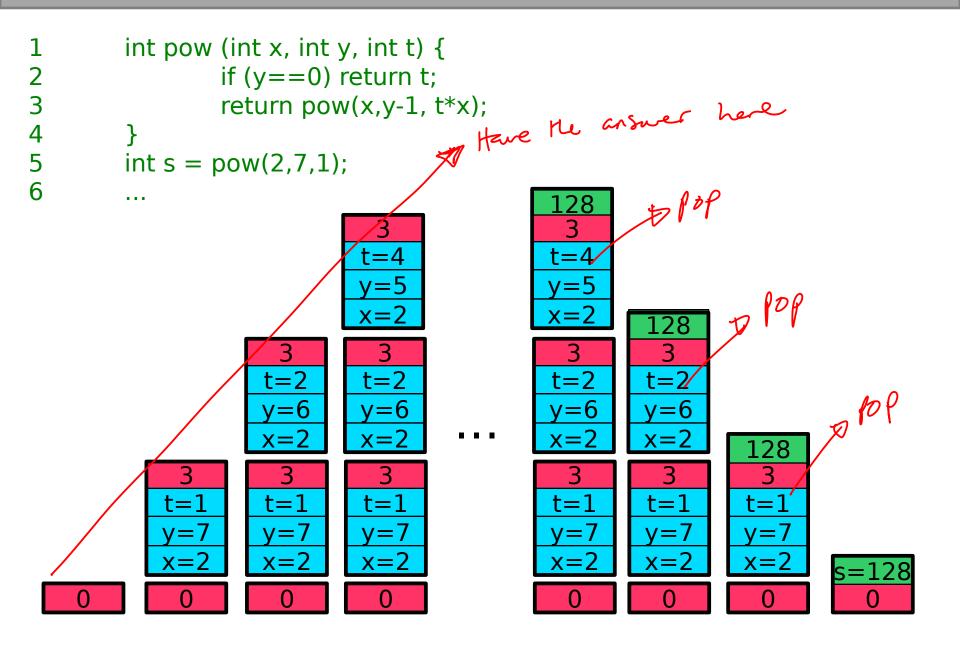
Nested Functions



Recursive Functions

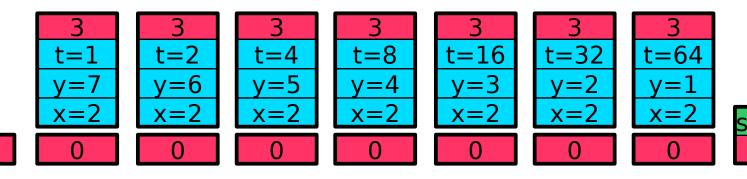
```
int pow (int x, int y) {
234567
                if (y==0) return 1;
               int p = pow(x,y-1);
                return x*p;
        int s=pow(2,7);
        . . .
                                  x=2
                                                p = 16
                                                 y=5
                          y=5
                                  y=5
                          x=2
                                  x=2
                                                x=2
                                                        p = 32
                           4
                                  y=6
                                                         y=6
                                                 y=6
                 y=6
                          y=6
         154
                                                x=2
                 x=3
                          x=2
                                  x=2
                                                         x=2
                                  y=7
                                                         y=7
                 y=7
                          y=7
         y=7
         x=2
                 x=2
                          x=2
                                                x=2
                                                         x=2
                                  x=2
```

Tail-Recursive Functions I



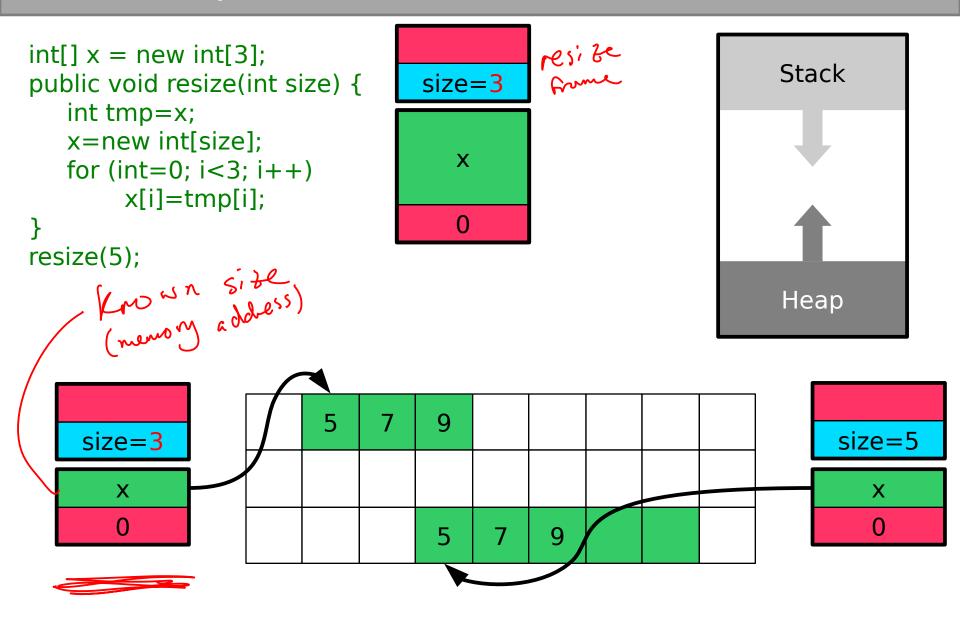
Tail-Recursive Functions II

```
int pow (int x, int y, int t) {
        if (y==0) return t;
        return pow(x,y-1, t*x);
}
int s = pow(2,7,1);
...
```



The problem with the stack You can only operate on the top. So resiting an entity (e.g. array) elsewhere not poss. Con't resize without popping!

The Heap

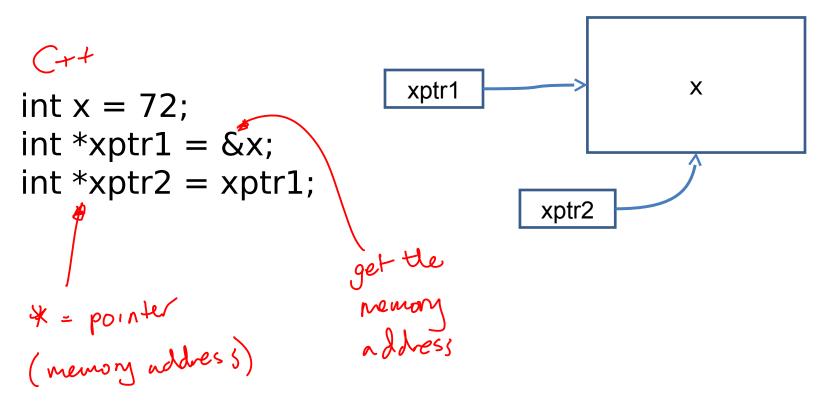


Memory and Pointers

- In reality the compiler stores a mapping from variable name to a specific memory address, along with the type so it knows how to interpret the memory (e.g. "x is an int so it spans 4 bytes starting at memory address 43526").
- Lower level languages often let us work with memory addresses directly. Variables that store memory addresses are called pointers or sometimes references
- Manipulating memory directly allows us to write fast, efficient code, but also exposes us to bigger risks
 - Get it wrong and the program 'crashes'.

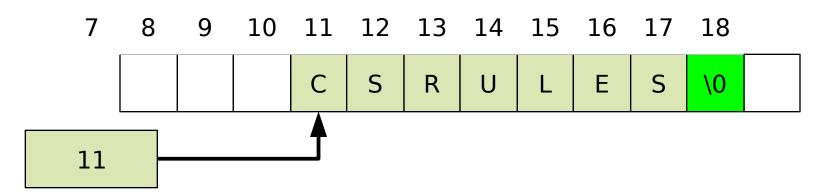
Pointers: Box and Arrow Model

- A pointer is just the memory address of the first memory slot used by the variable
- The pointer type tells the compiler how many slots the whole object uses



Example: Representing Strings I

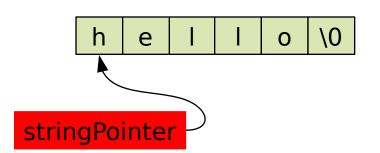
- A single character is fine, but a text string is of variable length how can we cope with that?
- We simply store the start of the string in memory and require it to finish with a special character (the NULL or terminating character, aka '\0')
- So now we need to be able to store memory addresses → use pointers



 We think of there being an array of characters (single letters) in memory, with the string pointer pointing to the first element of that array

Example: Representing Strings II

```
char letterArray[] = {'h','e','l','l','o','\0'};
char *stringPointer = &(letterArray[0]);
printf("%s\n",stringPointer);
letterArray[3]='\0';
printf("%s\n",stringPointer);
```



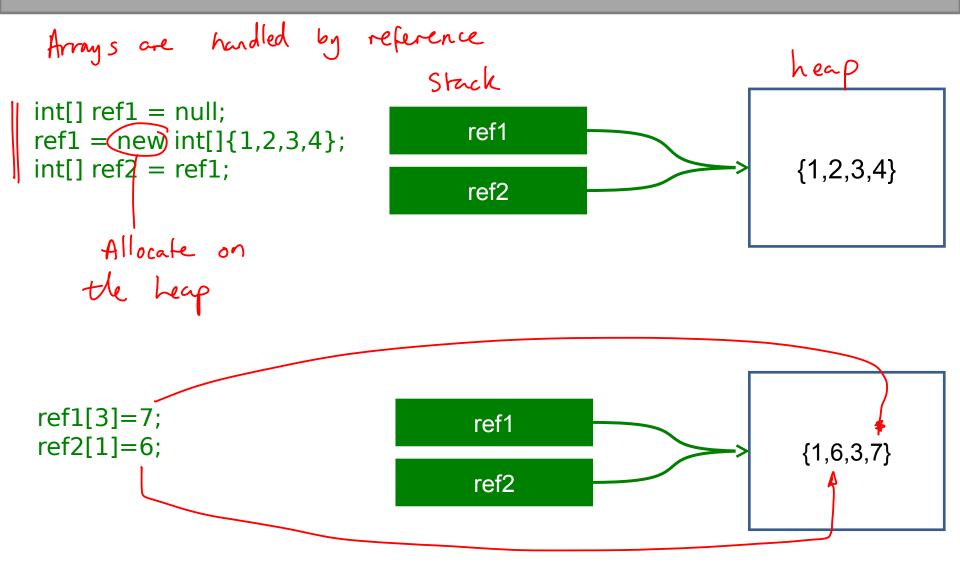
References

- Pointers are useful but dangerous
- References can be thought of as restricted pointers
 - Still just a memory address
 - But the compiler limits what we can do to it
- C, C++: pointers and references
- Java: references <u>only</u>
- ML: references <u>only</u>

References vs Pointers

	Pointers	References
Represents a memory address	Yes	Yes
Can be arbitrarily assigned	Yes	No
Can be assigned to established object	Yes	Yes
Can be tested for validity	No A	Yes 4
1. Valid ad 2. NMLL 3. Random	duess (invalid)	1. Valid address 2. Null

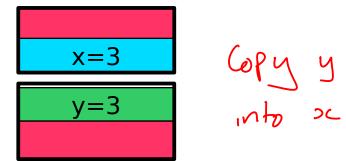
References Example (Java)



Argument Passing

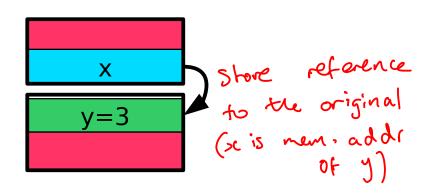
 Pass-by-value. Copy the object into a new value in the stack

```
void test(int x) {...}
int y=3;
test(y);
```



Pass-by-reference. Create a reference to the object and pass that.

```
void test(int &x) {...}
int y=3;
test(y);
```



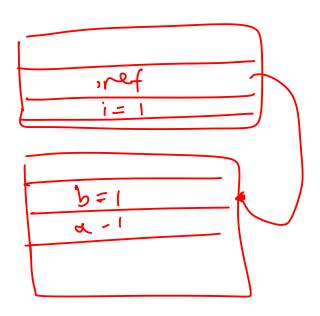
Passing Procedure Arguments In Java

```
class Reference {
 public static void update(int i, int[] array) {
   i++;
   array[0]++;
                                                                    Stach
                                                       test-array
 public static void main(String[] args) {
                                               Main
   int test i = 1;
   int[] test array = \{1\};
   update(test i, test array);
   System.out.println(test array[0]);
```

Passing Procedure Arguments In C++

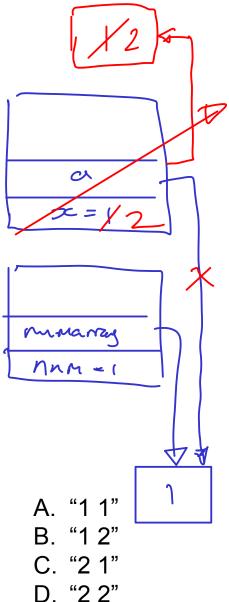
```
void update(int i, int &iref){
  i++;
  iref++;
}

int main(int argc, char** argv) {
  int a=1;
  int b=1;
  update(a,b);
  printf("%d %d\n",a,b);
}
```



Check...

```
public static void myfunction2(int x, int[] a) {
      x=1;
       x=x+1;
       a = new int[]{1};
       a[0]=a[0]+1;
public static void main(String[] arguments) {
       int num=1;
       int numarray[] = \{1\};
       myfunction2(num, numarray);
       System.out.println(num+" "+numarray[0]);
```



Lecture 3: Creating Classes

What Not to Do

- Your ML has doubtless been one big file where you threw together all the functions and value declarations
- Lots of C programs look like this :-(
- We could emulate this in OOP by having one class and throwing everything into it

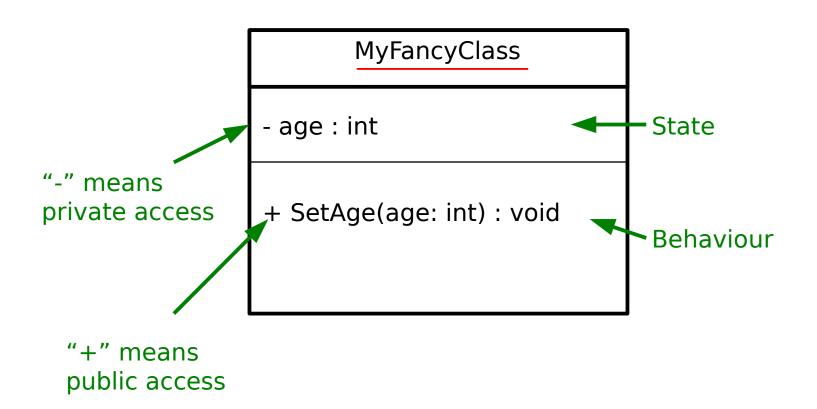
We can do (much) better

Identifying Classes

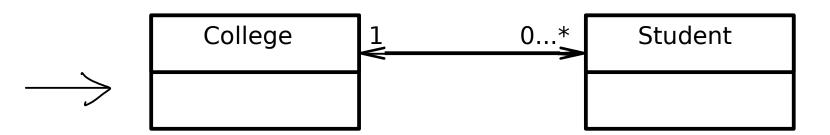
- We want our class to be a grouping of conceptually-related state and behaviour
- One popular way to group is using grammar
 - Noun → Object
 - Verb → Method

"A <u>simulation</u> of the <u>Earth</u>'s orbit around the <u>Sun</u>"

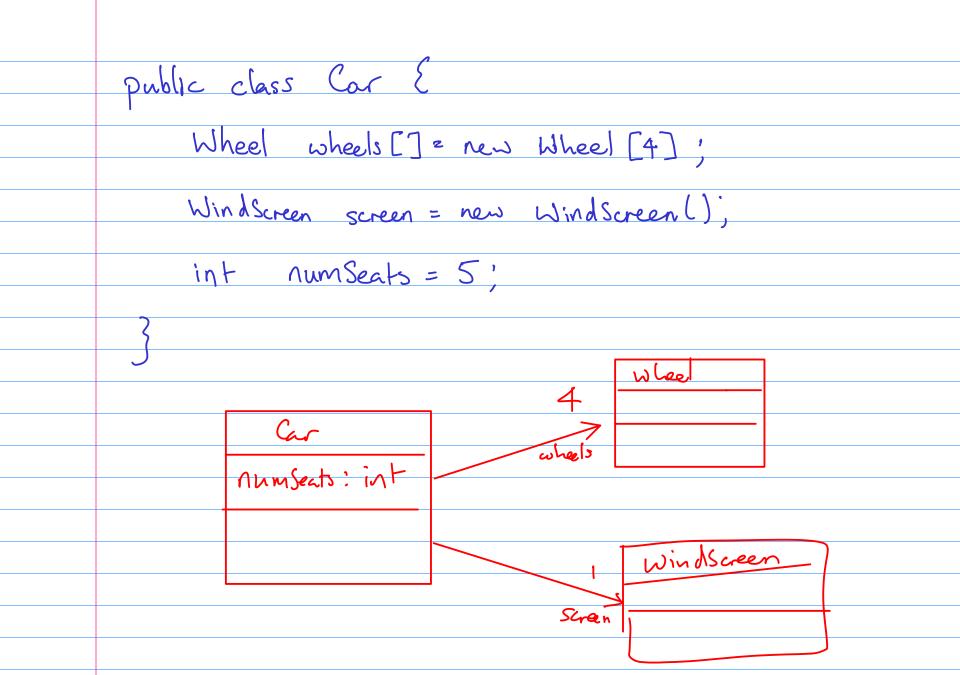
UML: Representing a Class Graphically



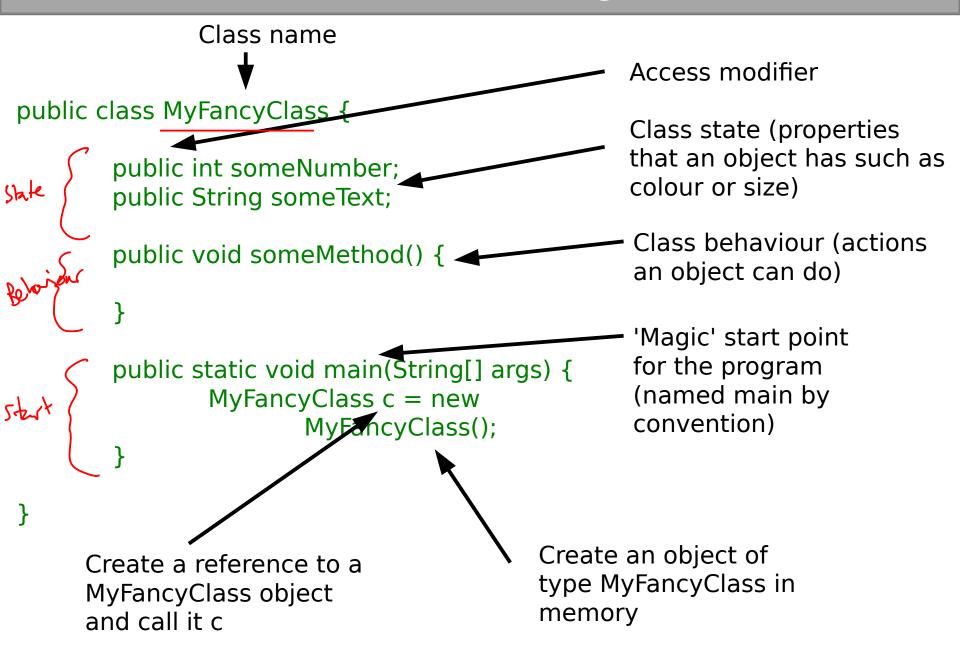
The has-a Association



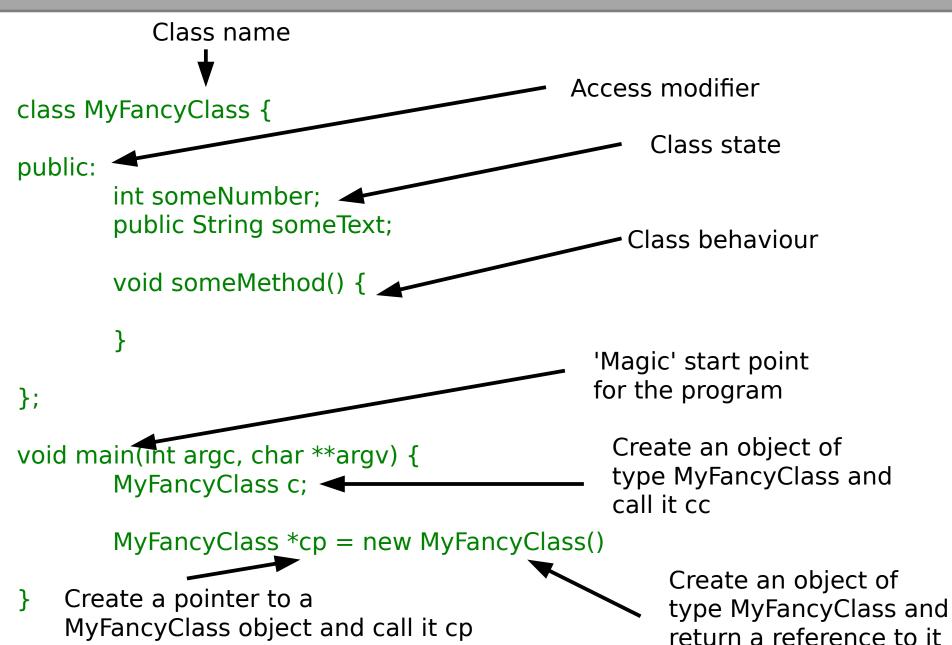
- Arrow going left to right says "a College has zero or more students"
- Arrow going right to left says "a Student has exactly 1 College"
- What it means in real terms is that the College class will contain a variable that somehow links to a set of Student objects, and a Student will have a variable that references a College object.
- Note that we are only linking classes: we don't start drawing arrows to primitive types.



Anatomy of an OOP Program (Java)



Anatomy of an OOP Program (C++)



Accessing the state and Behaviour Public class Test & Test t = new Test(); int x = 5; t.x=7; t.setx(3);void setx (int y) { X=Y;

OOP Concepts

- OOP provides the programmer with a number of important concepts:
 - Modularity
 - Code Re-Use
 - Encapsulation
 - Inheritance
 - Polymorphism
- Let's look at these more closely...

Modularity and Code Re-Use

- You've long been taught to break down complex problems into more tractable subproblems.
- Each class represents a sub-unit of code that (if written well) can be developed, tested and updated independently from the rest of the code.
- Indeed, two classes that achieve the same thing (but perhaps do it in different ways) can be swapped in the code
- Properly developed classes can be used in other programs without modification.



Encapsulation I

```
class Student {
 int age;
};
void main() {
 Student s = new Student();
 s.age = 21;
 Student s2 = new Student();
 s2.age=-1;
 Student s3 = new Student();
 s3.age=10055;
```

Encapsulation II

```
class Student {
  private int age;
                              Sanity
checking
of inputs
  boolean SetAge(int a) {
    if (a \ge 0 \&\& a < 130) {
       age=a;
       return true;
    return false;
  int GetAge() {return age;}
void main() {
  Student s = new Student();
  s.SetAge(21);
```

Encapsulation III

```
class Location {
    private float x;
    private float y;

    float getX() {return x;}
    float getY() {return y;}

    void setX(float nx) {x=nx;}
    void setY(float ny) {y=ny;}
}
class Location {
    private Vector2D v;

    float getX() {return v.getX();}
    float getY() {return v.getY();}

void setX(float nx) {v.setX(nx);}
    void setY(float ny) {v.setY(ny);}
```

Access Modifiers

	Everyone	Subclass	Same package (Java)	Same Class
private				X
package (Java)			X	X
protected		X	X	X
public	X	X	X	X

Immutability

- Everything in ML was immutable (ignoring the reference stuff). Immutability has a number of advantages:
 - Easier to construct, test and use
 - Can be used in concurrent contexts
 - Allows lazy instantiation
- We can use our access modifiers to create immutable classes

Creating Parameterised Types

These just require a placeholder type

```
class Vector3D<T> {
  private T x;
  private T y;

T getX() {return x;}
T getY() {return y;}

void setX(T nx) {x=nx;}
  void setY(T ny) {y=ny;}
}
```

Class-Level Data and Functionality I

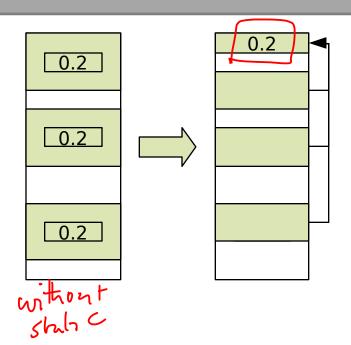
 A static field is created only once in the program's execution, despite being declared as part of a class

```
public class ShopItem {
    private float mVATRate;
    private static float sVATRate;
    ....
}

One of these created every time a new ShopItem is instantiated. Nothing keeps them all in sync.
```

Only one of these created ever. Every ShopItem object references it.

Class-Level Data and Functionality II



- Auto synchronised across instances
- Space efficient

Also static methods:

```
public class Whatever {
  public static void main(String[] args) {
    ...
  }
}
```

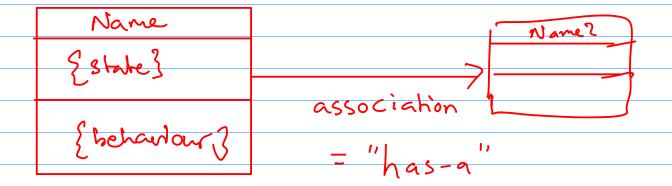
Why use Static Methods?

- Easier to debug (only depends on static state)
- Self documenting
- Groups related methods in a Class without requiring an object
- The compiler can produce more efficient code since no specific object is involved

```
public class Math {
  public float sqrt(float x) {...}
  public double sin(float x) {...}
  public double cos(float x) {...}
  public static float sqrt(float x) {...}
  public static float sin(float x) {...}
  public static float cos(float x) {...}
  public static float sqrt(float x) {...}
  public static float cos(float x) {...}
  public static float sqrt(float x) {...}
  public static float sqrt(float x) {...}
  public static float cos(float x) {...}
  public static float sqrt(float x) {...}
  public static float cos(float x) {...}
  public static float cos
```



- Oof types (classes) are groupings of state
- A good class represents a single concept



- We create multiple objects from each class
- · Encapsulation is a key concept that encourages us to separate implementation from interface

- Encapsulation allows us to refactor (rewrite) our classes without everyone who uses them changing their code
- · Encapsulation allows us to sonity check inputs.
- Static things are those that make sense for the concept but which are the same for every instance, i.e. Instance—independent state or behaviour.

Lecture 4: Inheritance

Inheritance I

```
class Student {
   public int age;
   public String name;
   public int grade;
}

class Lecturer {
   public int age;
   public String name;
   public int salary;
}
```

- There is a lot of duplication here
- Conceptually there is a hierarchy that we're not really representing
- Both Lecturers and Students are people (no, really).
- We can view each as a kind of specialisation of a general person
 - They have all the properties of a person
 - But they also have some extra stuff specific to them

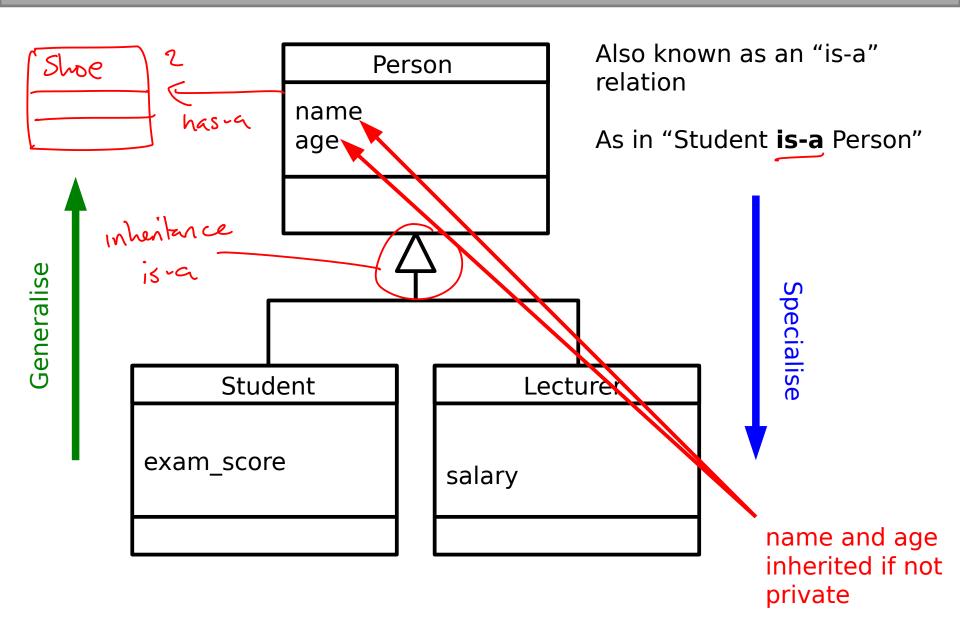
Inheritance II

```
class Person {
  public int age;
 Public String name;
class Student extends Person {
  public int grade;
class Lecturer extends Person {
  public int salary;
```

- We create a base class (Person) and add a new notion: classes can inherit properties from it
 - Both state and functionality
- We say:
 - Person is the superclass of Lecturer and Student
 - Lecturer and Student subclass Person

inhentance

Representing Inheritance Graphically



UML Amows You Need Fo Know
Association > "has -a"
as in "X has-a Y
Inhentance
→ "is-="1"
as in "X is a Y"

Java Oddity: Object	
"Magic' class Object	
Every class extends Object	
Object	
Myclass public class My Class {	

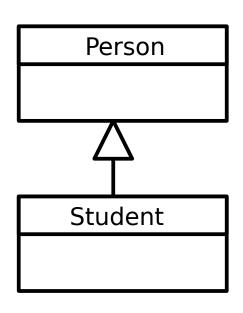
Casting

 Many languages support type casting between numeric types

```
int i = 7; float f = (float) i; f = 7.0 \text{ s. up cast} double d = 3.2; int i2 = (int) d; f = 7.0 \text{ s. up cast} int i2 = (int) d; f = 7.0 \text{ s. up cast} f = 7.0 \text{ s. up cast}
```

 With inheritance it is reasonable to type cast an object to any of the types above it in the inheritance tree...

Widening



- Student is-a Person
- Hence we can use a Student object anywhere we want a Person object
- Can perform widening conversions (up the tree)

```
Student s = new Student()

Person p = (Person) s;

Student s = new Student();

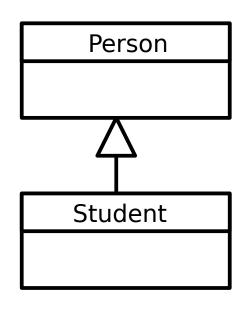
print(s);

Person part

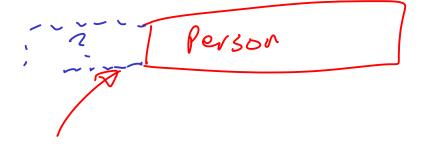
Small Person part

Implicit cast
```

Narrowing



- Narrowing conversions move down the tree (more specific)
- Need to take care...



Person p = new Person();

Student s = (Student) p;



FAILS. Not enough info In the real object to represent a Student Student s = new Student(); Person p = (Person) s; Students s2 = (Student) p;



OK because underlying object really is a Student

of all people
Linked List < Person) p
Linked List < Person) p = new Lnked List < Person > ()
1p. add (new Student()); 1p. add (new Lechnerl));
for (inti=o; 1< 1p.size(); i++) { person p = 1p.get(i)
}

Fields and Inheritance

```
class Person {
 public String mName;
 protected int mAge;
 private double mHeight;
class Student extends Person {
 public void do something()
  mName="Bob";
  mAge=70; ✓
  mHeight=1.70;)
        won't compile
```

Student inherits this as a public variable and so can access it

Student inherits this as a protected variable and so can access it

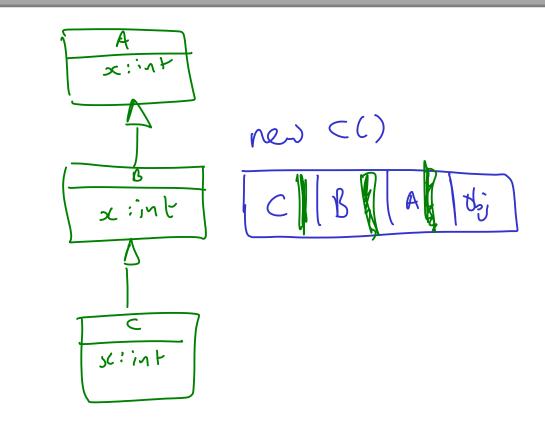
Student inherits this but as a **private** variable and so cannot access it directly

Person .

public - Everytting external can access protected - subclasses can access (+ anything in same package in Java) private - Nothing external can access (package) - Only things in same package

Fields and Inheritance: Shadowing

```
class A { public int x; }
class B extends A {
  public int x;
class C extends B {
 public int x;
 public void action() {
   // Ways to set the x in C
    x = 10;
    this.x = 10;
    // Ways to set the x in B
    super.x = 10;
    ((B)this).x = 10;
    // Ways to set the x in A
    ((A)this).x = 10;
```



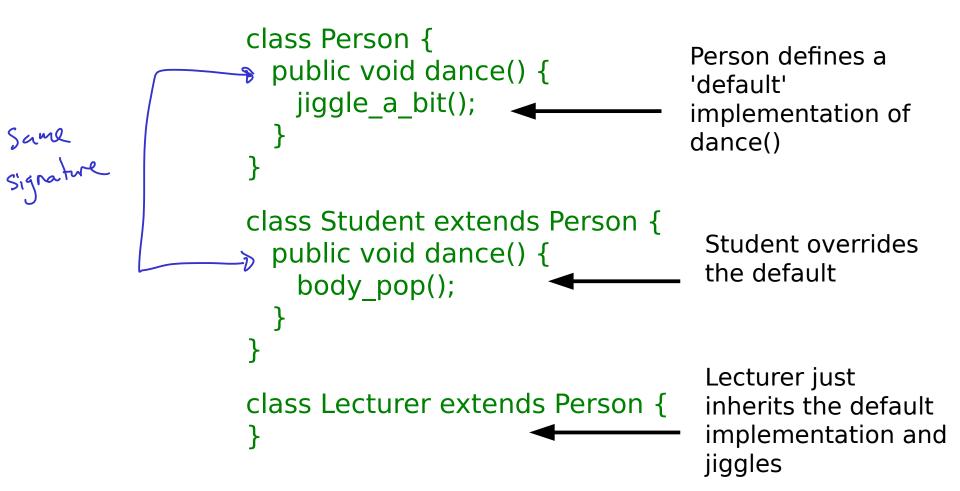
Super, super, super

this: The current object (ref to)

Super: Get a ref. to the superclass object

Methods and Inheritance: Overriding

We might want to require that every Person can dance. But the way a Lecturer dances is not likely to be the same as the way a Student dances...



Annotations Override public void dancel) §

Abstract Methods

- Sometimes we want to force a class to implement a method but there isn't a convenient default behaviour
- An abstract method is used in a base class to do this
- It has no implementation whatsoever

```
class abstract Person {
 public abstract void dance(); 
class Student extends Person {
 public void dance() {
   body pop();
class Lecturer extends Person {
  public void dance() {
   jiggle a bit();
```

Abstract Classes

Note that I had to declare the class abstract too. This is because it has a method without an implementation so we can't directly instantiate a Person.

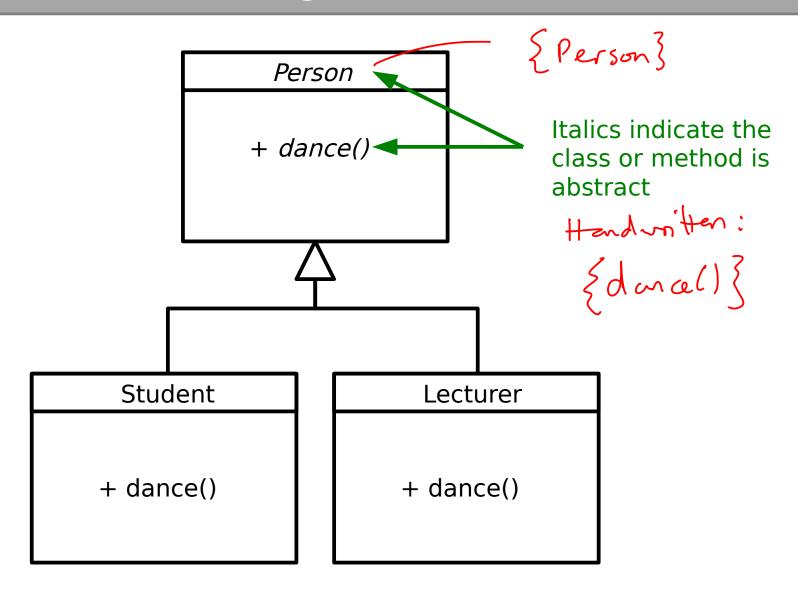
```
public abstract class Person {
   public abstract void dance();
}

Java

class Person {
   public:
     virtual void dance()=0;
   }
   C++
```

- All state and non-abstract methods are inherited as normal by children of our abstract class
- Interestingly, Java allows a class to be declared abstract even if it contains no abstract methods!

Representing Abstract Classes



Lecture 5: Polymorphism and Multiple Inheritance

Polymorphic Methods

```
Student s = new Student();
Person p = (Person)s;
p.dance();
```

Assuming Person has a default dance() method, what should happen here??

General problem: when we refer to an object via a parent type and both types implement a particular method: which method should it run?

Polymorphic Concepts I

- Static polymorphism
 - Decide at <u>compile-time</u>
 - Since we don't know what the true type of the object will be, we just run the parent method
 - Type errors give compile errors

```
Student s = new Student();

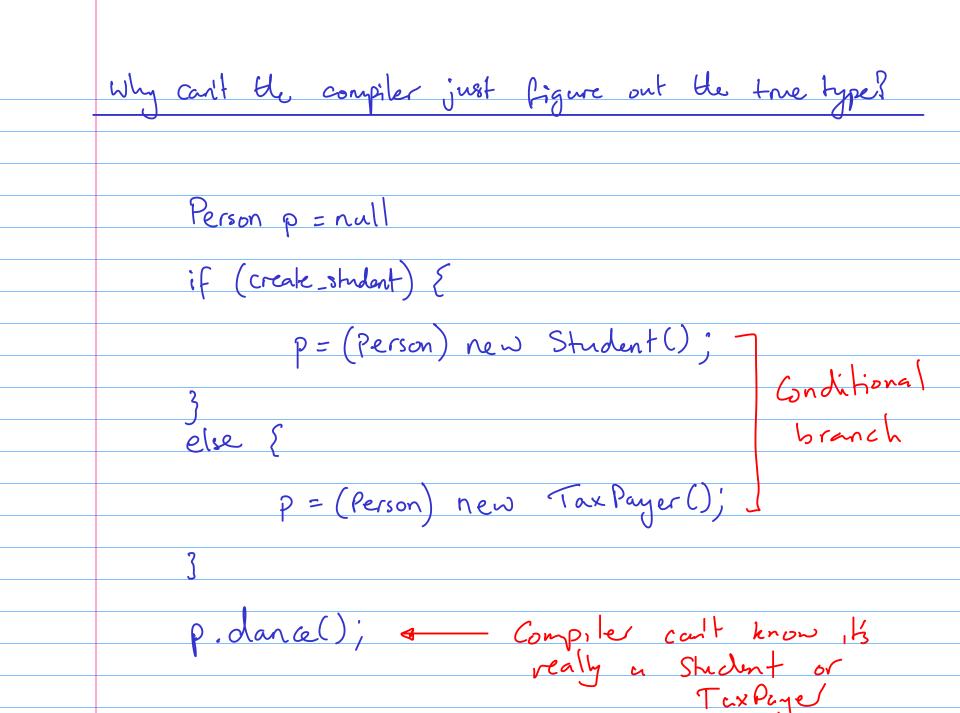
Person p = (Person)s;

p.dance();

p is a lesson

i. who lesson dence()
```

- Compiler says "p is of type Person"
- So p.dance() should do the default dance() action in Person



Statte Polymorphism You've Seen fun cons a ocs = a::ocs; val cons = fn: 'a > 'a list -> à list Cons 1 [2,3,4]: Compiler creates a cons for into Java Generics public class Stack < T> { new Stack < Integer > (): 2 Compiler fills in the type

Polymorphic Concepts II

- Dynamic polymorphism
 - Run the method in the child
 - Must be done at <u>run-time</u> since that's when we know the child's type
 - Type errors cause run-time faults (crashes!)

```
Student s = new Student();

Person p = (Person)s;

p.dance();

Looks in memory

>> Sees a Shdent

>> mus shdent dance()
```

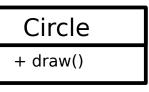
- Compiler looks in memory and finds that the object is really a Student
- So p.dance() runs the dance() action in <u>Student</u>

The Canonical Example I

- A drawing program that can draw circles, squares, ovals and stars
- It would presumably keep a list of all the drawing objects

Option 1

- Keep a list of Circle objects, a list of Square objects,...
- Iterate over each list drawing each object in turn
- What has to change if we want to add a new shape?

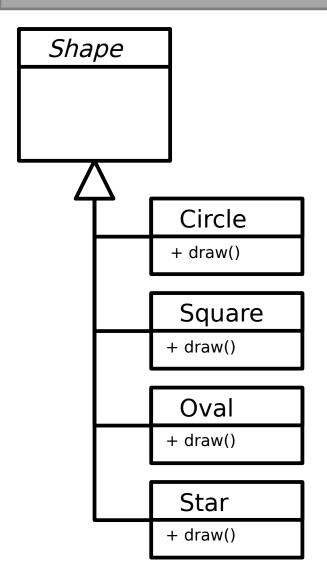


Square + draw()

Oval + draw()

Star + draw()

The Canonical Example II

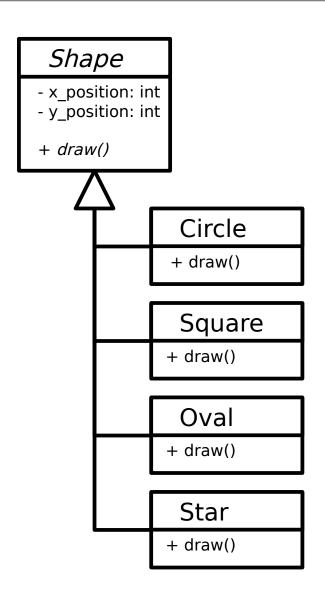


Option 2

- Keep a single list of Shape references
- Figure out what each object really is, narrow the reference and then draw() for every Shape s in myShapeList if (s is really a Circle) Circle c = (Circle)s; c.draw(); else if (s is really a Square) Square sq = (Square)s; sq.draw(); else if...

What if we want to add a new shape?

The Canonical Example III



Option 3 (Polymorphic)

- Keep a single list of Shape references
- Let the compiler figure out what to do with each Shape reference

For every Shape s in myShapeList s.draw();



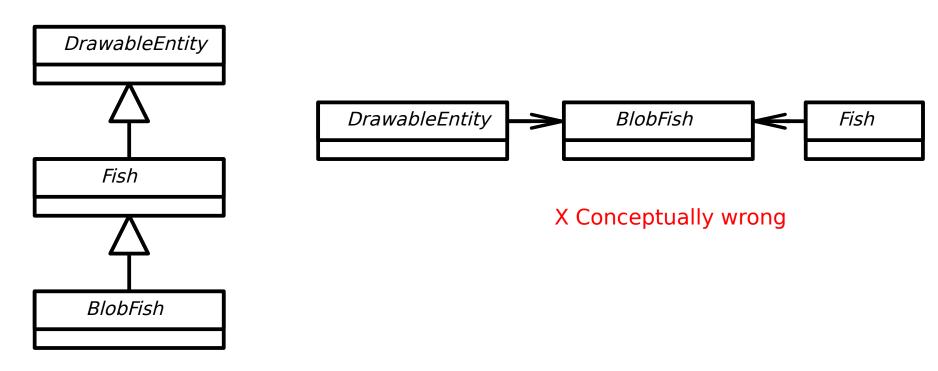
What if we want to add a new shape?

Implementations

- Java
 - All methods are dynamic polymorphic.
- Python
 - All methods are dynamic polymorphic.
- **■** C++
 - Only functions marked *virtual* are dynamic polymorphic
- Polymorphism in OOP is an extremely important concept that you need to make <u>sure</u> you understand...

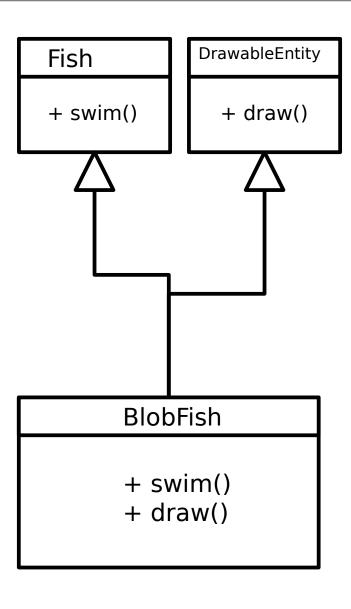
Harder Problems

 Given a class Fish and a class DrawableEntity, how do we make a BlobFish class that is a drawable fish?



X Dependency between two independent concepts

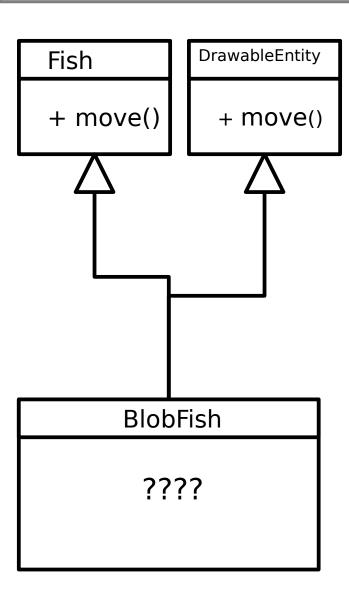
Multiple Inheritance



- If we multiple inherit, we capture the concept we want
- BlobFish inherits from both and is-a Fish and is-a DrawableEntity
- C++:

But...

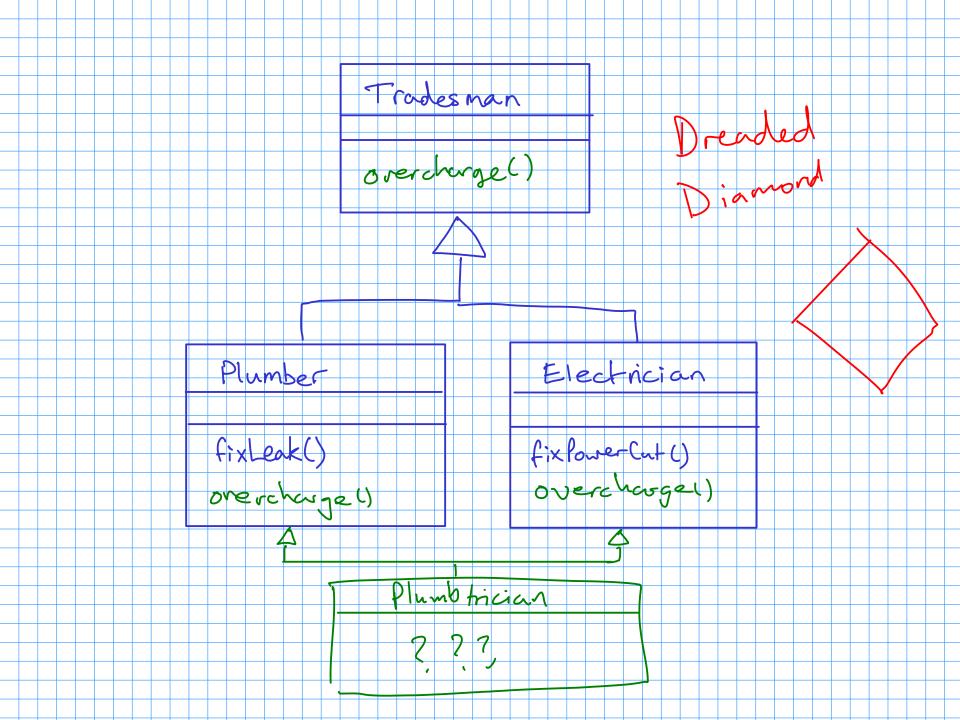
Multiple Inheritance Problems



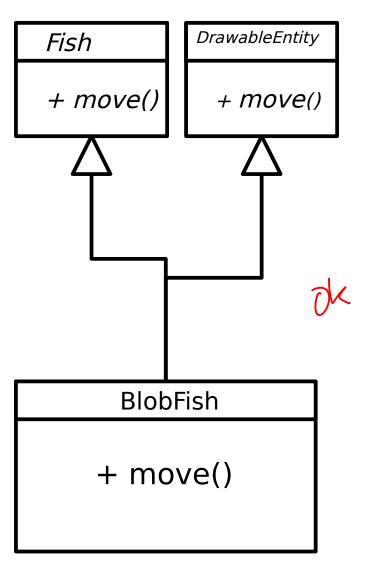
- What happens here? Which of the move() methods is inherited?
- Have to add some grammar to make it explicit
- C++;

```
BlobFish *bf = new BlobFish();
bf->Fish::move();
bf->DrawableEntity::move();
```

Yuk.



Fixing with Abstraction



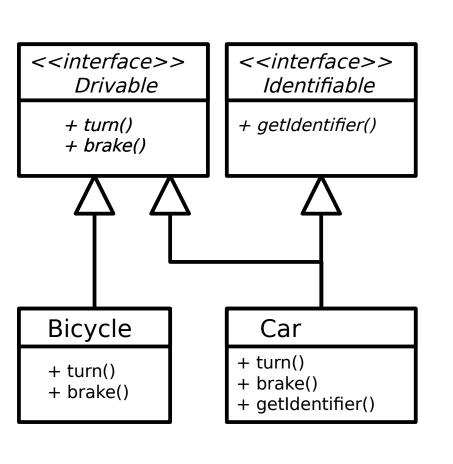
 Actually, this problem goes away if one or more of the conflicting methods is abstract

```
Ok => Mush have a move()
method
```

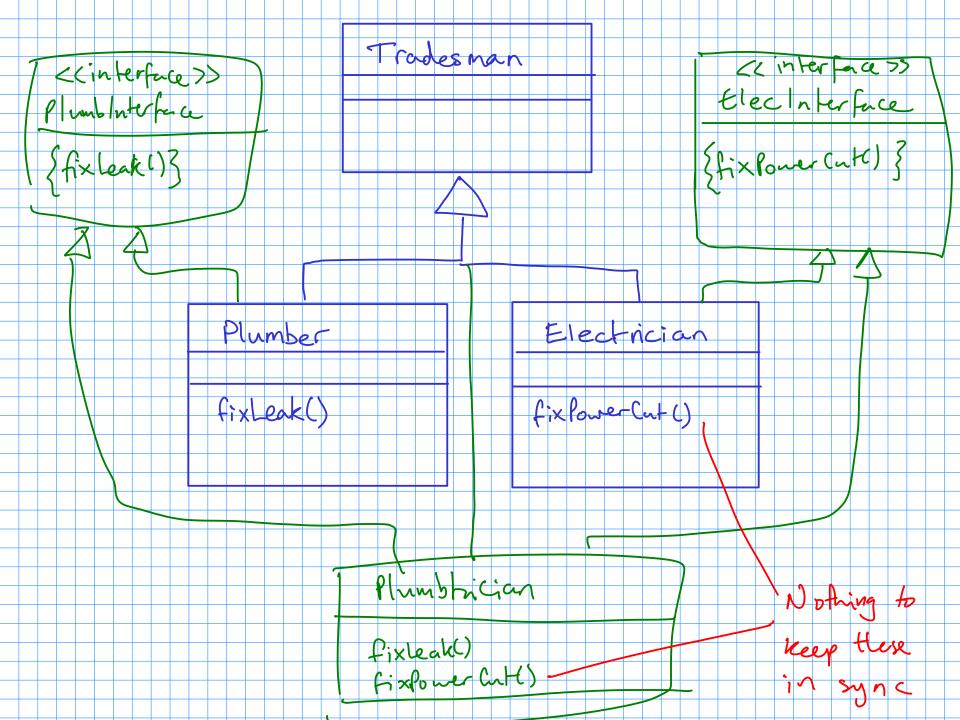
>No définition of move ()

Java's Take on it: Interfaces

- Classes can have at most one parent. Period.
- But special 'classes' that are totally abstract can do multiple inheritance – call these interfaces



```
Interface Drivable { . extend
  public void turn();
  public void brake();
Interface Identifiable { interfaces
  public void getIdentifier();
class Bicycle implements Drivable {
  public void turn() {...}
  public void brake() {... }
class Car implements Drivable, Identifiable {
  public void turn() {...}
  public void brake() {... }
  public void getIdentifier() {...}
```



Aterfaces (Java only) · Only abstract methods Not required to write abstract (implicit) · Only static AND Gnal state public interface car & · Allows a limited (safer) form public void drive(); of multiple inhentance

Lecture 6: Lifecycle of an Object

Constructors

- You will have noticed that the RHS looks rather like a function call, and that's exactly what it is.
- It's a method that gets called when the object is constructed, and it goes by the name of a constructor (it's not rocket science). It maps to the datatype constructors you saw in ML.
- We use constructors to initialise the state of the class in a convenient way
 - A constructor has the same name as the class
 - A constructor has no return type

Constructor Examples

```
Java
public class Person {
 private String mName;
 // Constructor
 public Person(String name) {
    mName=name;
      L No return type
 public static void main(
      String[] args) {
   Person p =
      new Person("Bob");
```

```
C++
class Person {
  private:
   std::string mName;
 public:
  Person(std::string &name) {
      mName=name;
};
int main (int argc,
         char ** argv) {
  Person p ("Bob");
```

Initialiser Blocks

Public class Test \mathcal{E} Private int x = 7;

 $\frac{x=0}{x}$

public Test () {

4

final Kerrsited final int mx= eiter constrictor or the int. block final LinkedList <Integer > mlist = ... Revent re from re-setting the reference but the object remains public final class X E.-3 Cannot extend the class

Default Constructor

```
public class Person {

public ferson() {

private String mName;

public static void main(String[] args) {

Person p = new Person();

}
```

- If you specify no constructor at all, Java fills in an empty one for you
- Here it creates Person() for us
- The default constructor takes no arguments (since it wouldn't know what to do with them!)

Multiple Constructors

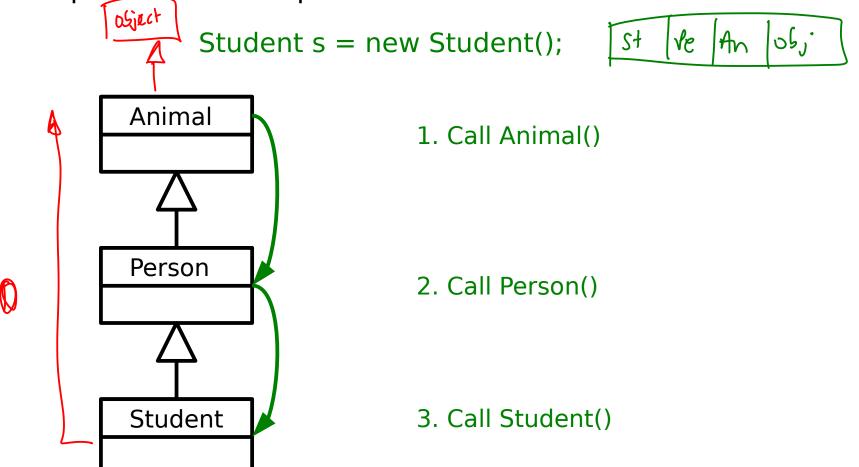
```
public class Student {
  private String mName;
  private int mScore;
  public Student(String s) {
    mName=s;
    mScore=0;
  public Student(String s, int sc) {
     mName=s:
     mScore=sc;
  public static void main(String[] args) {
   Student s1 = new Student("Bob");
   Student s2 = new Student("Bob",55);
```

- You can specify as many constructors as you like.
- Each constructor must have a different signature (argument list)

Calling One Constructor From Another
public Student (String name, int score) {
MNane = nane;
mScore = score;
)
3
public Student (string name) { this (name io);
this (name io);
A mat 1 Ma
3 Pris must be the
first (ine of your
constructor
W 1 3 1 10 10 10 10 10 10 10 10 10 10 10 10 1

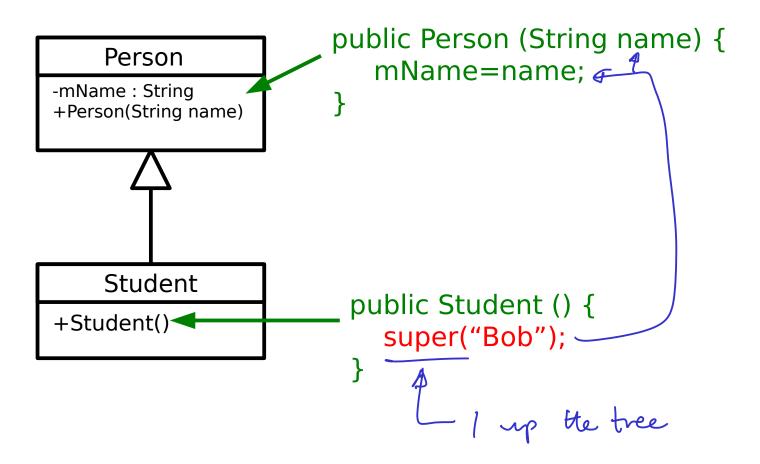
Constructor Chaining

 When you construct an object of a type with parent classes, we call the constructors of all of the parents in sequence



Chaining without Default Constructors

- What if your classes have explicit constructors that take arguments? You need to explicitly chain
- Use super in Java:



Destructors

- Most OO languages have a notion of a destructor too
 - Gets run when the object is destroyed
 - Allows us to release any resources (open files, etc) or memory that we might have created especially for the object

```
class FileReader {
                                                         int main(int argc, char ** argv) {
                   public:
                                                          // Construct a FileReader Object
                                                          FileReader *f = new FileReader():
                    // Constructor
                     FileReader() {
                      f = fopen("myfile","r");
                                                          // Use object here
C++
                                                          ... 🥕
                      Destructor
                                                          // Destruct the object
                      fileReader() {
                                                          delete f;
                       fclose(f);
                   private:
                     FILE *file;
```

Cleaning Up

 A typical program creates lots of objects, not all of which need to stick around all the time

Approach 1:

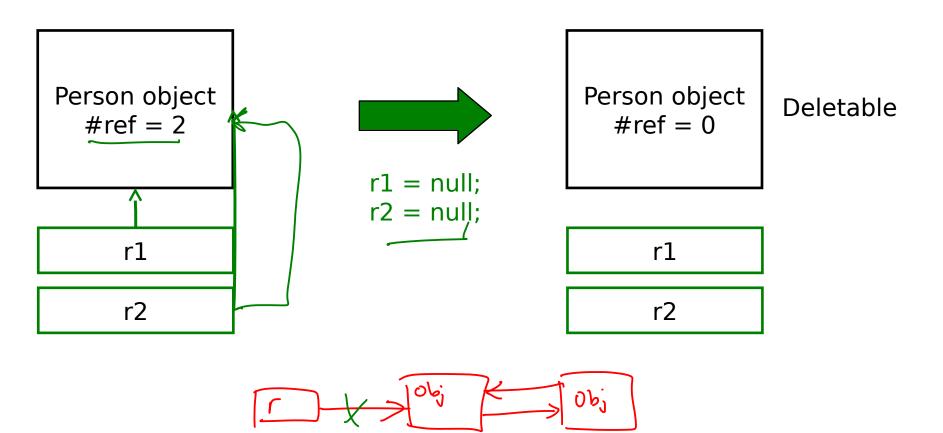
- Allow the programmer to specify when objects should be deleted from memory
- Lots of control, but what if they forget to delete an object?
 - A "memory leak"

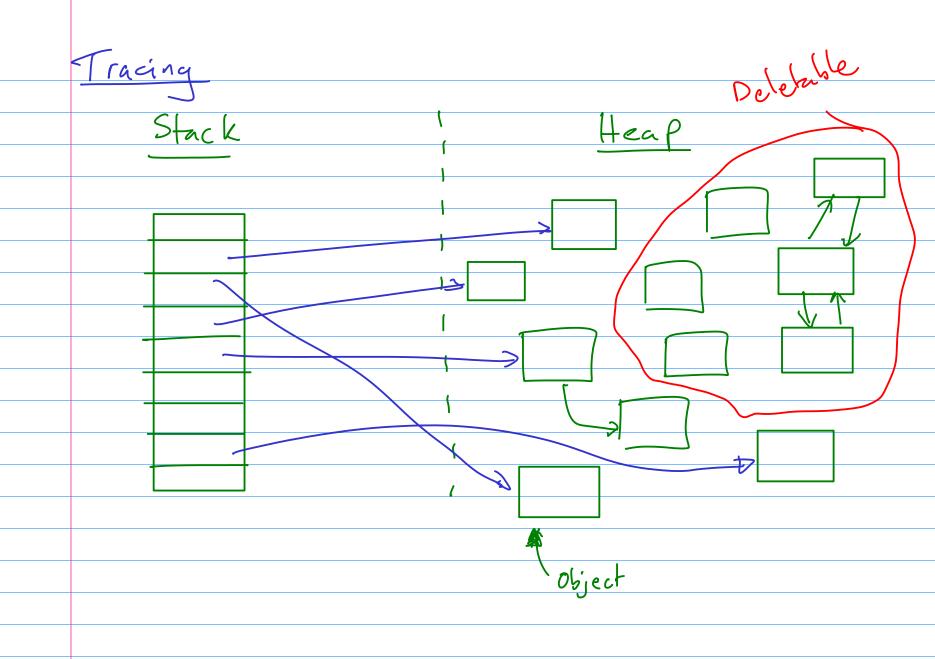
Approach 2:

- Delete the objects automatically (Garbage collection)
- But how do you know when an object will never be used again and can be deleted??

Cleaning Up (Java) I

 Java reference counts. i.e. it keeps track of how many references point to a given object. If there are none, the programmer can't access that object ever again so it can be deleted





Cleaning Up (Java) II

- Actual deletion occurs through a garbage collector
 - A separate process that periodically scans the objects in memory for any with a reference count of zero, which it then deletes.
 - Running the garbage collector is obviously not free. If your program creates a lot of short-term objects, you will soon notice the collector running
 - Gives noticeable pauses to your application while it runs.
 - But minimises memory leaks (it does not prevent them...)

Cleaning Up (Java) III

- One problem with GC is we have no idea when an object will actually be deleted. The GC may even decide to defer the deletion until a future run.
- This causes issues for destructors it might be ages before a resource is closed and available again!
- Therefore Java doesn't have destructors
- It does have finalizers that gets run when the GC deletes an object
 - BUT there's no guarantee an object will <u>ever</u> get garbage collected in Java...
 - Garbage Collection != Destruction

Lecture 7: Error Handling

Return Codes

 The traditional imperative way to handle errors is to return a value that indicates success/failure/error

```
public int divide(double a, double b) {
  if (b==0.0) return -1; // error
  double result = a/b;
  return 0; // success
}
...

if ( divide(x,y)<0) System.out.println("Failure!!");</pre>
```

- Problems:
 - Could ignore the return value
 - Have to keep checking what the return values are meant to signify, etc.
 - The actual result often can't be returned in the same way

Pass by reference int divide (double n, double d, double & result) { if (d== 0.0) return -1; else { result = n/d; return 0; double 1:00 divide (a, b, r)

Could also try to signal an error though magic values double divide (double n, double d) {

if (d = = 0) return Double. MAX VALUE; double d = d; rde (a,b); d = divide (d,c); danble r = 0.0; if (divide(a, b, r) < 0) § 1/ handle evor if (divide (r,c,r) <0) { 1. Ignore the rehm 2 Forget

Deferred Error Handling

- A similar idea (with the same issues) is to set some state in the system that needs to be checked for errors.
- C++ does this for streams:

```
ifstream file( "test.txt" );
if ( file.good() )
{
     cout << "An error occurred opening the file" << endl;
}</pre>
```

```
public Vector 2D pos = null;
Vector2D norm (Vector2D V) [
     if (V.getX() == 0.0 lk v.getY() == 0.0) {

return (null;)
}
     else {
                      or absence
            donble mag = Math. sqrt (v.getx() * v.getx() +
                                   V. getY() * V-ge { Y() )
           return new Vector 21) (V.get X()/mag,
                                   v.get Y()/mag).
```

Non-Exam

Ophional CT>

Exceptions

- An exception is an object that can be thrown or raised by a method when an error occurs and caught or handled by the calling code
- Example usage:

Flow Control During Exceptions

 When an exception is thrown, any code left to run in the try block is skipped

```
\stackrel{\blacktriangle}{\sim} double z=0.0;
boolean failed=false;
   z = \frac{\text{divide}(5,0)}{z = 1.0};
z = 1.0;
This is skipped
  catch(DivideByZeroException d) {
 failed=true;
>z=3.0;
  System.out.println(z+" "+failed);
```

Creating Exceptions

 Just extend Exception (or RuntimeException if you need it to be unchecked). Good form to add a detail message in the constructor but not required.

```
public class DivideByZero extends Exception {}

public class ComputationFailed extends Exception {
   public ComputationFailed(String msg) {
      super(msg);
   }
}
```

 You can also add more data to the exception class to provide more info on what happened (e.g. store the numerator and denominator of a failed division)

Throwing Exceptions

- An exception is an object that has Exception as an ancestor
- So you need to create it (with new) before throwing

```
double divide(double x, double y) throws DivideByZeroException {
  if (y==0.0) throw new DivideByZeroException();
  else return x/y;
}
```

Multiple Handlers

A try block can result in a range of different exceptions.
 We test them in sequence

```
try {
    FileReader fr = new FileReader("somefile");
    Int r = fr.read();
}
catch(FileNoteFound fnf) {
    // handle file not found with FileReader
}
catch(IOException d) {
    // handle read() failed
}
```

Exception Hierarchies

You can use inheritance hierarchies

```
public class MathException extends Exception {...}
public class InfiniteResult extends MathException {...}
public class DivByZero extends MathException {...}
```

And catch parent classes

```
try {
    ...
}
catch(InfiniteResult ir) {
    // handle an infinite result
}
catch(MathException me) {
    // handle any MathException or DivByZero
}
```

Types of Error Defects - Bug - Not handle - able · Invalid Conditions out of our immediate control - These are putentially fixable as we mn

Checked vs Unchecked Exceptions

- Checked: must be handled or passed up.
 - Used for recoverable errors
 - Java requires you to declare checked exceptions that your method throws
 - Java requires you to catch the exception when you call the function

double somefunc() throws SomeException {}

- Unchecked: not expected to be handled. Used for programming errors
 - Extends RuntimeException
 - Good example is NullPointerException

finally

 With resources we often want to ensure that they are closed whatever happens

```
try {
    fr,read();
    fr.close();
}
catch(IOException ioe) {
    // read() failed but we must still close the FileReader
    fr.close();
}
```

finally II

 The finally block is added and will always run (after any handler)

```
try {
   fr,read();
}
catch(IOException ioe) {
   // read() failed
}
finally {
   fr.close();
}
```

Evil I: Exceptions for Flow Control

- At some level, throwing an exception is like a GOTO
- Tempting to exploit this

```
try {
  for (int i=0; i++) {
    System.out.println(myarray[i]);
  }
} catch (ArrayOutOfBoundsException ae) {
  // This is expected
```

- This is not good. Exceptions are for exceptional circumstances only
 - Harder to read
 - May prevent optimisations

Evil II: Blank Handlers

- Checked exceptions must be handled
- Constantly having to use try...catch blocks to do this can be annoying and the temptation is to just gaffer-tape it for now

```
try {
   FileReader fr = new FileReader(filename);
}
catch (FileNotFound fnf) {
}
```

 ...but we never remember to fix it and we could easily be missing serious errors that manifest as bugs later on that are extremely hard to track down Evil III: Circumventing Exception Handling public void fine thous Exception {

Advantages of Exceptions

Advantages:

- Class name can be descriptive (no need to look up error codes)
- Doesn't interrupt the natural flow of the code by requiring constant tests
- The exception object itself can contain state that gives lots of detail on the error that caused the exception
- Can't be ignored, only handled

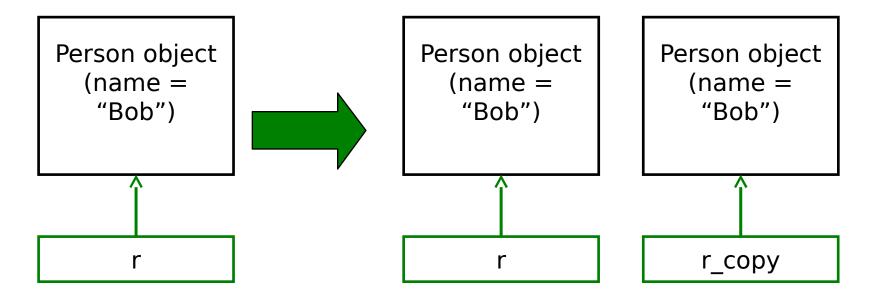


1. Separate Error Handling and Regular Code Exceptions Return codes code Code Code code

Lecture 8: Copying Objects

Cloning I

Sometimes we really do want to copy an object



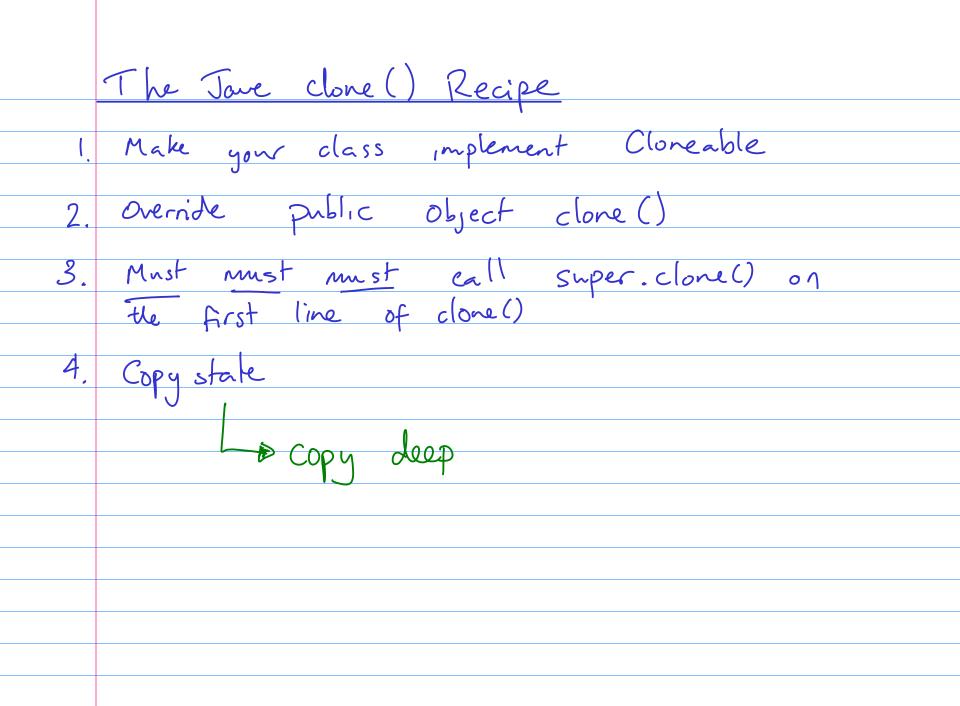
- Java calls this cloning
- We need special support for it

	Copy Constructors
•	
1,	Greate a constructor that takes a ref. to
	Greate a constructor that takes a ref. to an object to be copied.
2.	You call the copy cons. of the parent (using
	You call the copy cons. of the parent (using super)
3	Copy the state in your class
	,
	Copy shallow
	<u> </u>
	Lo Copy deep

Cloning II

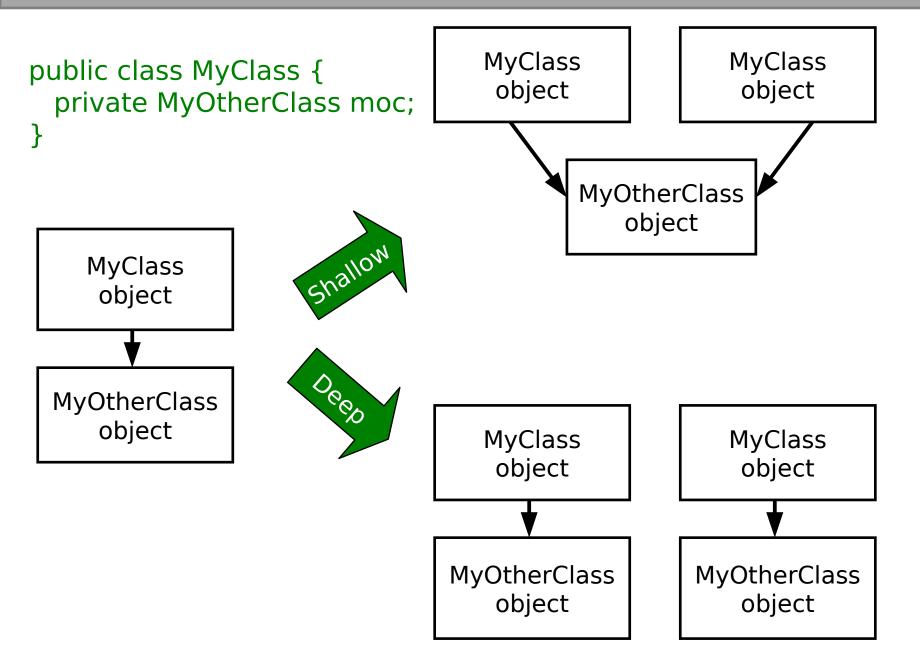
- Every class in Java ultimately inherits from the
 Object class
 - This class contains a clone() method so we just call this to clone an object, right?
 - This can go horribly wrong if our object contains reference types (objects, arrays, etc)

The Jave clone () Recipe
(Implement Cloneable
2. Overide public object clone();
3. Call super. clone() on the first line
4. Deep copy ony referenced objects as appropriate



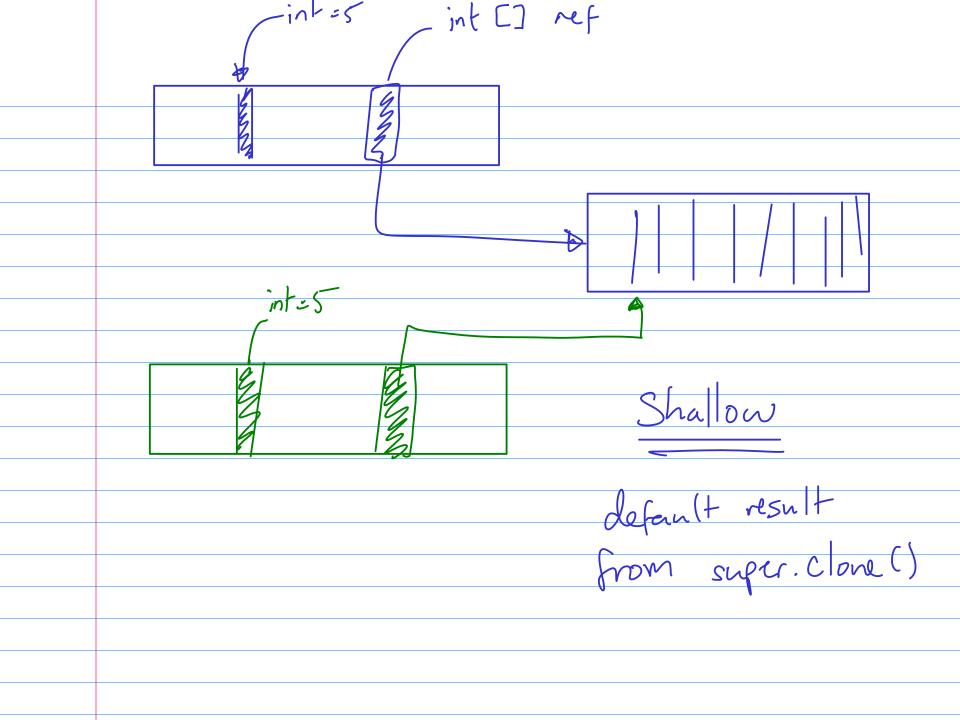
The clone() logic if this chass is not Cloneable thow Clone Not Supported Exception Do a Sit-for-bit copy int=5 literal memory copy

Shallow and Deep Copies



Java Cloning

- So do you want shallow or deep?
 - The default implementation of clone() performs a shallow copy
 - But Java developers were worried that this might not be appropriate: they decided they wanted to know for <u>sure</u> that we'd thought about whether this was appropriate
- Java has a Cloneable interface
 - If you call clone on anything that doesn't extend this interface, it fails



Clone Example I

```
public class Velocity {
  public float vx;
  public float vy;
  public Velocity(float x, float y) {
    VX=X;
    \vee \vee = \vee;
public class Vehicle {
  private int age;
  private Velocity vel;
  public Vehicle(int a, float vx, float vy) {
    age=a;
    vel = new Velocity(vx,vy);
```

Clone Example II

```
public class Vehicle implements Cloneable {
 private int age;
 private Velocity vel;
 public Vehicle(int a, float vx, float vy) {
    age=a;
    vel = new Velocity(vx,vy);
 }
 public Object clone() {
    return super.clone();
};
```

Clone Example III

```
public class Velocity implement Cloneable {
   public Object clone() {
      return super.clone();
};
public class Vehicle implements Cloneable {
 private int age;
 private Velocity v;
 public Student(int a, float vx, float vy) {
    age=a;
   vel = new Velocity(vx,vy);
 }
 public Object clone() {
   Vehicle cloned = (Vehicle) super.clone();
    cloned.vel = (Velocity)vel.clone();
    return cloned;
```

Cloning Arrays

 Arrays have build in cloning but the contents are only cloned shallowly

```
int intarray[] = new int[100];
Vector3D vecarray = new Vector3D[10];
...
int intarray2[] = intarray.clone();
Vector3D vecarray2 = vecarray.clone();
```

Covariant Return Types

The need to cast the clone return is annoying

```
public Object clone() {
    Vehicle cloned = (Vehicle) super.clone();
    cloned.vel = (Velocity)vel.clone();
    return cloned;
}
```

 Recent versions of Java allow you to override a method in a subclass and change its return type to a subclass of the original's class

```
class C {
    A mymethod() {}
}
class B extends A {}

class D extends C {
    B mymethod() {}
}
```

Marker Interfaces

- If you look at what's in the Cloneable interface, you'll find it's empty!! What's going on?
- Well, the clone() method is already inherited from Object so it doesn't need to specify it
- This is an example of a Marker Interface
 - A marker interface is an empty interface that is used to label classes
 - This approach is found occasionally in the Java libraries

Copy Constructors

- Another way to create copies of objects is to define a copy constructor that takes in an object of the same type and manually copies the data
- See examples sheet

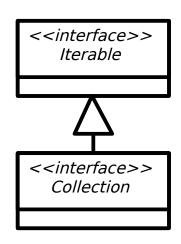
```
public class Vehicle {
  private int age;
  private Velocity vel;
  public Vehicle(int a, float vx, float vy) {
    age=a;
    vel = new Velocity(vx,vy);
  public Vehicle(Vehicle v) {
       age=v.age;
       vel = v.vel.clone();
```

Lecture 8: Java Collections

Java Class Library

- Java the platform contains around 4,000 classes/interfaces
 - Data Structures
 - Networking, Files
 - Graphical User Interfaces
 - Security and Encryption
 - Image Processing
 - Multimedia authoring/playback
 - And more...
- All neatly(ish) arranged into packages (see API docs)

Java's Collections Framework

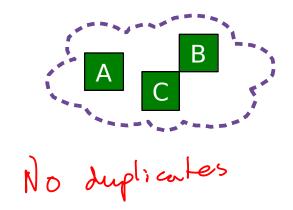


- Important chunk of the class library
- A collection is some sort of grouping of things (objects)
- Usually when we have some grouping we want to go through it ("iterate over it")
- The Collections framework has two main interfaces: Iterable and Collections. They define a set of operations that all classes in the Collections framework support
- add(Object o), clear(), isEmpty(), etc.

Sets

<<interface>> Set

- A collection of elements with no duplicates that represents the mathematical notion of a set
- <u>Tree</u>Set: objects stored in order
- HashSet: objects in unpredictable order but fast to operate on (see Algorithms course)



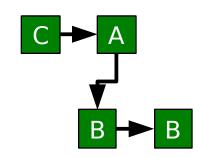
```
Sorted
```

```
TreeSet<Integer> ts = new TreeSet<Integer>();
ts.add(15);
ts.add(12);
ts.contains(7); // false
ts.contains(12); // true
ts.first(); // 12 (sorted)
```

Lists

<<interface>> List

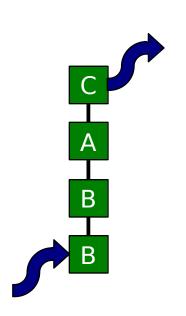
- An ordered collection of elements that may contain duplicates
- LinkedList: linked list of elements
- ArrayList: array of elements (efficient access)
- Vector: Legacy, as ArrayList but threadsafe



Queues

<<interface>> Queue

- An ordered collection of elements that may contain duplicates and supports removal of elements from the head of the queue
- offer() to add to the back and poll() to take from the front
- LinkedList: supports the necessary functionality
- PriorityQueue: adds a notion of priority to the queue so more important stuff bubbles to the top



```
LinkedList<Double> || = new LinkedList<Double>(); ||.offer(1.0); ||.offer(0.5); ||.poll(); // 1.0 ||.poll(); // 0.5
```

Maps

<<interface>> Map

- Like dictionaries in ML
- Maps key objects to value objects
- Keys must be unique
- Values can be duplicated and (sometimes) null.
- TreeMap: keys kept in order
- HashMap: Keys not in order, efficient (see Algorithms)

```
K1 A B B
```

```
TreeMap<String, Integer> tm = new TreeMap<String,Integer>();
tm.put("A",1);
tm.put("B",2);
tm.get("A"); // returns 1
tm.get("C"); // returns null
tm.contains("G"); // false
```

Iteration

for loop

```
LinkedList<Integer> list = new LinkedList<Integer>();
...
for (int i=0; i<list.size(); i++) {
    Integer next = list.get(i);
}</pre>
```

foreach loop (Java 5.0+)

Iterators

What if our loop changes the structure?

```
for (int i=0; i<list.size(); i++) {
    If (i==3) list.remove(i);
}</pre>
```

Java introduced the Iterator class

```
Iterator<Integer> it = list.iterator();
while(it.hasNext()) {Integer i = it.next();}
for (; it.hasNext(); ) {Integer i = it.next();}
```

Safe to modify structure

```
while(it.hasNext()) {
   it.remove();
}
```

The Origins of Generics

```
// Make a TreeSet object
TreeSet ts = new TreeSet();
// Add integers to it
ts.add(new Integer(3));
// Loop through
iterator it = ts.iterator();
while(it.hasNext()) {
    Object o = it.next();
    Integer i = (Integer)o;
```

- The original Collections framework just dealt with collections of Objects
 - Everything in Java "is-a"
 Object so that way our
 collections framework will
 apply to any class
 - But this leads to:
 - Constant casting of the result (ugly)
 - The need to know what the return type is
 - Accidental mixing of types in the collection

The Origins of Generics II

```
// Make a TreeSet object
TreeSet ts = new TreeSet();
// Add integers to it
ts.add(new Integer(3));
ts.add(new Person("Bob"));
// Loop through
iterator it = ts.iterator();
                                      Going to fail for the
while(it.hasNext()) {
                                      second element!
    Object o = it.next();
                                      (But it will compile:
    Integer i = (Integer)o;
                                      the error will be at
                                      runtime)
```

The Generics Solution

- Java implements type erasure
 - Compiler checks through your code to make sure you only used a single type with a given Generics object
 - Then it deletes all knowledge of the parameter, converting it to the old code invisibly

The C++ Templates Solution

 Compiler first generates the class definitions from the template

};

class MyClass_float {

```
class MyClass<T> {
  T membervar;
};

class MyClass_int {
  int membervar;
};

class MyClass_double {
  double membervar;
}
```

Generics and SubTyping

```
// Object casting
Animal
               Person p = new Person();
               Animal o = (Animal) p;
               // List casting
               List<Person> plist = new LinkedList<Person>();
Person
               List<Animal> alist = (List<Animal>)plist;
               So a list of Persons is a list of Animals, yes?
      alist, add (new Hippol));
```

```
import java.util.LinkedList;

public class Naughty {

    public static void main(String[] args) {
        LinkedList temp = new LinkedList<Double>();
        temp.add(9.0);
        printAll(temp); // Should fail
    }

    public static void printAll(LinkedList<Object> list) {
        for (Object d : list)
            System.out.println((Double)d);
    }
}
```

Lecture 10: Comparing Objects

Comparing Primitives

```
> Greater Than
>= Greater than or equal to
== Equal to
!= Not equal to
< Less than
<= Less than or equal to</pre>
```

- Clearly compare the value of a primitive
- But what does (ref1==ref2) do??
 - Test whether they point to the same object?
 - Test whether the objects they point to have the same state?

Reference Equality

- r1 = -r2, r1! = r2
- These test reference equality
- i.e. do the two references point of the same chunk of memory?

```
Person p1 = new Person("Bob");
Person p2 = new Person("Bob");

(p1==p2);

(p1!=p2);

True (references differ)

(p1==p1);

True
```

Value Equality

- Use the equals() method in Object
- Default implementation just uses reference equality (==) so we have to override the method

```
public EqualsTest {
  public int x = 8;
  @Override
  public boolean equals(Object o) {
     EqualsTest e = (EqualsTest)o;
     return (this.x = = e.x);
                    Bocault
  public static void main(String args[]) {
     EqualsTest t1 = new EqualsTest();
     EqualsTest t2 = new EqualsTest();
     System.out.println(t1==t2);
    System.out.println(t1.equals(t2));
```

Aside: Use The Override Annotation

It's so easy to mistakenly write:

```
public EqualsTest {
  public int x = 8;
  public boolean equals(EqualsTest e) {
     return (this.x = = e.x);
  public static void main(String args[]) {
     EqualsTest t1 = new EqualsTest();
     EqualsTest t2 = new EqualsTest();
     Object o1 = (Object) t1;
     Object o2 = (Object) t2;
     System.out.println(t1.equals(t2));
     System.out.println(o1.equals(o2));
```

Aside: Use The Override Annotation

Annotation would have picked up the mistake:

```
public EqualsTest {
  public int x = 8;
  @Override
  public boolean equals(EqualsTest e) {
     return (this.x = =e.x);
  public static void main(String args[]) {
     EqualsTest t1 = new EqualsTest();
     EqualsTest t2 = new EqualsTest();
     Object o1 = (Object) t1;
     Object o2 = (Object) t2;
     System.out.println(t1.equals(t2));
     System.out.println(o1.equals(o2));
```

Java Quirk: hashCode()

- Object also gives classes hashCode()
- Code assumes that if equals(a,b) returns true, then a.hashCode() is the same as b.hashCode()
- So you should override hashCode() at the same time as equals()

Comparable < T > Interface I

int compareTo(T obj);

- Part of the Collections Framework
- Doesn't just tell us true or false, but smaller, same, or larger: useful for sorting.
- Returns an integer, r:
 - r<0 This object is less than obj
 - r==0 This object is equal to obj
 - r>0 This object is greater than obj

Comparable<T> Interface II

```
public class Point implements Comparable < Point > {
   private final int mX;
   private final int mY;
   public Point (int, int y) { mX=x; mY=y; }
   // sort by y, then x
   public int compareTo(Point p) {
      if ( mY>p.mY) return 1;
      else if (mY<p.mY) return -1;
      else {
        if (mX>p.mX) return 1;
        else if (mX<p.mX) return -1;
        else return 0.
// This will be sorted automatically by y, then x
Set<Point> list = new TreeSet<Point>();
```

Comparator<T> Interface I

int compare(T obj1, T obj2)

- Also part of the Collections framework and allows us to specify a specific ordering for a particular job
- E.g. a Person might have natural ordering that sorts by surname. A Comparator could be written to sort by age instead...

Comparator<T> Interface II

```
public class Person implements Comparable<Person> {
   private String mSurname;
   private int mAge;
   public int compareTo(Person p) {
      return mSurname.compareTo(p.mSurname);
public class AgeComparator implements Comparator<Person> {
  public int compare(Person p1, Person p2) {
    return (p1.mAge-p2.mAge);
ArrayList<Person> plist = ...;
Collections.sort(plist); // sorts by surname
Collections.sort(plist, new AgeComparator()); // sorts by age
```

Operator Overloading

 Some languages have a neat feature that allows you to overload the comparison operators. e.g. in C++

```
class Person {
  public:
    Int mAge
    bool operator==(Person &p) {
      return (p.mAge==mAge);
    };
}

Person a, b;
b == a; // Test value equality
```

Lecture 11: Design Patterns

Design Patterns

- A Design Pattern is a general reusable solution to a commonly occurring problem in software design
- Coined by Erich Gamma in his 1991 Ph.D. thesis
- Originally 23 patterns, now many more. Useful to look at because they illustrate some of the power of OOP (and also some of the pitfalls)
- We will only consider a subset

UML Revisited - private + public # protected Classname state methods B is-a A B has-a A dosthing() Q -pseudocode mlist addi) &

The Open-Closed Principle

Classes should be open for extension but closed for modification

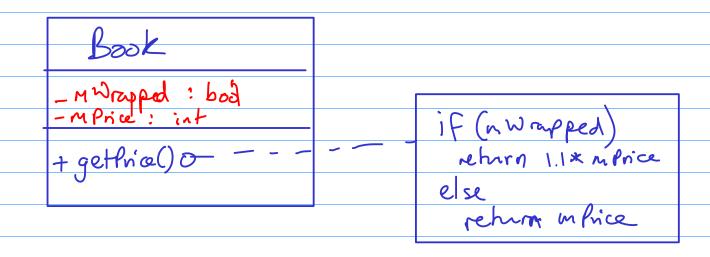
- i.e. we would like to be able to modify the behaviour without touching its source code
- This rule-of-thumb leads to more reliable large software and will help us to evaluate the various design patterns

Decorator

Abstract problem: How can we add state or methods at runtime?

Example problem: How can we efficiently support gift-wrapped books in an online bookstore?

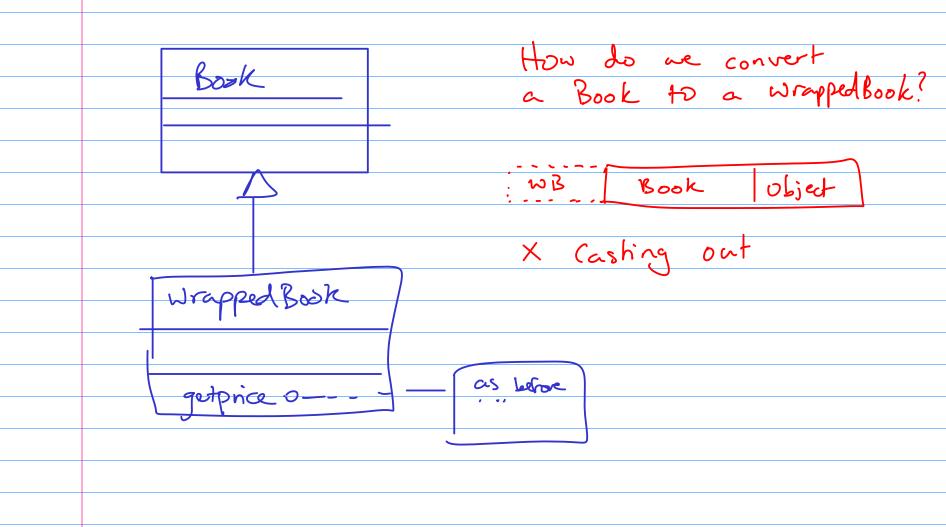
Solution 1: Add variables to the established Book class that describe whether or not the product is to be gift wrapped.



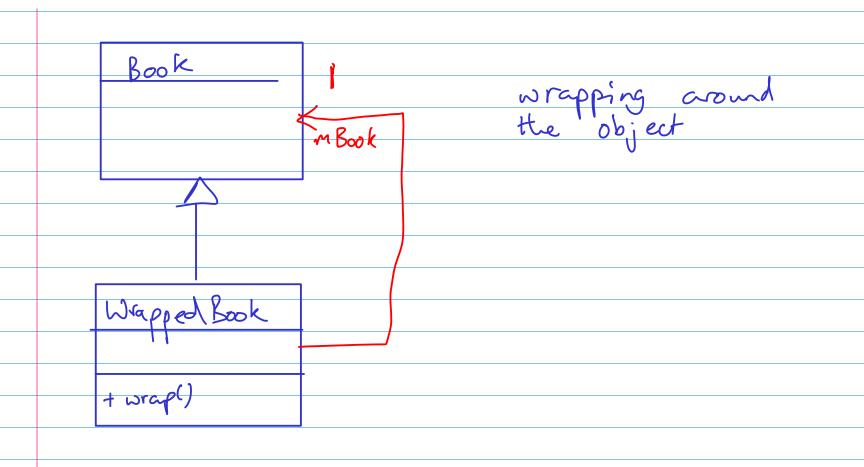
X violates open-closed

X Wastefil - every book wastes a boolean

Solution 2: Extend Book to create WrappedBook.

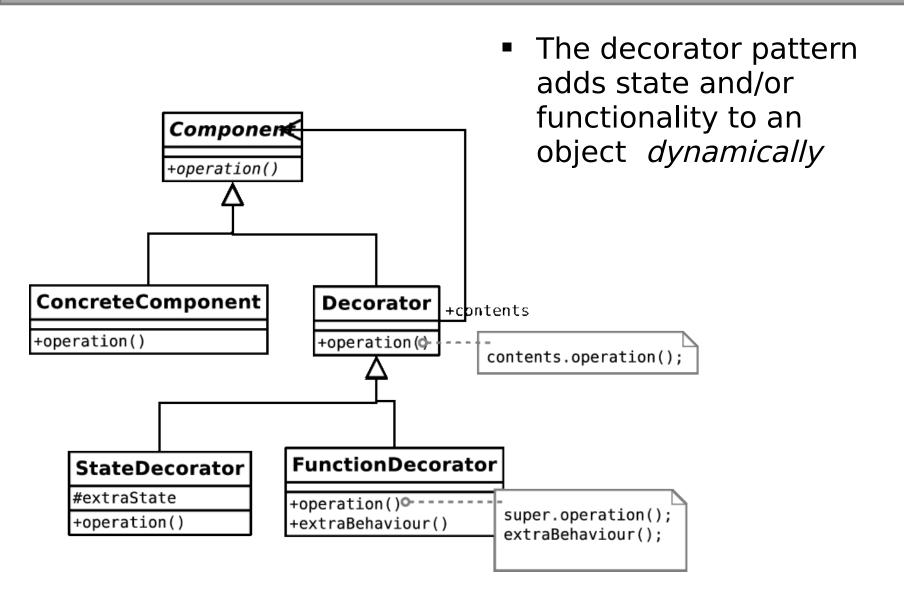


Solution 3: (Decorator) Extend Book to create WrappedBook and also add a member reference to a Book object. Just pass through any method calls to the internal reference, intercepting any that are to do with shipping or price to account for the extra wrapping behaviour.



Deorahir class Buffered Reader Buffered leader br = new Buffered Reader ("myfile.+xt"); FileReader follows instructions exactly + gets things byte by byte. Visit the disk for every byte BR gives you a lookahead buffer => more efficiently 2 Extra furctionality => readline()

Decorator in General



Singleton

Abstract problem: How can we ensure only one instance of an object is created by developers using our code?

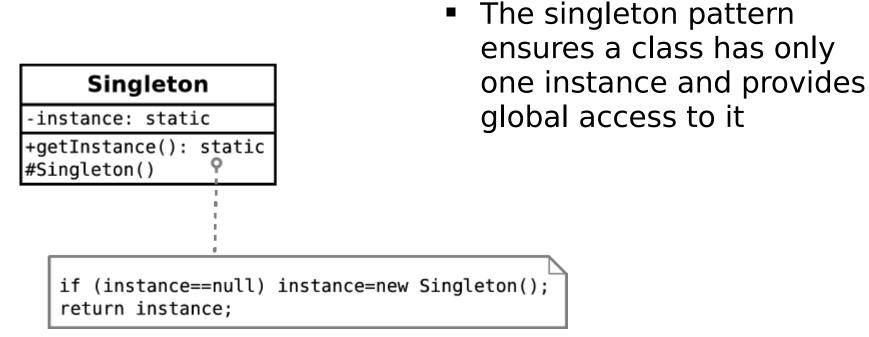
Example problem: You have a class that encapsulates accessing a database over a network. When instantiated, the object will create a connection and send the query. Unfortunately you are only allowed one connection at a time.

Singleton

- 1. Make the constructor private
- 2. Create a single static instance
- 3. Make a <u>static</u> getter for your instance

java.ubl. Runtime Rustine, get Runtimel) freeMemony ()

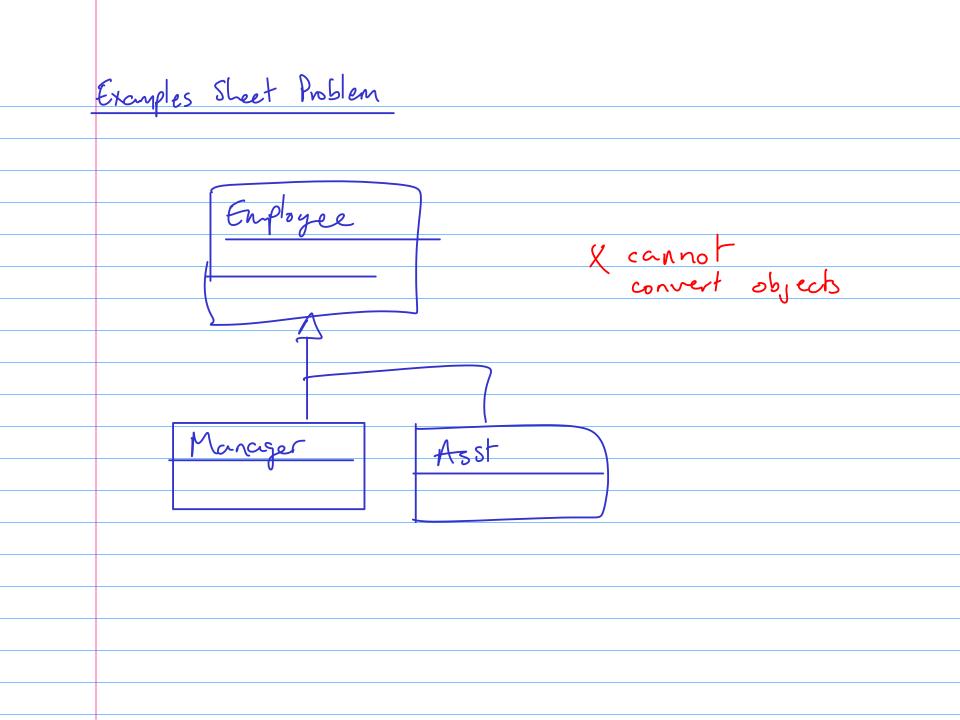
Singleton in General



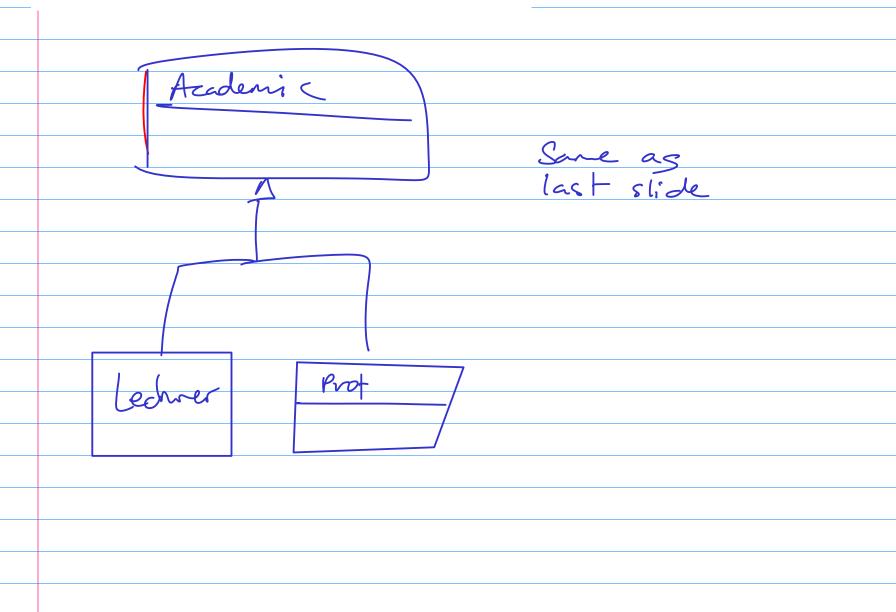
State

Abstract problem: How can we let an object alter its behaviour when its internal state changes?

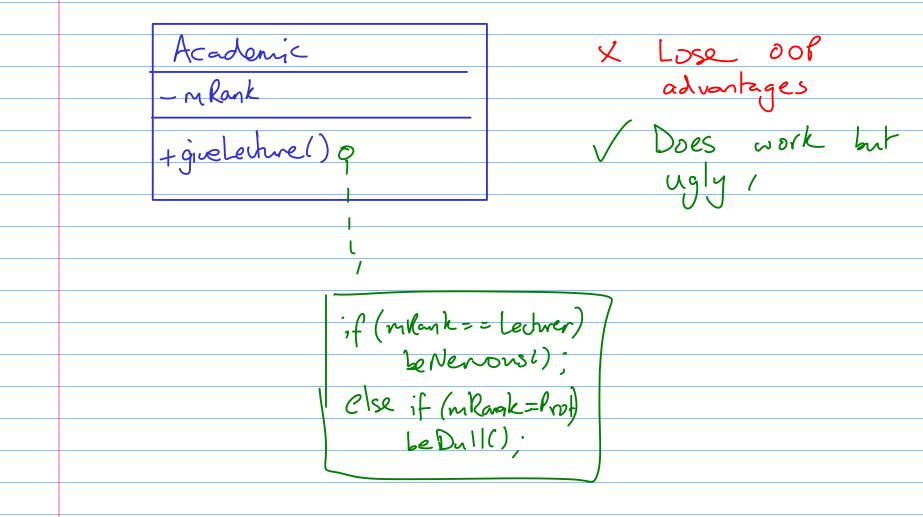
Example problem: Representing academics as they progress through the rank



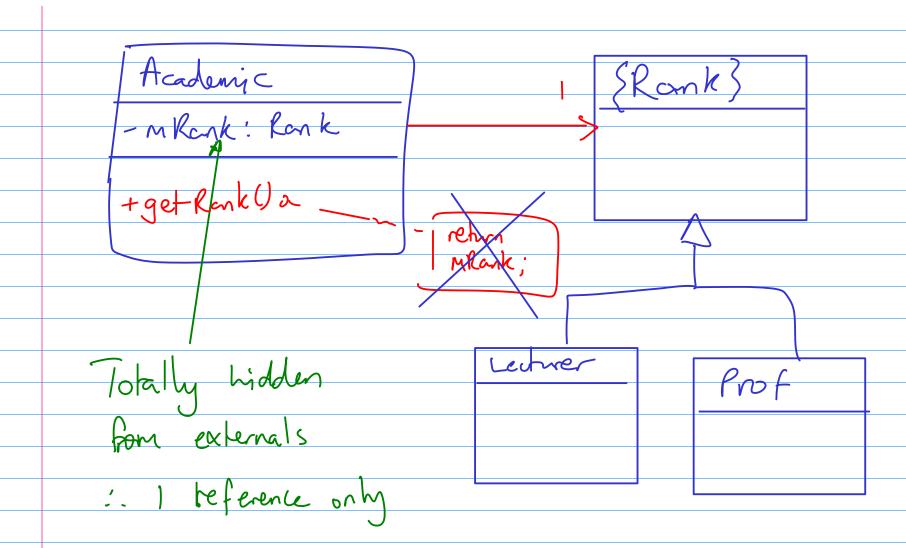
Solution 1: Have an abstract Academic class which acts as a base class for Lecturer, Professor, etc.



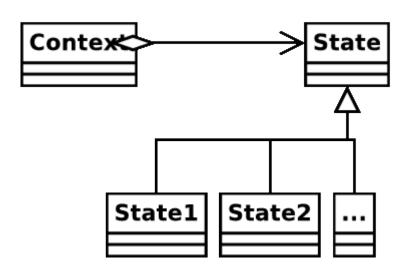
Solution 2: Make Academic a concrete class with a member variable that indicates rank. To get rank-specific behaviour, check this variable within the relevant methods.



Solution 3: (State) Make Academic a concrete class that has-a AcademicRank as a member. Use AcademicRank as a base for Lecturer, Professor, etc., implementing the rank-specific behaviour in each..



State in General



 The state pattern allows an object to cleanly alter its behaviour when internal state changes

Strategy

Abstract problem: How can we select an algorithm implementation at runtime?

Example problem: We have many possible change-making implementations. How do we cleanly change between them?

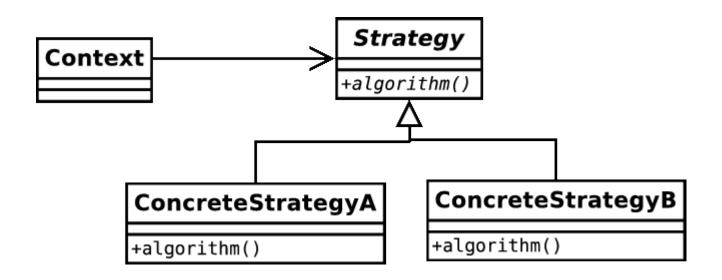
Solution 1: Use a lot of ifelse statements in the getChange() method.

Solution 2: (Strategy) Create an abstract ChangeFinder class. Derive a new class for each of our algorithms.

java. util. Comparator

Strategy in General

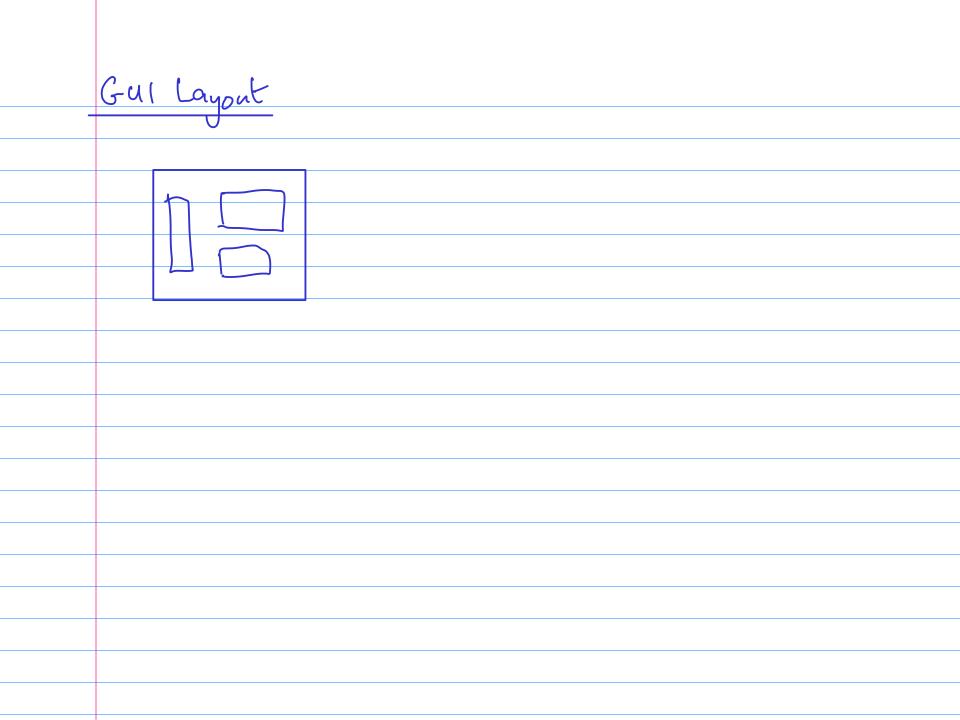
 The strategy pattern allows us to cleanly interchange between algorithm implementations



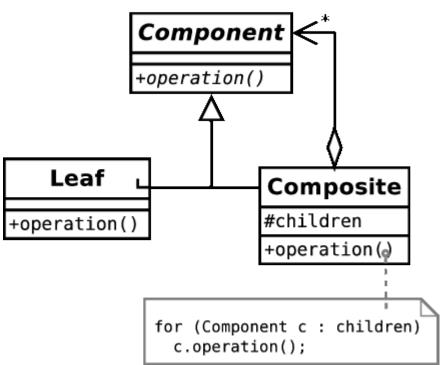
Composite

Abstract problem: How can we treat a group of objects as a single object?

Example problem: Representing a DVD box-set as well as the individual films without duplicating info and with a 10% discount



Composite in General



 The composite pattern lets us treat objects and groups of objects uniformly

Observer

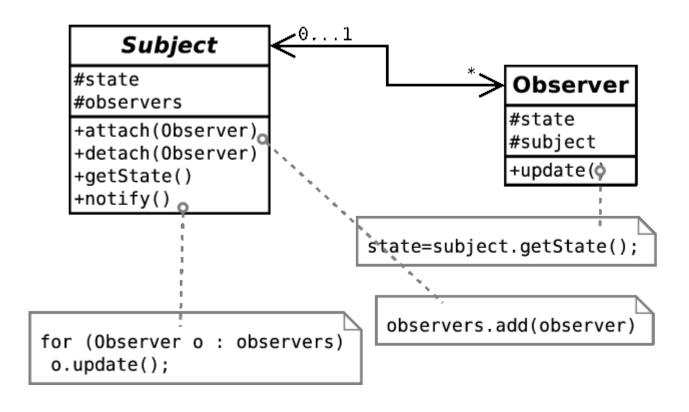
Abstract problem: When an object changes state, how can any interested parties know?

Example problem: How can we write phone apps that react to accelerator events?

	Observer Pattern
_	

Observer in General

 The observer pattern allows an object to have multiple dependents and propagates updates to the dependents automatically.



Factory Pattern
<u> </u>

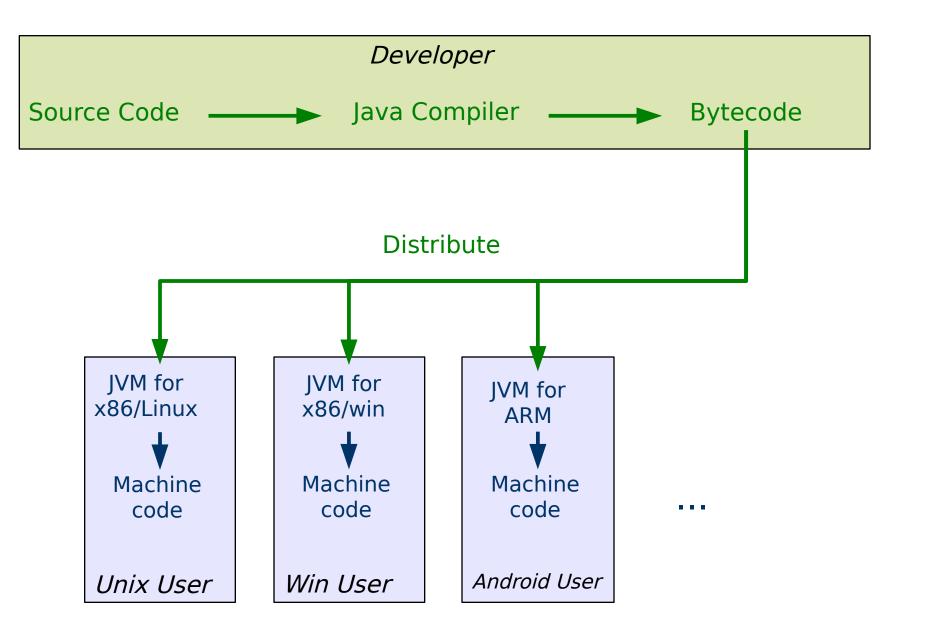
Interpreter to Virtual Machine

- Java was born in an era of internet connectivity. SUN wanted to distribute programs to internet machines
 - But many architectures were attached to the internet
 how do you write one program for them all?
 - And how do you keep the size of the program small (for quick download)?
- Could use an interpreter (→ Javascript). But:
 - High level languages not very space-efficient
 - The source code would implicitly be there for anyone to see, which hinders commercial viability.
- Went for a clever hybrid interpreter/compiler

Java Bytecode I

- SUN envisaged a hypothetical Java Virtual Machine (JVM). Java is compiled into machine code (called bytecode) for that (imaginary) machine. The bytecode is then distributed.
- To use the bytecode, the user must have a JVM that has been specially compiled for their architecture.
- The JVM takes in bytecode and spits out the correct machine code for the local computer. i.e. is a bytecode interpreter

Java Bytecode II



Java Bytecode III

- + Bytecode is compiled so not easy to reverse engineer
- + The JVM ships with tons of libraries which makes the bytecode you distribute small
- + The toughest part of the compile (from human-readable to computer readable) is done by the compiler, leaving the computer-readable bytecode to be translated by the JVM (→ easier job → faster job)
- Still a performance hit compared to fully compiled ("native") code