

L28: Advanced functional programming

Exercise 1

Due on February 9th

For these questions, you may assume that all the definitions given in Figure 1 are available in System $F\omega$.

1. Give a typing derivation for $\Lambda\alpha :: *. \lambda x : \alpha. \langle x, x \rangle$ (3 marks)
2. Algorithm J is defined recursively over the structure of terms. The case for function application (M N) is as follows:

$$\begin{aligned} J \ (\Gamma, M \ N) = \beta & \\ \text{where } A = J \ (\Gamma, M) & \\ \text{and } B = J \ (\Gamma, N) & \\ \text{and unify ' } (\{A = B \rightarrow \beta\}) \text{ succeeds} & \\ \text{and } \beta \text{ is fresh} & \end{aligned}$$

Give similar cases to handle the following constructs:

- (i) Constructing a pair ($\langle M, N \rangle$)
- (ii) Projecting the first element of a pair (**fst** M)
- (iii) Constructing a sum using **inl** (**inl** M)
- (iv) Destructing a sum (**case** L **of** x.M | y.N)

(5 marks)

3. The following OCaml type represents ternary trees:

```
type 'a ttree =
| Empty : 'a ttree
| Tree : 'a * 'a ttree * 'a ttree * 'a ttree -> 'a ttree
```

- (i) Give an encoding of this type in System $F\omega$ consisting of a type operator (**TTree**) of kind $*\Rightarrow*$, a function for each constructor (**empty** and **tree**), and a function (**foldTTree**) for folding over trees.

(3 marks)

- (ii) Write a function

```
totalNatTree : TTree Nat -> Nat
```

that computes the total of a ternary tree of Nats in System $F\omega$.

(1 marks)

4. Using existentials, products, the List type, the Nat type and the Option type, implement a queue data structure in System F ω with a type corresponding to the following OCaml signature:

```
type 'a t
val empty : 'a t
val enqueue : 'a -> 'a t -> 'a t
val dequeue : 'a t -> 'a t * 'a option
val size : 'a t -> int
```

(6 marks)

5. (i) Why does the type checker reject the following program? Specifically, what potentially unsafe behaviour is the type checker's rejection of the program intended to prevent?

```
let rec make size elem =
  if size <= 0 then []
  else elem :: (make (size - 1) elem)
```

```
let foo =
  let make_three = make 3 in
    make_three "hello",
    make_three 5
```

(3 marks)

- (ii) Replace `make` with a function of the same type which would make the program unsafe if the type checker did not reject it.

(1 mark)

- (iii) Adjust the definition of `make_three` so that the program is accepted by the type-checker, without changing the result of the program.

(1 mark)

- (iv) The following program is also rejected by the type-checker:

```
let make size =
  Printf.printf "Making lists of length %d\n" size;
let rec loop size elem =
  if size <= 0 then []
  else elem :: (loop (size - 1) elem)
in
  loop size
```

```
let foo =
  let make_three = make 3 in
    make_three "hello",
    make_three 5
```

It is not possible to make this program pass the type-checker by only changing `make_three` without affecting the program's observable behaviour (i.e. the number of times the message is printed).

However, using a record field with a universal type, it is possible to wrap the result of `make` so that it can be used polymorphically.

Using this technique, adjust the program so that it passes the type-checker without affecting its observable behaviour.

(2 marks)

$$\begin{aligned}
\text{Nat} &= \forall \alpha :: *. \alpha \rightarrow (\alpha \rightarrow \alpha) \rightarrow \alpha \\
\text{zero} &= \Lambda \alpha :: *. \lambda z : \alpha. \lambda s : \alpha \rightarrow \alpha. z \\
\text{succ} &= \lambda n : \text{Nat}. \Lambda \alpha :: *. \lambda z : \alpha. \lambda s : \alpha \rightarrow \alpha. s (n [\alpha] z s) \\
\text{add} &= \lambda m : \text{Nat}. \lambda n : \text{Nat}. m [\text{Nat}] n \text{succ} \\
\text{List} &= \lambda \alpha :: *. \forall \varphi :: * \Rightarrow *. \varphi \alpha \rightarrow (\alpha \rightarrow \varphi \alpha \rightarrow \varphi \alpha) \rightarrow \varphi \alpha \\
\text{nil} &= \Lambda \alpha :: *. \Lambda \varphi :: * \Rightarrow *. \lambda \text{nil} : \varphi \alpha. \lambda \text{cons} : \alpha \rightarrow \varphi \alpha \rightarrow \varphi \alpha. \text{nil} \\
\text{cons} &= \Lambda \alpha :: *. \lambda x : \alpha. \lambda xs : \text{List } \alpha. \\
&\quad \Lambda \varphi :: * \Rightarrow *. \lambda \text{nil} : \varphi \alpha. \lambda \text{cons} : \alpha \rightarrow \varphi \alpha \rightarrow \varphi \alpha. \\
&\quad \text{cons } x (xs [\varphi] \text{nil } \text{cons}) \\
\text{foldList} &= \Lambda \alpha :: *. \Lambda \beta :: *. \\
&\quad \lambda c : \alpha \rightarrow \beta \rightarrow \beta. \lambda n : \beta. \lambda l : \text{List } \alpha. \\
&\quad l [\lambda \gamma :: *. \beta] n c \\
\text{append} &= \Lambda \alpha :: *. \\
&\quad \lambda l : \text{List } \alpha. \lambda r : \text{List } \alpha. \\
&\quad \text{foldList } [\alpha] [\text{List } \alpha] (\text{cons } [\alpha]) l r \\
\text{reverse} &= \Lambda \alpha :: *. \\
&\quad \lambda l : \text{List } \alpha. \\
&\quad \text{foldList } [\alpha] [\text{List } \alpha \rightarrow \text{List } \alpha] \\
&\quad (\lambda x : \alpha. \lambda k : \text{List } \alpha \rightarrow \text{List } \alpha. \lambda xs : \text{List } \alpha. k (\text{cons } [\alpha] x xs)) \\
&\quad (\lambda xs : \text{List } \alpha. xs) l (\text{nil } [\alpha]) \\
\text{Option} &= \lambda \alpha :: *. \forall \varphi :: * \Rightarrow *. \varphi \alpha \rightarrow (\alpha \rightarrow \varphi \alpha) \rightarrow \varphi \alpha \\
\text{none} &= \Lambda \alpha :: *. \Lambda \varphi :: * \Rightarrow *. \lambda \text{none} : \varphi \alpha. \lambda \text{some} : \alpha \rightarrow \varphi \alpha. \text{none} \\
\text{some} &= \Lambda \alpha :: *. \lambda x : \alpha. \\
&\quad \Lambda \varphi :: * \Rightarrow *. \lambda \text{none} : \varphi \alpha. \lambda \text{some} : \alpha \rightarrow \varphi \alpha. \\
&\quad \text{some } x \\
\text{foldOption} &= \Lambda \alpha :: *. \Lambda \beta :: *. \\
&\quad \lambda s : \alpha \rightarrow \beta. \lambda n : \beta. \lambda o : \text{Option } \alpha. \\
&\quad o [\lambda \gamma :: *. \beta] n s
\end{aligned}$$

Figure 1: Definitions in System $F\omega$